The (Almost) Complete Guide to Tree Pattern Containment

Wojciech Czerwiński

Wim Martens

Paweł Parys

Marcin Przybyłko

Tree Pattern

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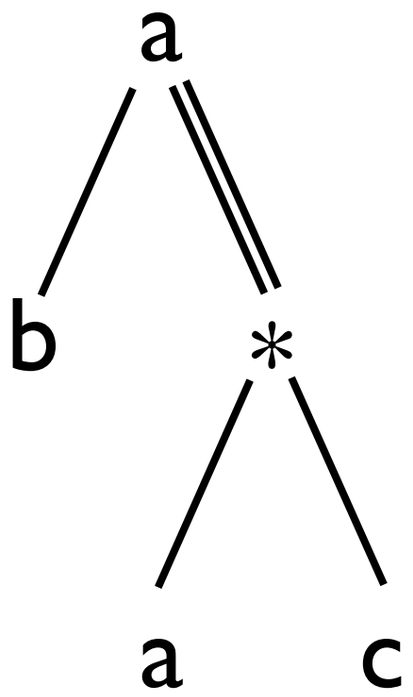

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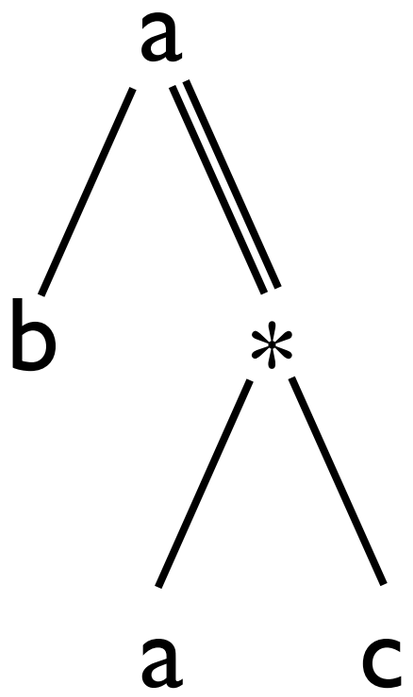

