

Reachability Problem for Fixed Dimensional Vector Addition Systems with States

Wojciech Czerwiński

Sławomir Lasota

Ranko Lazić

Jérôme Leroux

Filip Mazowiecki

Towards Understanding
the Reachability Problem
for
Vector Addition Systems
with States

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Plan

Plan

- definition

Plan

- definition
- state of art

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- definition
- state of art
- motivation

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- definition
- state of art
- motivation
- our results

Plan

- definition
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- 3D example

Plan

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- 3D example
- 4D example

Vector Addition Systems with States (VASS)

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Reachability problem

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Reachability problem

Input: finite set of transitions T in $Q \times \mathbb{Z}^d \times Q$
source s , target t in $Q \times \mathbb{N}^d$

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by finitely many transitions from T
inside the positive quadrant?

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Assume s and t is 0 vector

State of art

State of art

General VASSes

State of art

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Ackermann (Leroux, Schmitz `19)

State of art

General VASSes

Ackermann (Leroux, Schmitz `19)

Tower-hard (CLLLM `19)

State of art

General VASSes

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Fixed dimension d-VASSes

State of art

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Fixed dimension d-VASSes

F_{d+4} (Leroux, Schmitz `19)

State of art

General VASSes

Ackermann (Leroux, Schmitz `19)

Tower-hard (CLLLM `19)

Fixed dimension d -VASSes

F_{d+4} (Leroux, Schmitz `19)

$(d-13)$ -ExpSpace-hard (CLLLM `19)

Small dimension

Small dimension

unary

binary

1-VASS

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--

2-VASS

--

--

Small dimension

unary

binary

1-VASS

NL

2-VASS

Small dimension

unary

binary

1-VASS

NL

NP

(HKQW `10)

2-VASS

Small dimension

unary

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NL

NP
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2-VASS

PSpace
(BFGHM '15)

Small dimension

unary

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← short paths

Small dimension

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← short paths

linear path schemes

Small dimension

unary

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← short paths

linear path schemes

$u_0 (v_1)^* u_1 (v_2)^* \dots u_{k-1} (v_k)^* u_k$

Small dimension

unary

binary

1-VASS

NL

NP
(HKQW `10)

2-VASS

NL
(ELT `16)

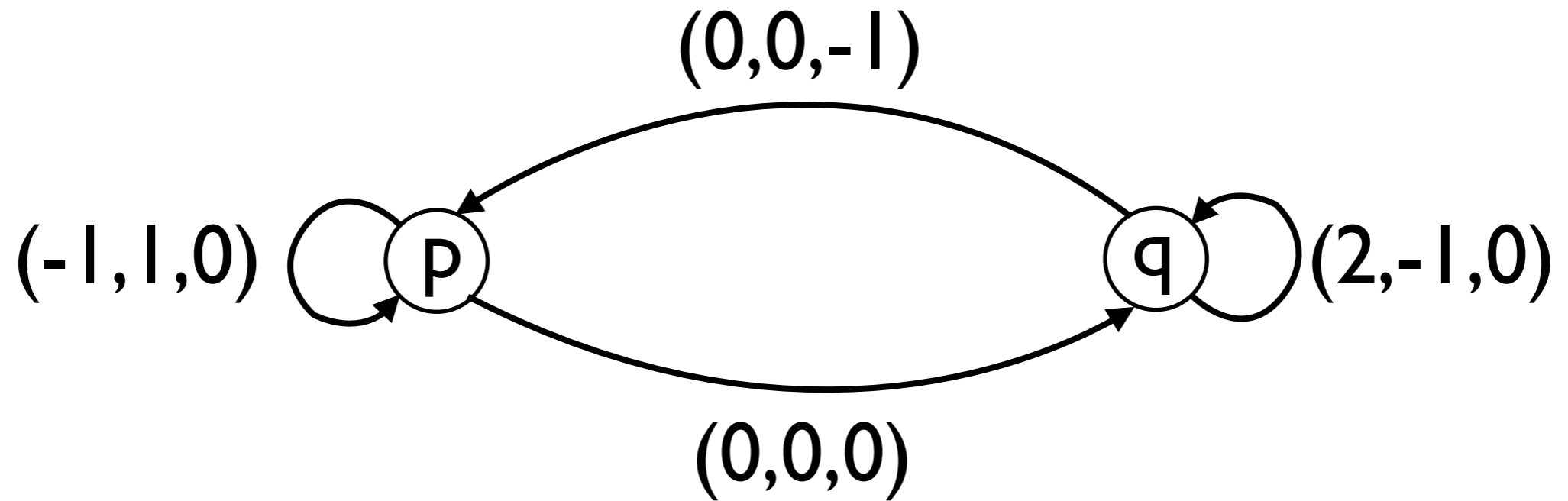
PSpace
(BFGHM `15)

