Star exercises - series I Languages, automata and computations II

Deadline: January 11, 2019 By email to: wczerwin@mimuw.edu.pl

1. UFA on finite words

A finite automaton \mathcal{A} is unambiguous if for every word accepted by \mathcal{A} there is exactly one accepting run of it. Decide whether there is a polynomial P such that for every nondeterministic finite automata \mathcal{A} and \mathcal{B} over Σ with at most nstates, such that $L(\mathcal{B}) = \Sigma^* \setminus L(\mathcal{A})$, there is an unambiguous finite automaton \mathcal{C} accepting $L(\mathcal{A})$ with at most P(n) states. **Remark:** We don't know the solution.

2. UFA on infinite trees

Decide the same question as in exercise 1., but for Büchi automata over infinite words. **Remark:** We don't know the solution.

3. Top-down tree automata

Is the following problem decidable: given a regular language of finite trees over some ranked alphabet, is this language recognizable by some deterministic topdown tree automaton?

4. Distance automata with more counters

Consider the following extension of a distance automaton. Instead of having a set of costly transitions, we have two counters $\{1,2\}$ and each transition is labelled by an instruction from the following toolkit:

- do nothing;
- increment counter 1;
- reset counter 1;
- reset counter 1 and increment counter 2;
- reset both counters.

The value of a run is the biggest value attained by any counter. Prove that limitedness is decidable for these automata, using the limitedness game.

5. Separation

Prove that the following problem is decidable:

- Input: Regular word languages $L, K \subseteq \Sigma^*$, given say by deterministic automata.
- Question: Is there a language of star height 1 which contains L but is disjoint with K? A language of star height 1 is a language which can be defined by a regular expression, without complement, where the Kleene star is allowed, but it cannot be nested.

As a hint, use the previous exercise.