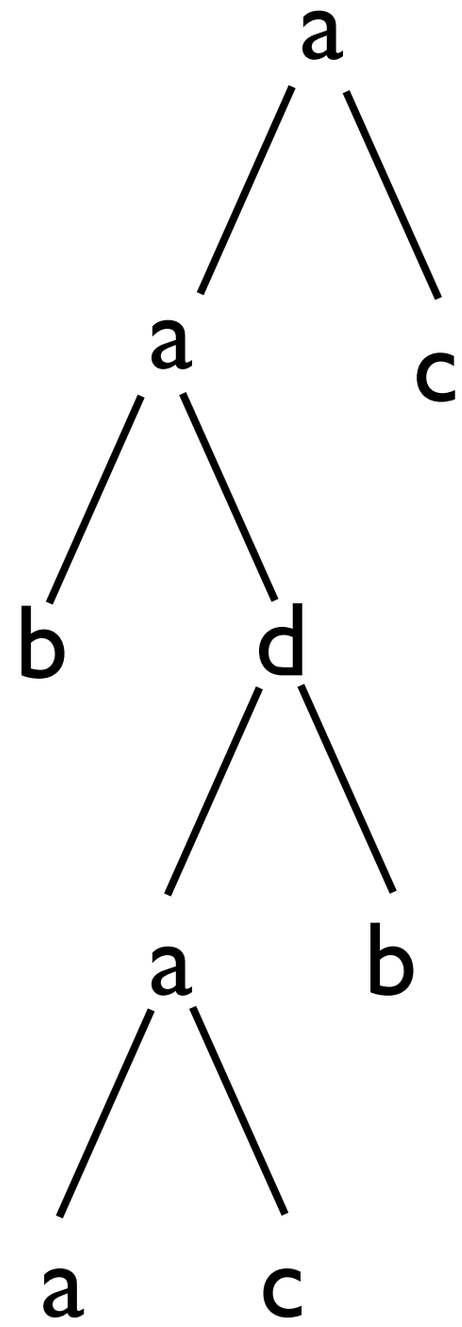

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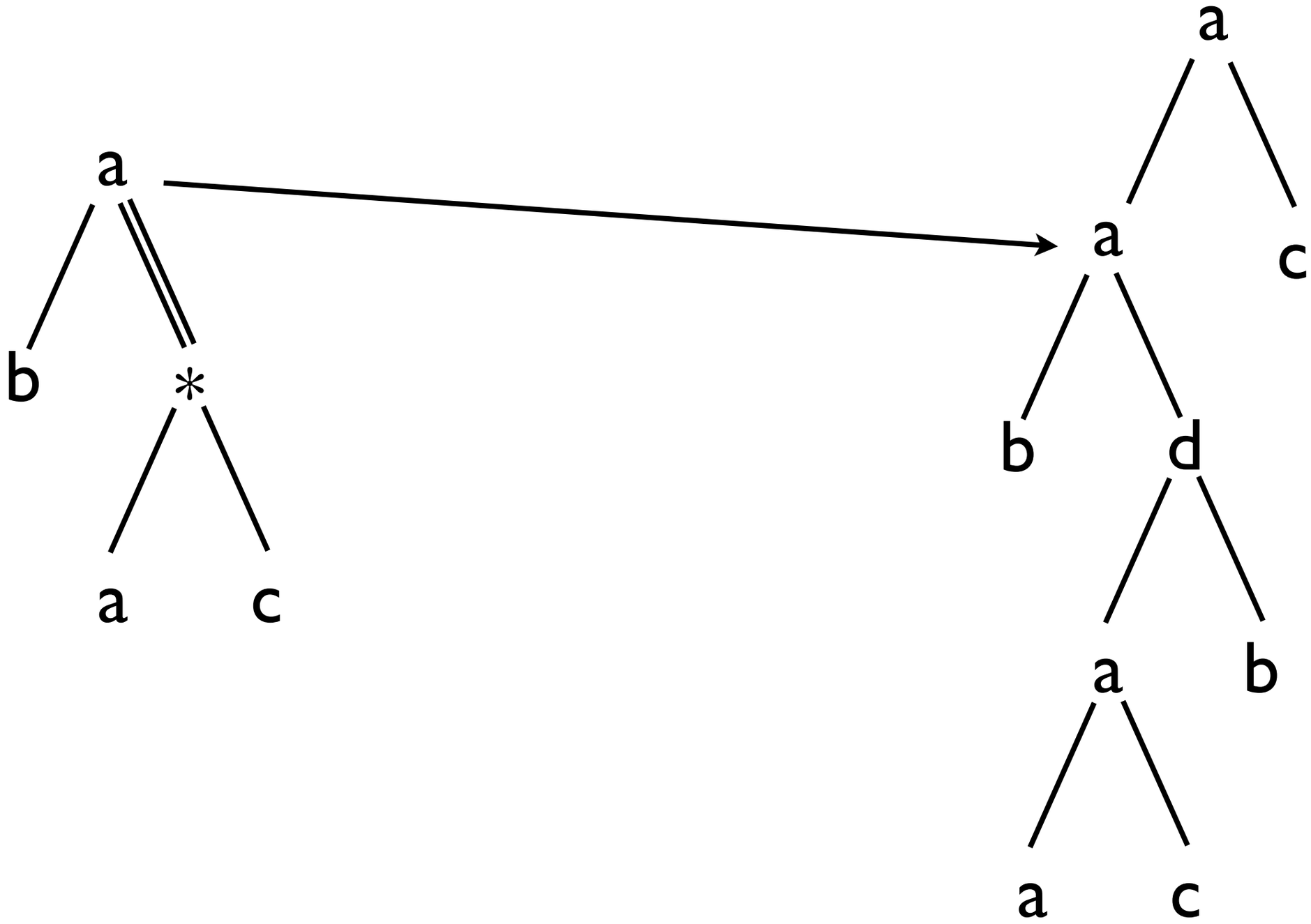

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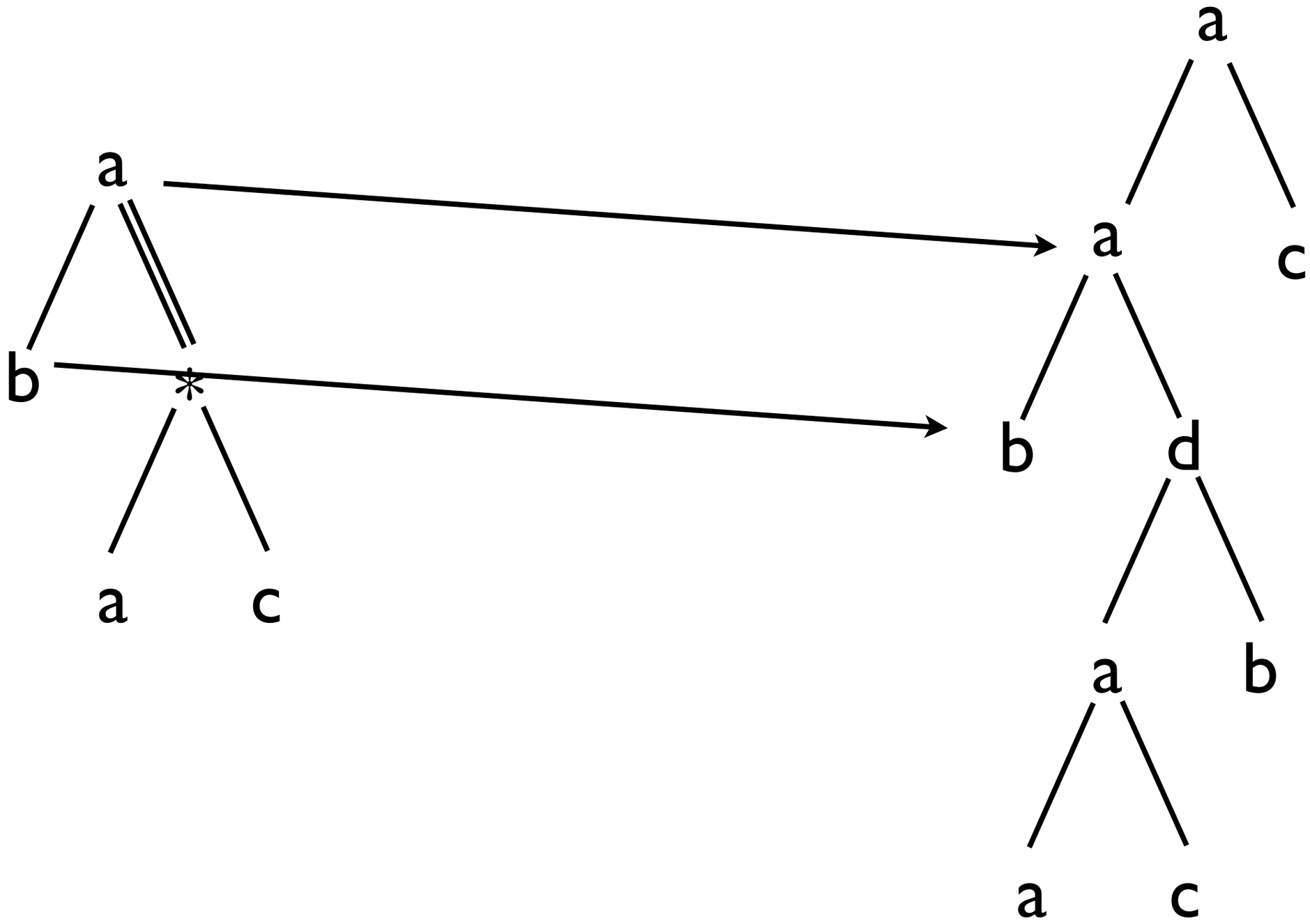