← short paths
linear path schemes

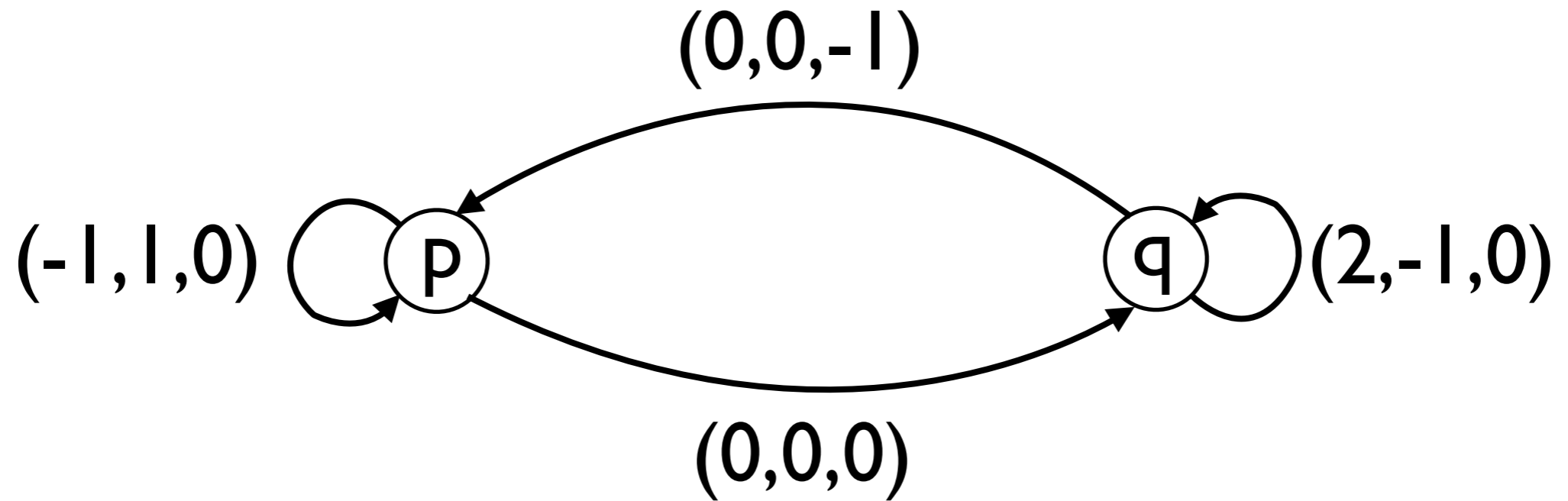
$u_0 (v_1)^* u_1 (v_2)^* \dots u_{k-1} (v_k)^* u_k$

Hopcroft-Pansiot example

Hopcroft-Pansiot example

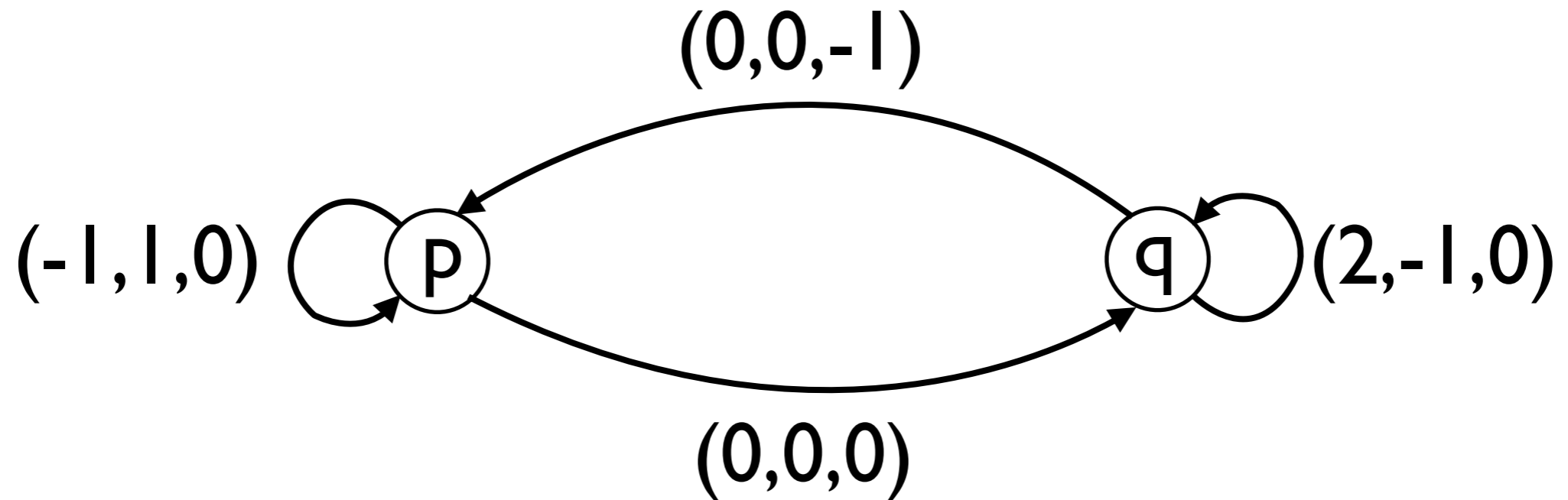


Hopcroft-Pansiot example



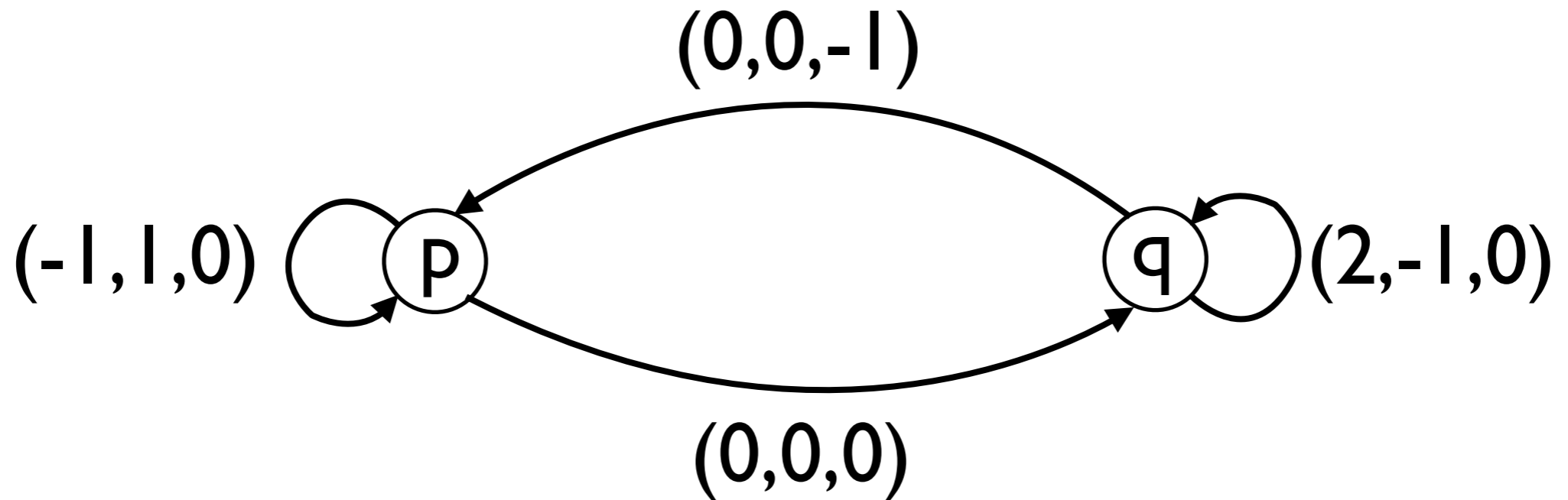
$p(k, 0, n)$

Hopcroft-Pansiot example