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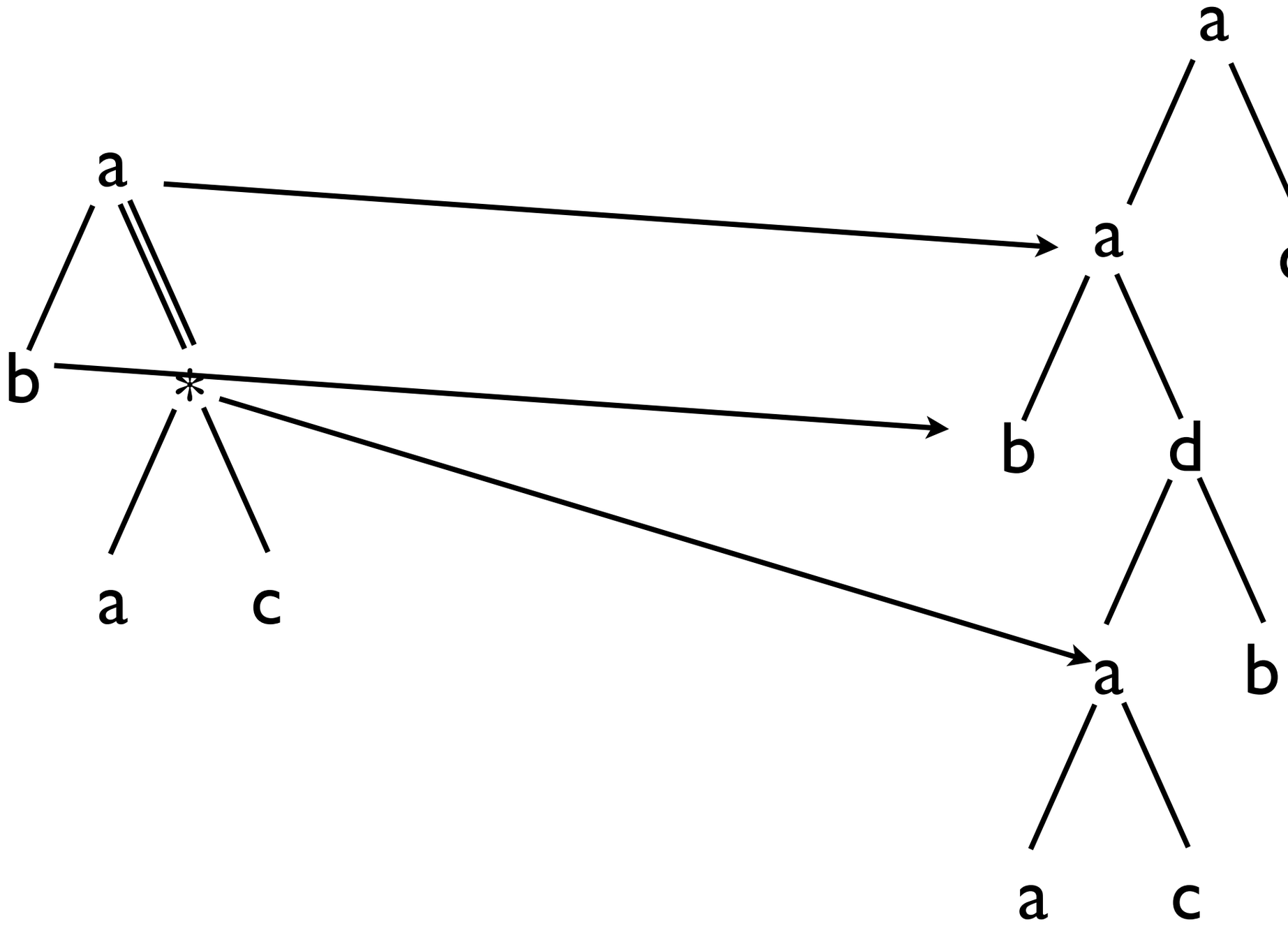