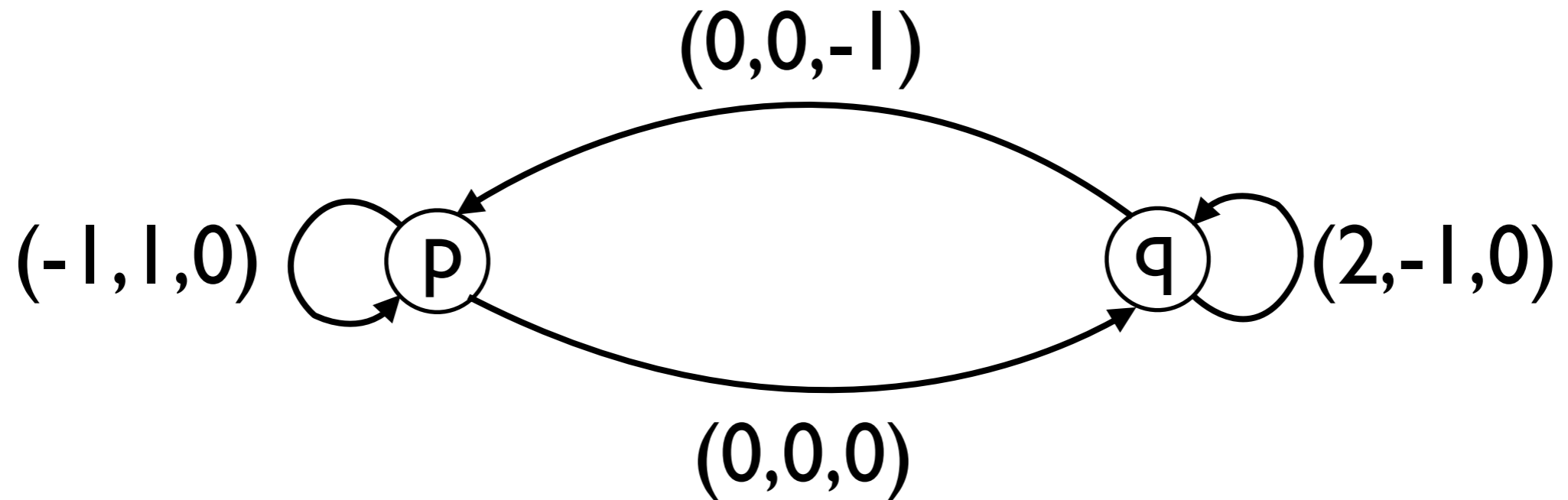
$p(k, 0, n) \longrightarrow p(0, k, n)$

Hopcroft-Pansiot example



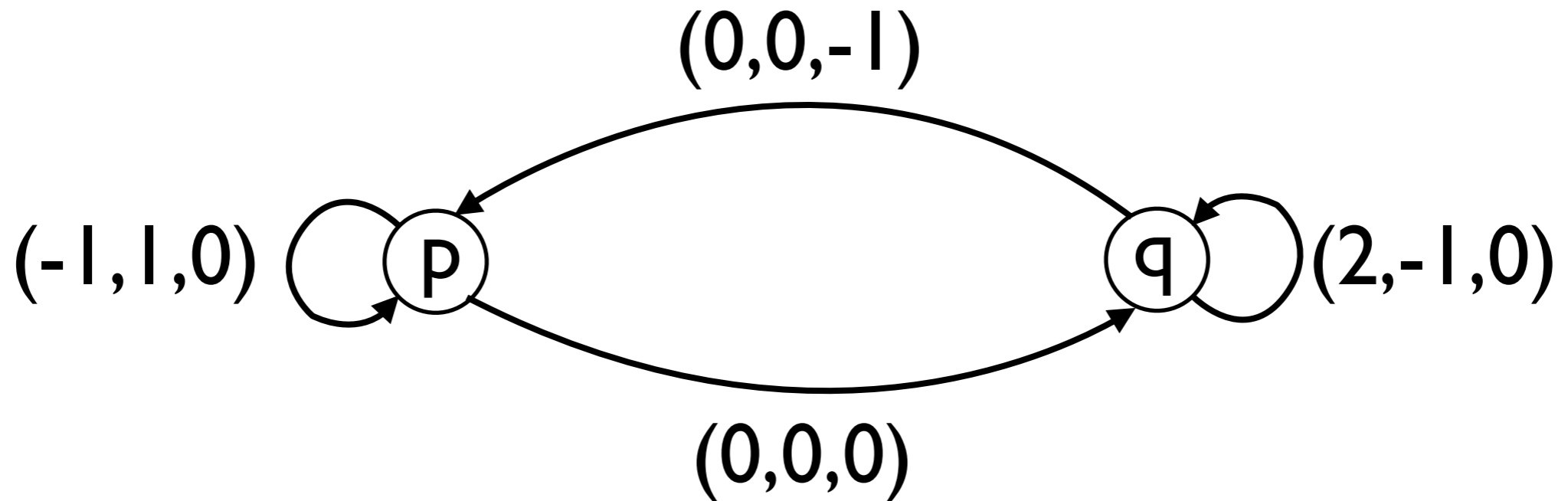
$p(k, 0, n) \longrightarrow p(0, k, n) \longrightarrow q(0, k, n)$

Hopcroft-Pansiot example



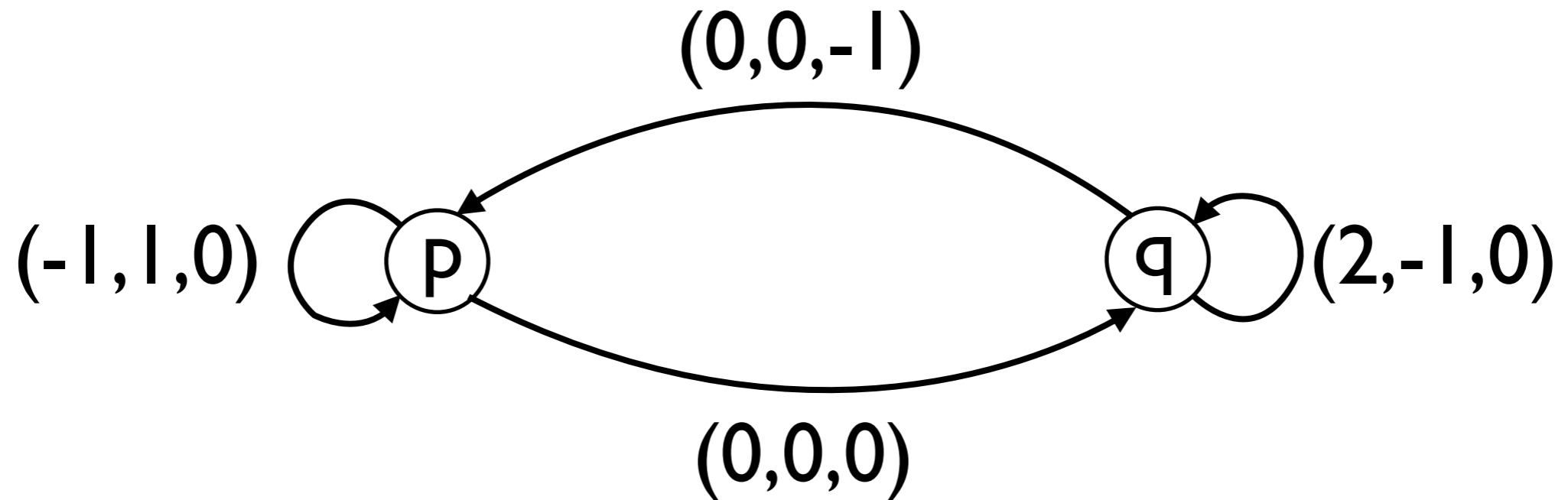
$$p(k, 0, n) \longrightarrow p(0, k, n) \longrightarrow q(0, k, n) \longrightarrow q(2k, 0, n)$$

Hopcroft-Pansiot example



$p(k, 0, n) \longrightarrow p(0, k, n) \longrightarrow q(0, k, n) \longrightarrow q(2k, 0, n) \longrightarrow p(2k, 0, n-1)$

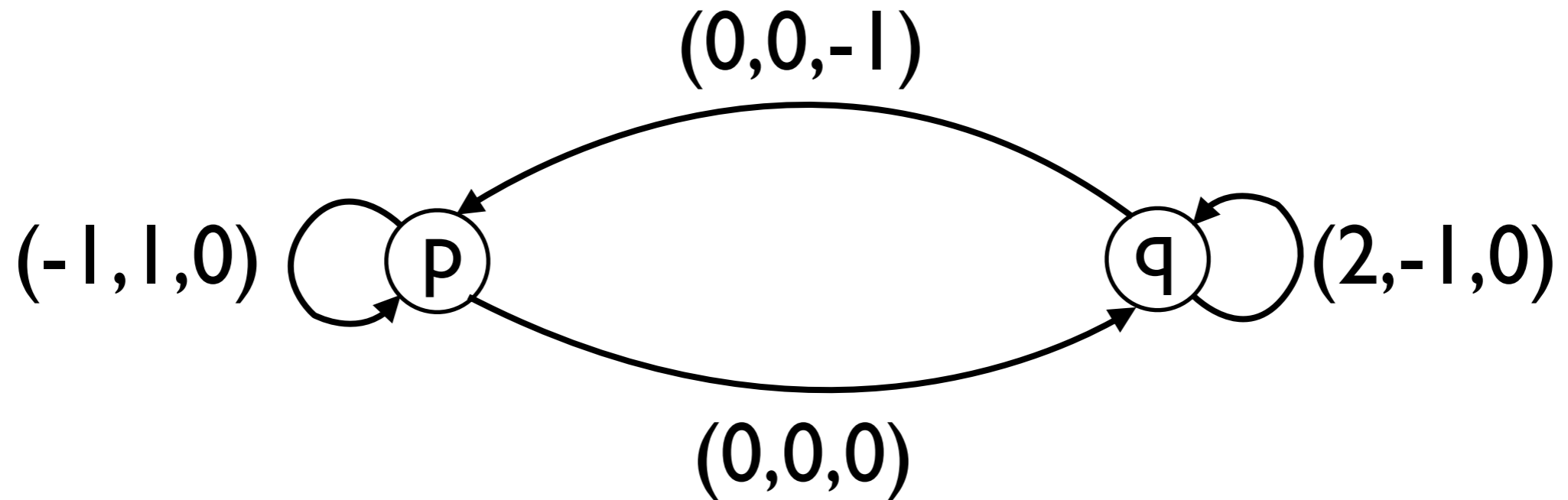
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$p(1, 0, n)$