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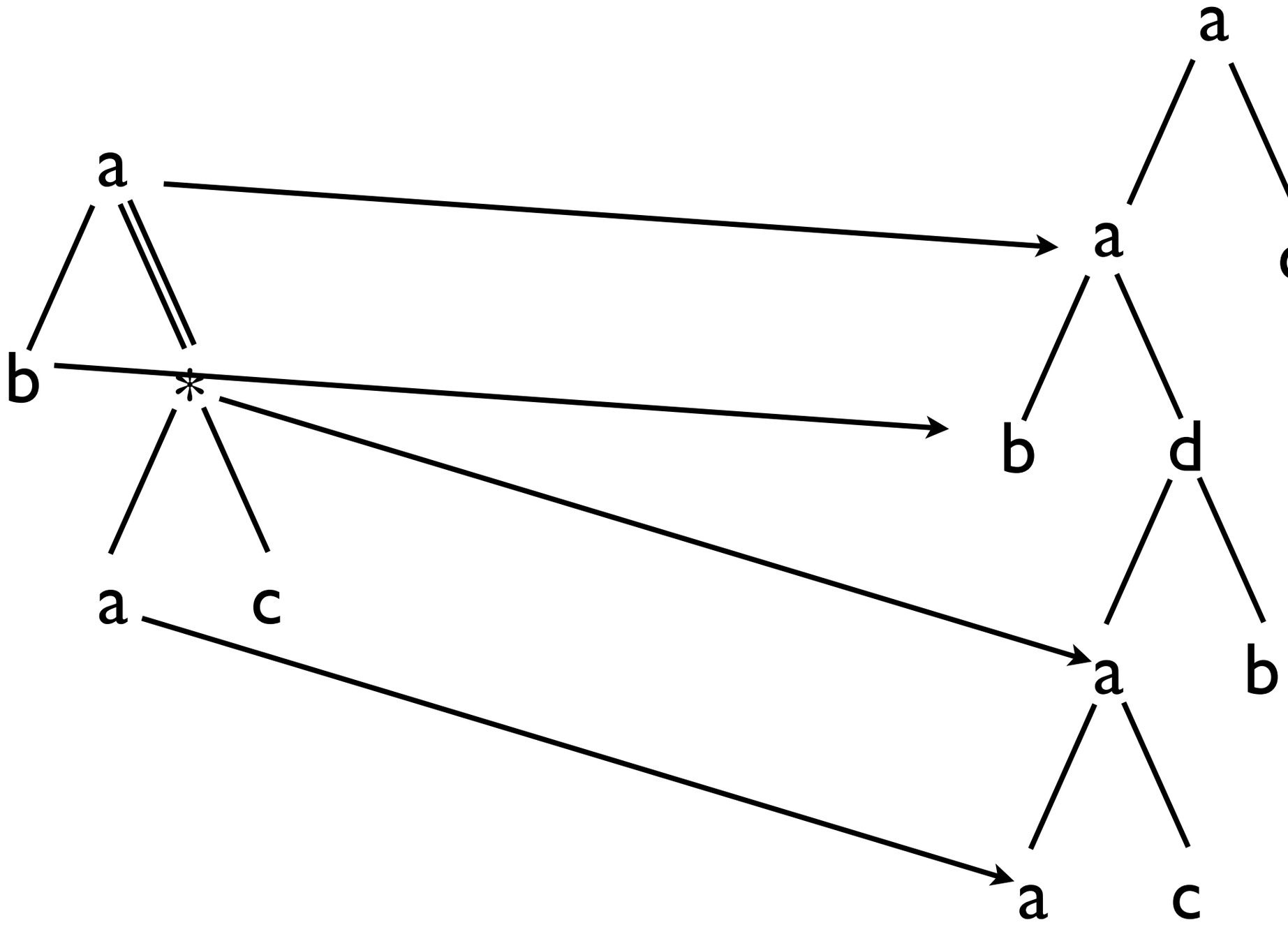

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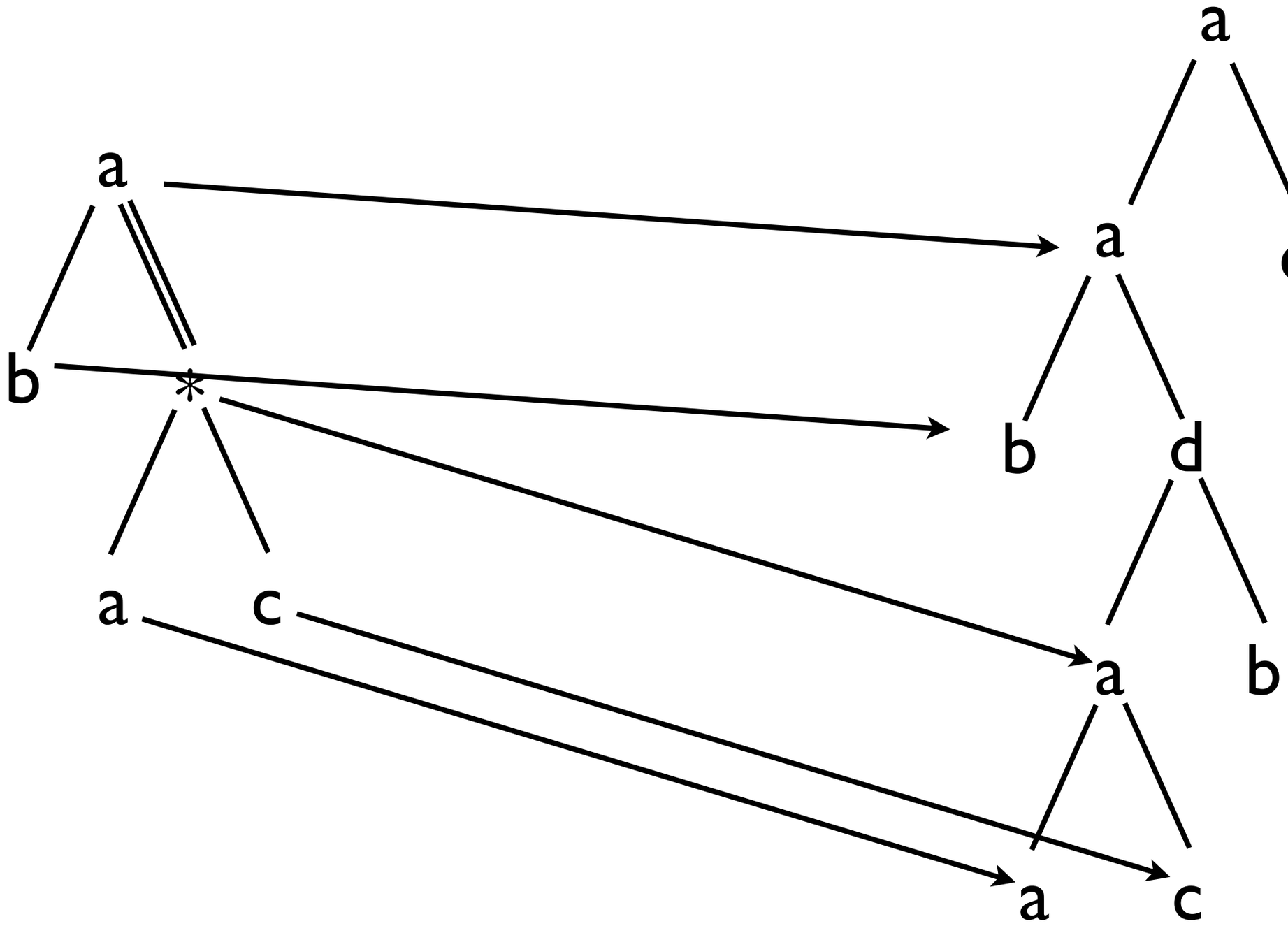

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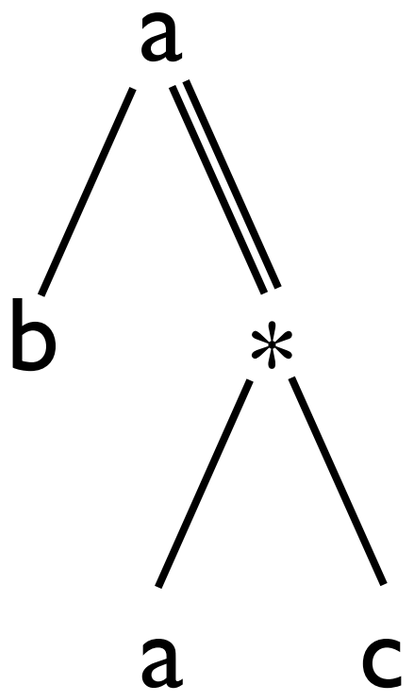

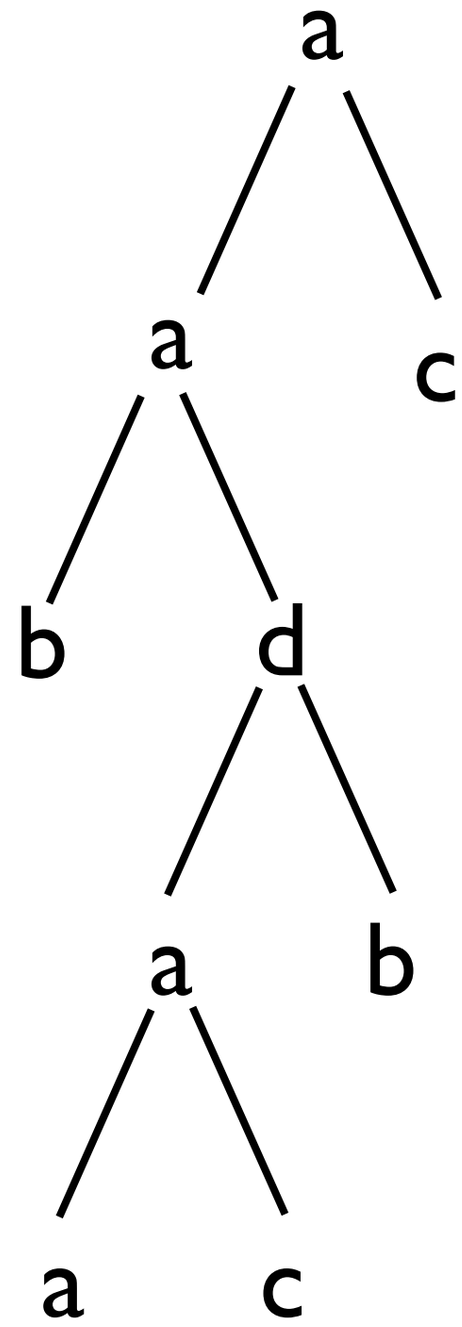

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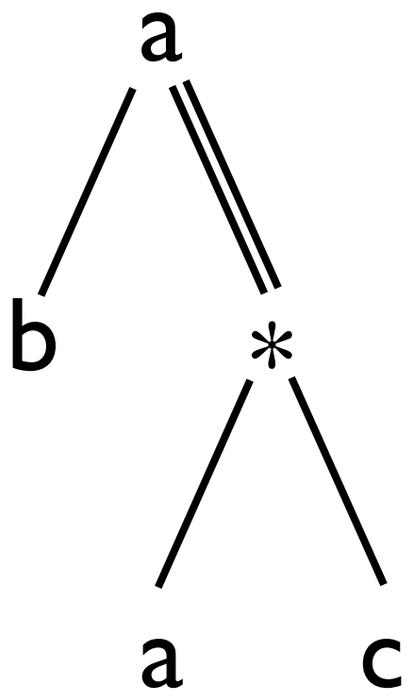
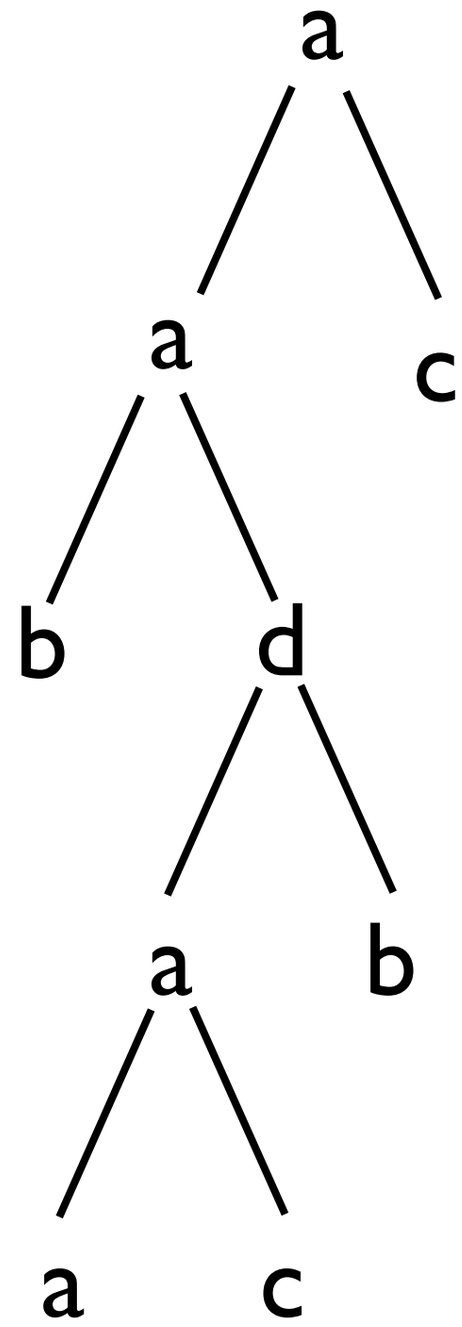