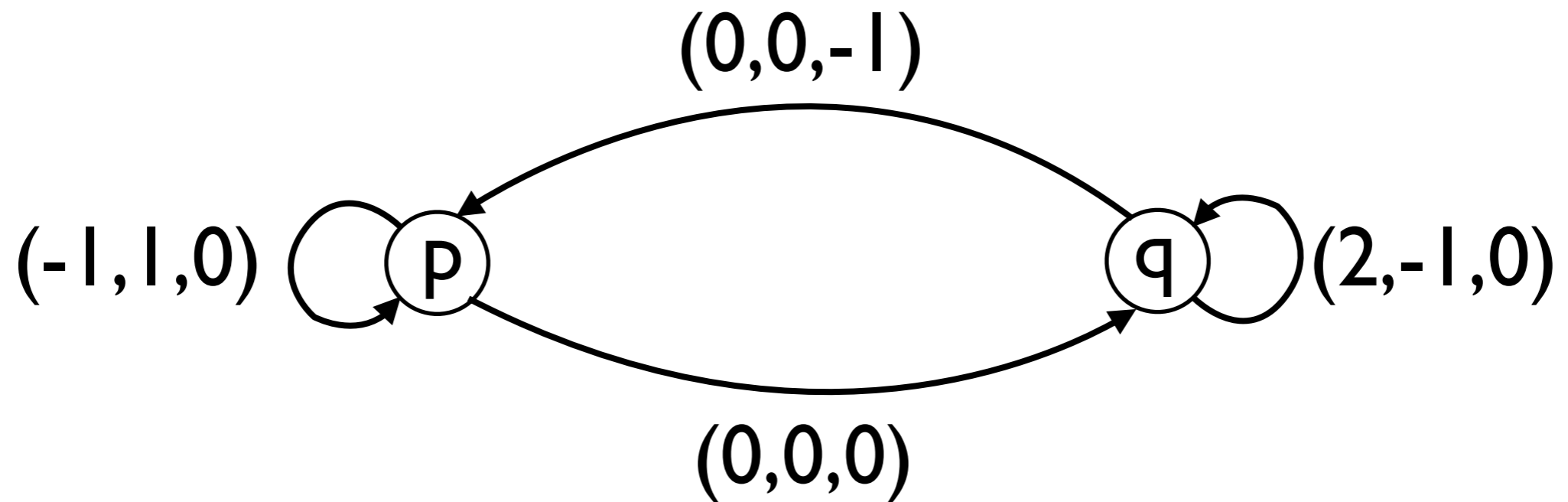
Hopcroft-Pansiot example



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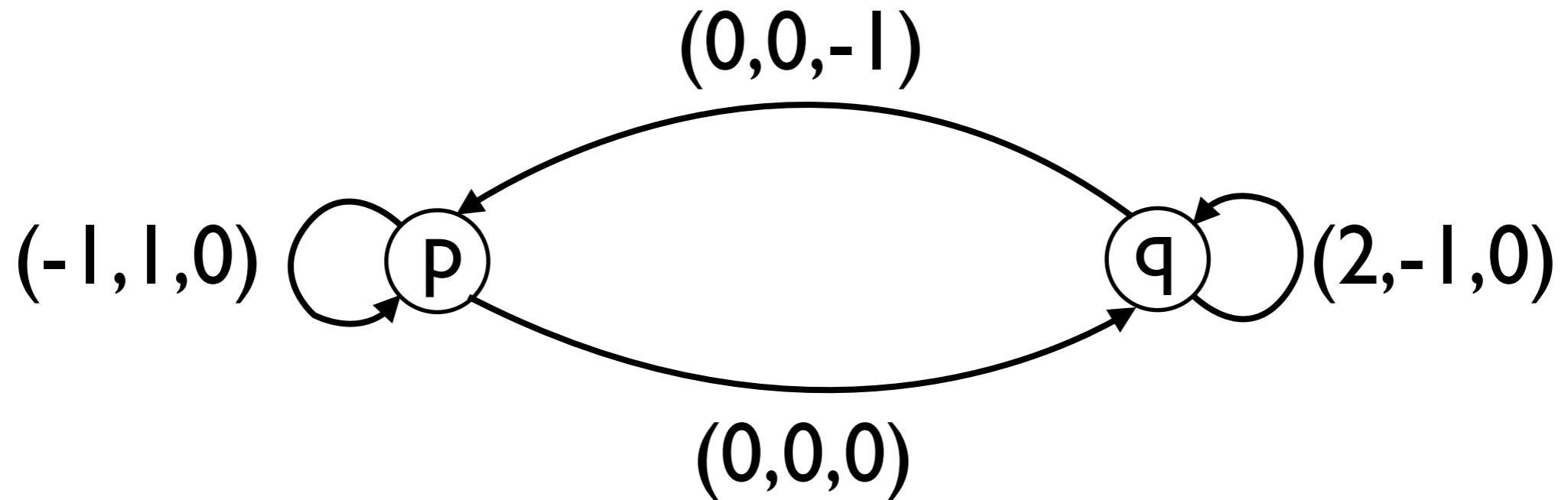
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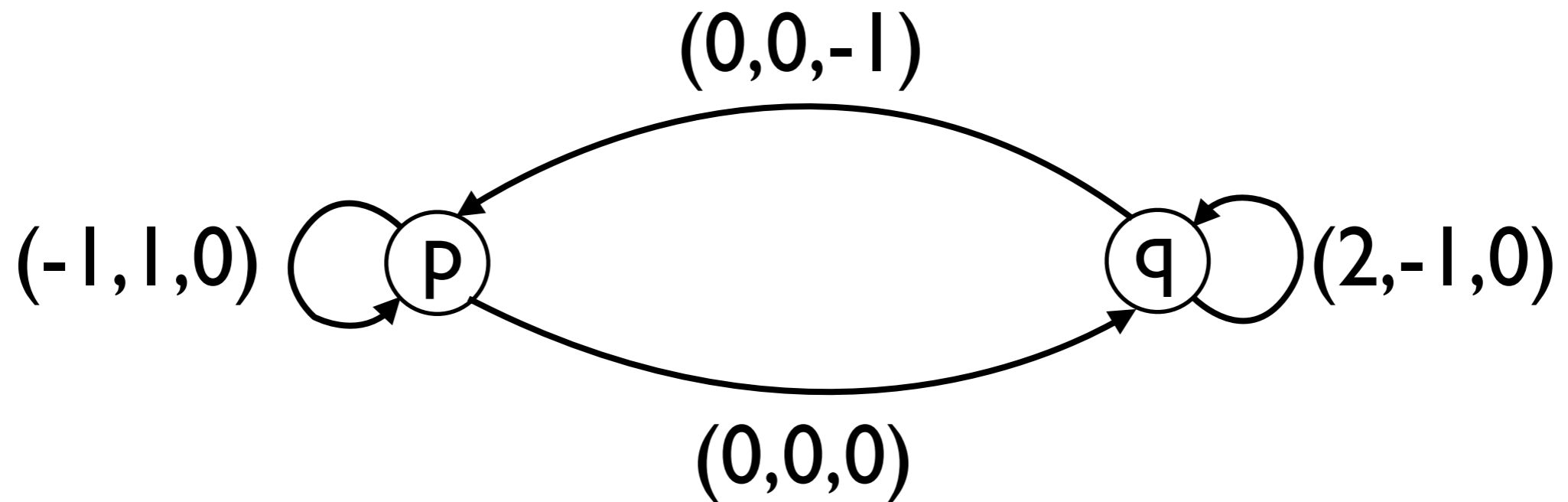
Hopcroft-Pansiot example



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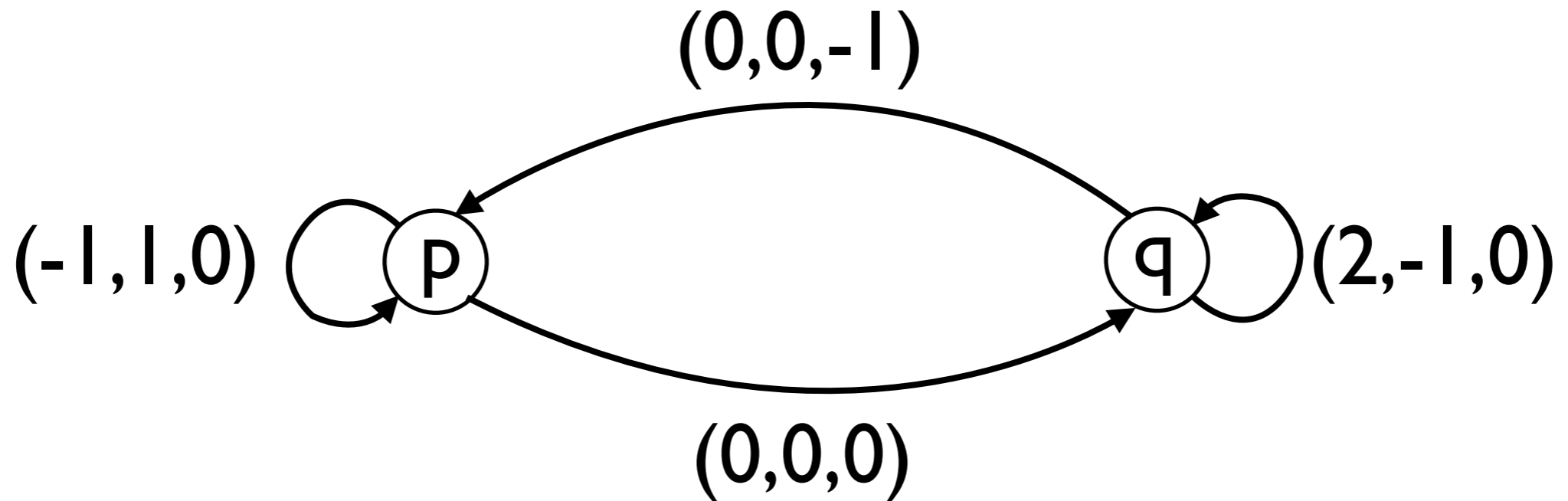
Hopcroft-Pansiot example