$L_w(p)$



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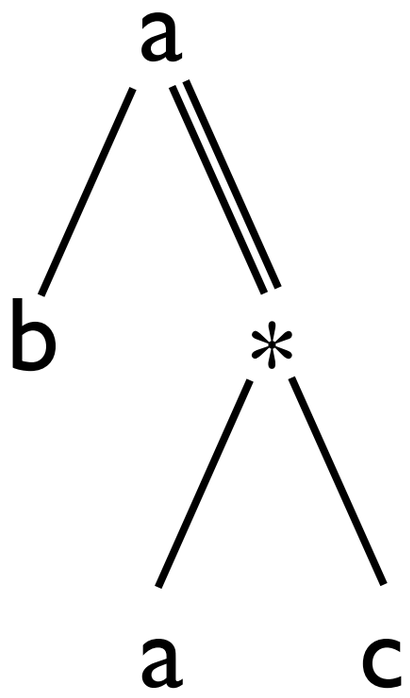

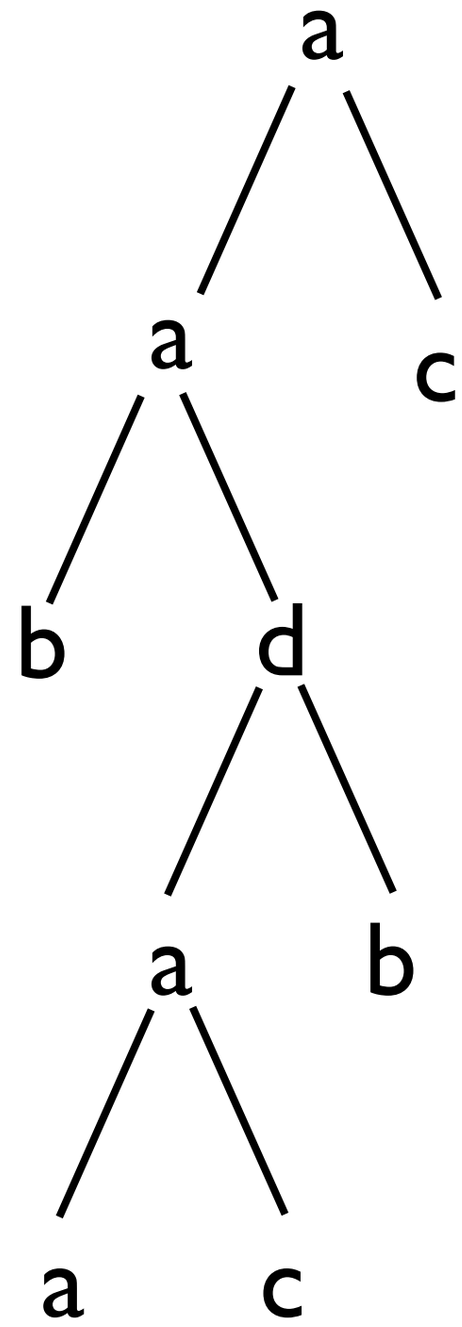

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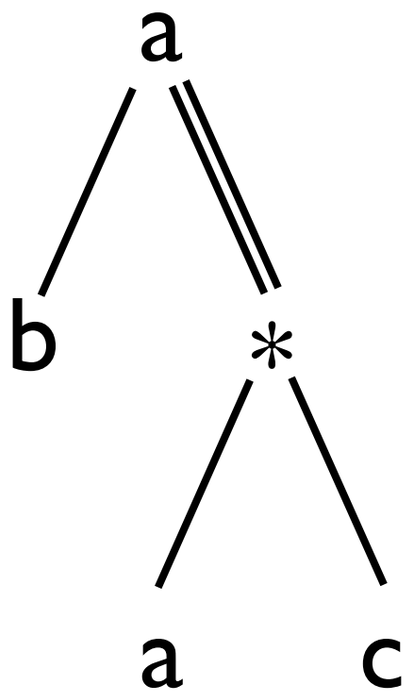

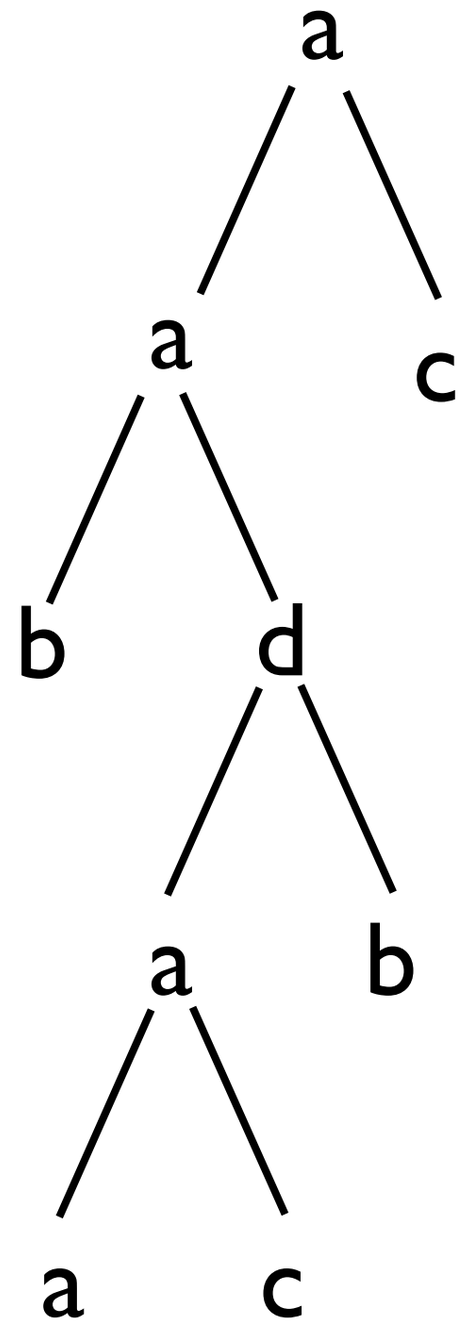

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does p enforce q among trees in L(d)?

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- (some) results transfer to graph structured data

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Goal: uniform overview, solve hard cases,
track the PTIME borderline

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- a lot of mess in the landscape

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Technique: new triomino-tiling method

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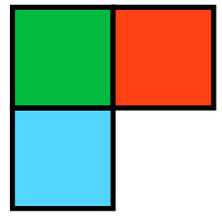

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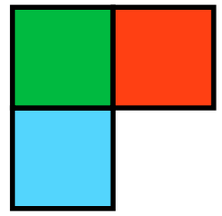

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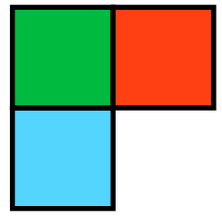

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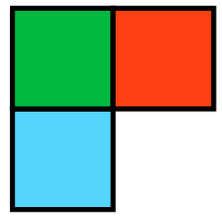

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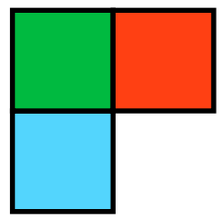

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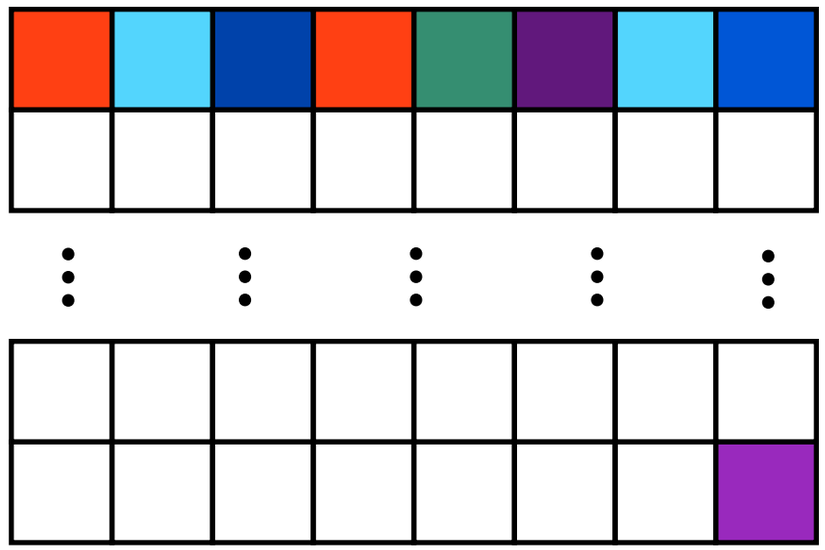

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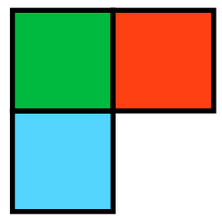

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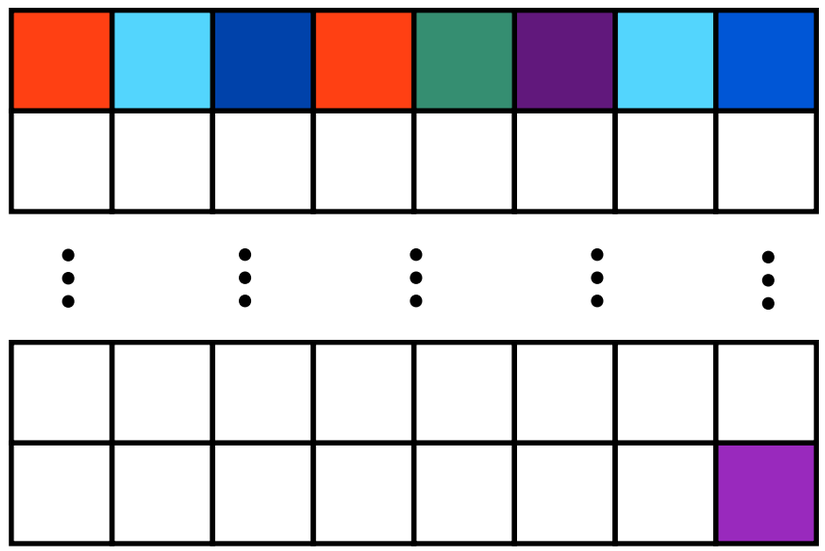