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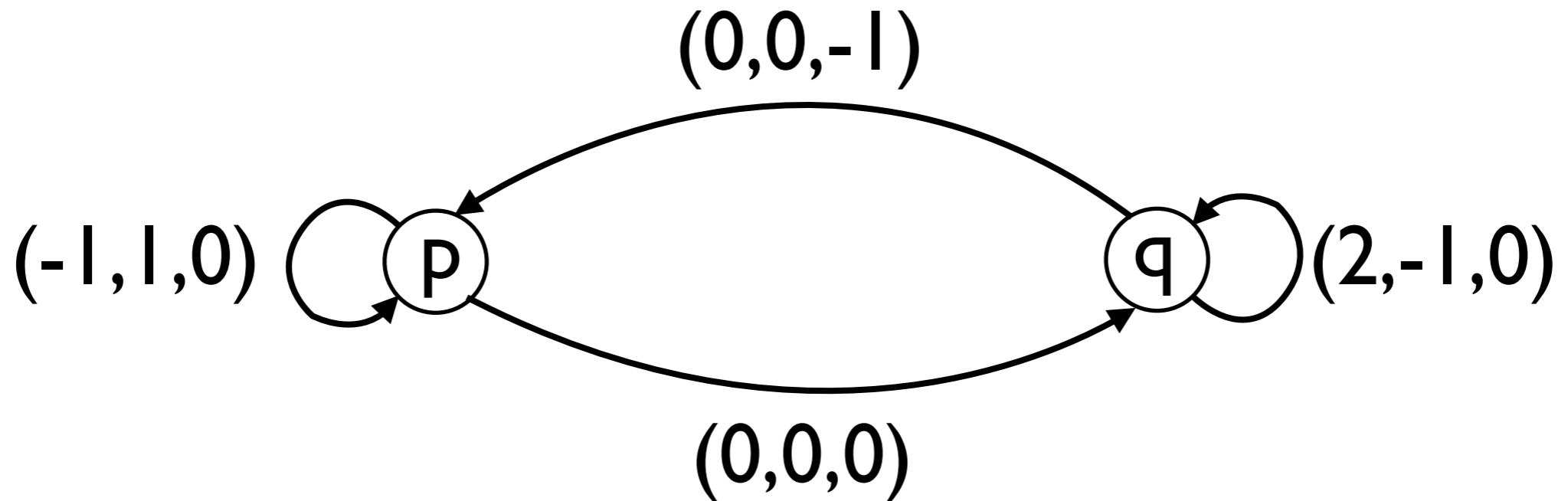
Hopcroft-Pansiot example



$$p(k, 0, n) \longrightarrow p(0, k, n) \longrightarrow q(0, k, n) \longrightarrow q(2k, 0, n) \longrightarrow p(2k, 0, n-1)$$

$$p(1, 0, n) \longrightarrow p(2, 0, n-1) \longrightarrow \dots \longrightarrow p(2^n, 0, 0)$$