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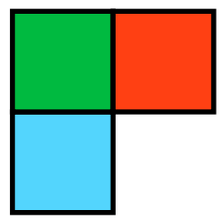

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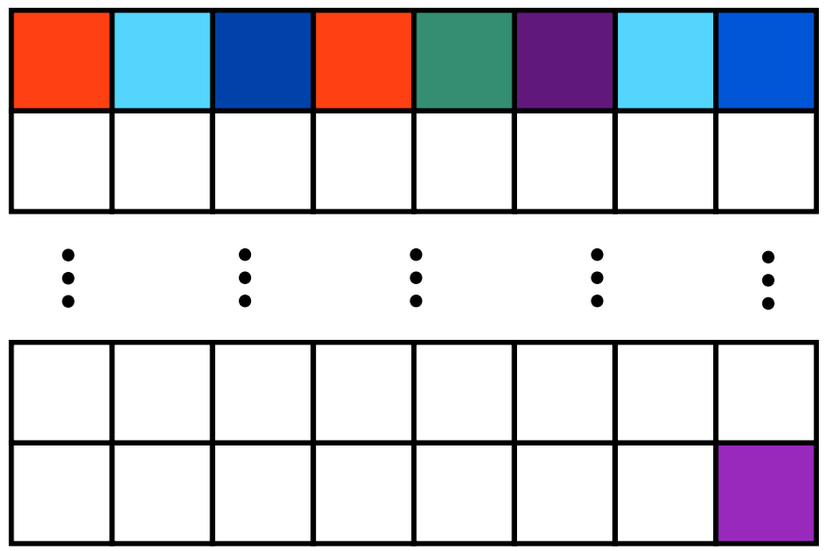

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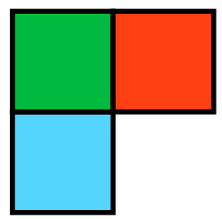

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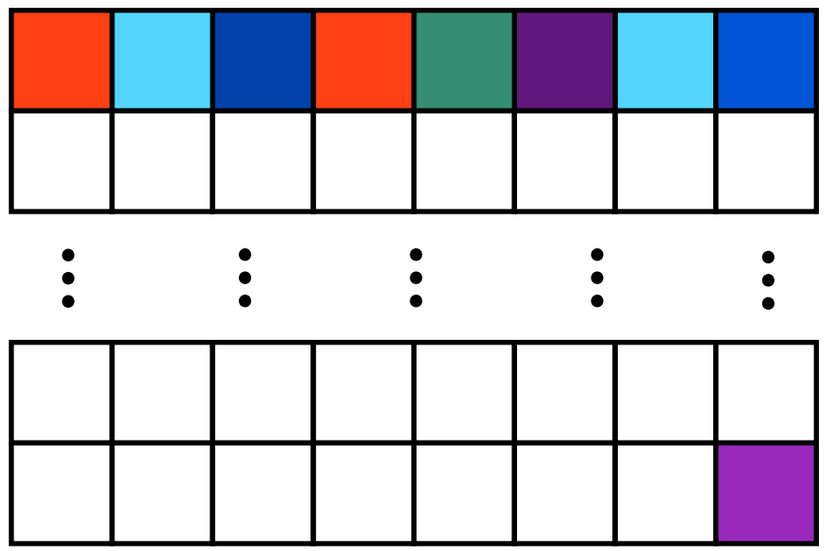

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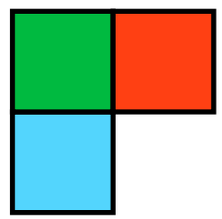

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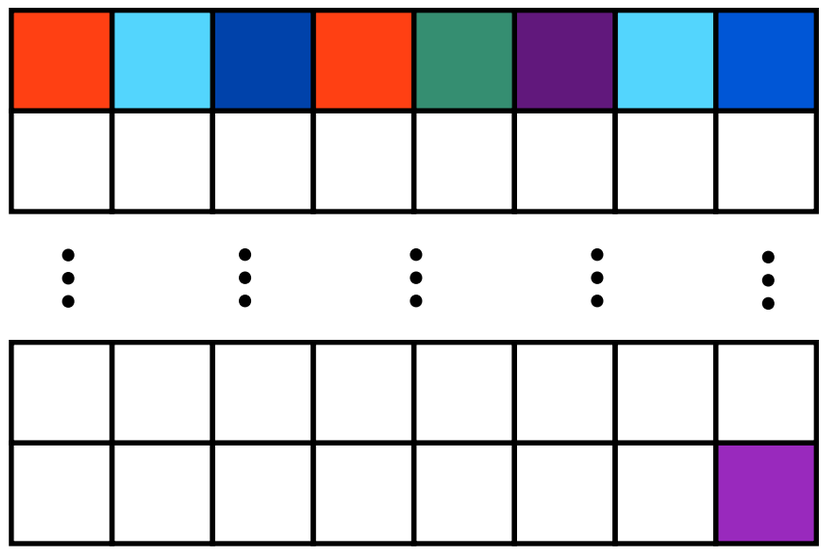

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Right pattern and DTD: implement forbiden triominos via forbiding a******b

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open for a decade - now complete classification without DTD

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Message: such a uniform technique works in all easy cases

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Open question

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- inclusion of PQ in TPQ(/,//) with DTD

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- PTIME or not?

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- PTIME or not?
- probably hard to solve

Thank you!