Hopcroft-Pansiot example



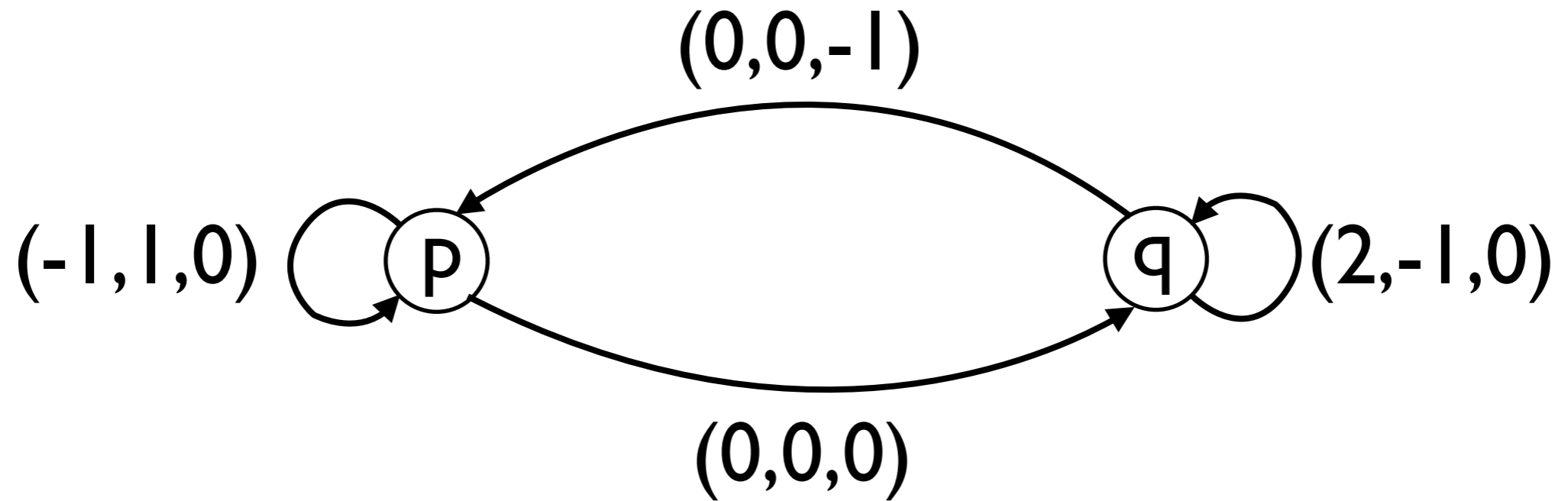
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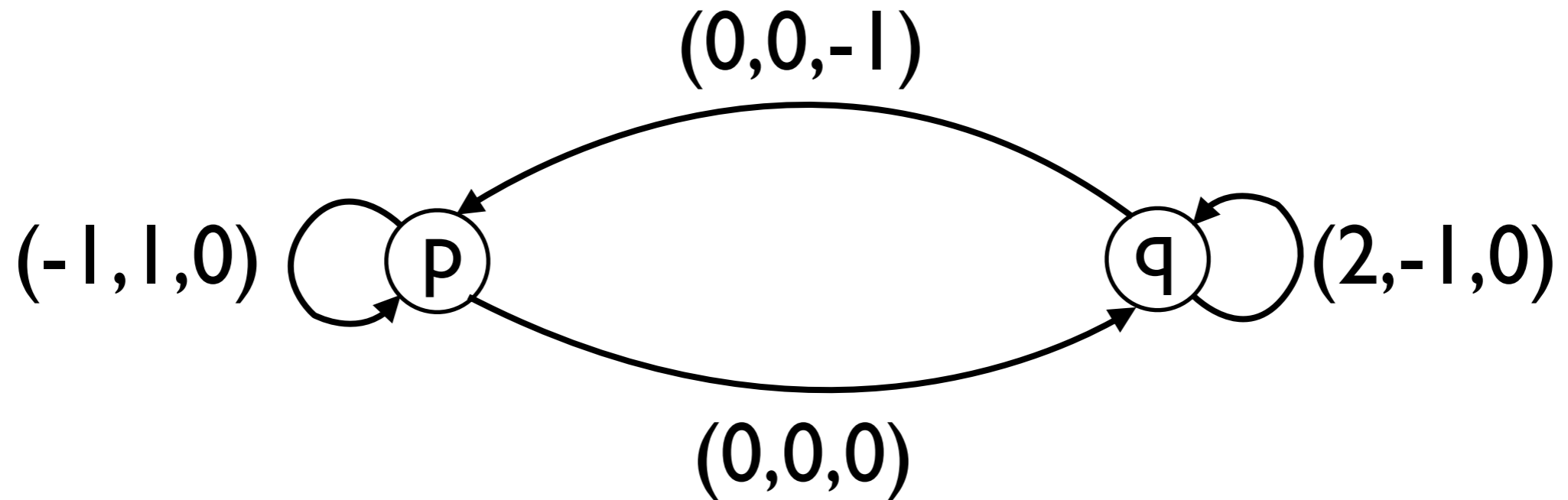
Each $p(x, y, 0)$ for $x+y=2^n$ is reachable

Hopcroft-Pansiot example

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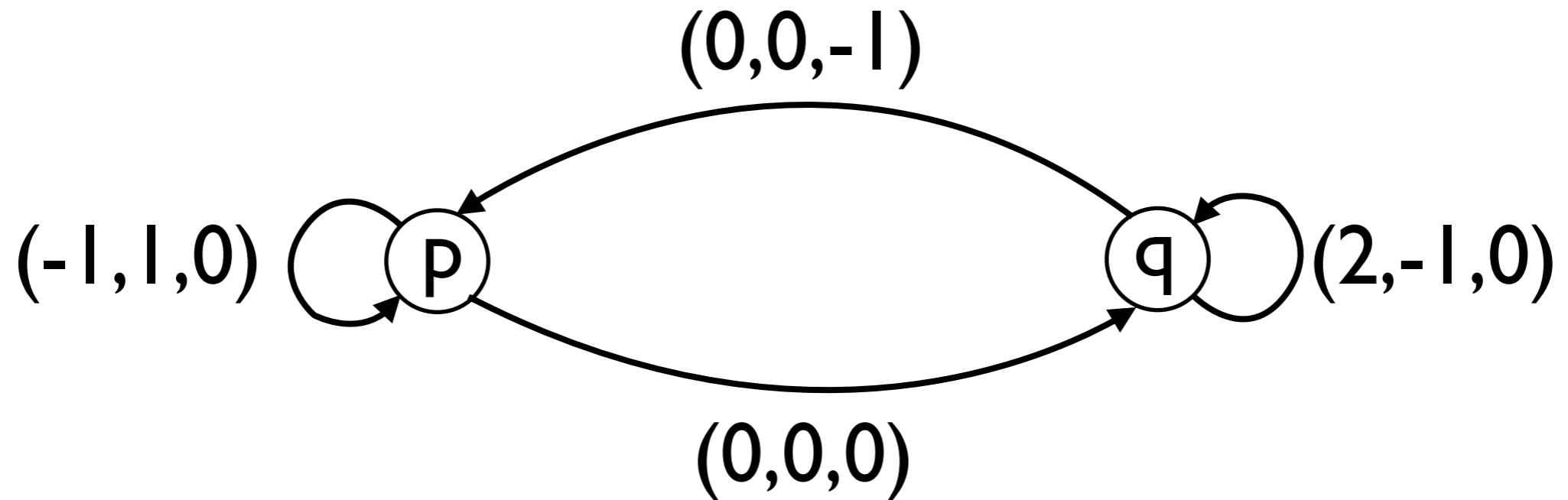


Hopcroft-Pansiot example



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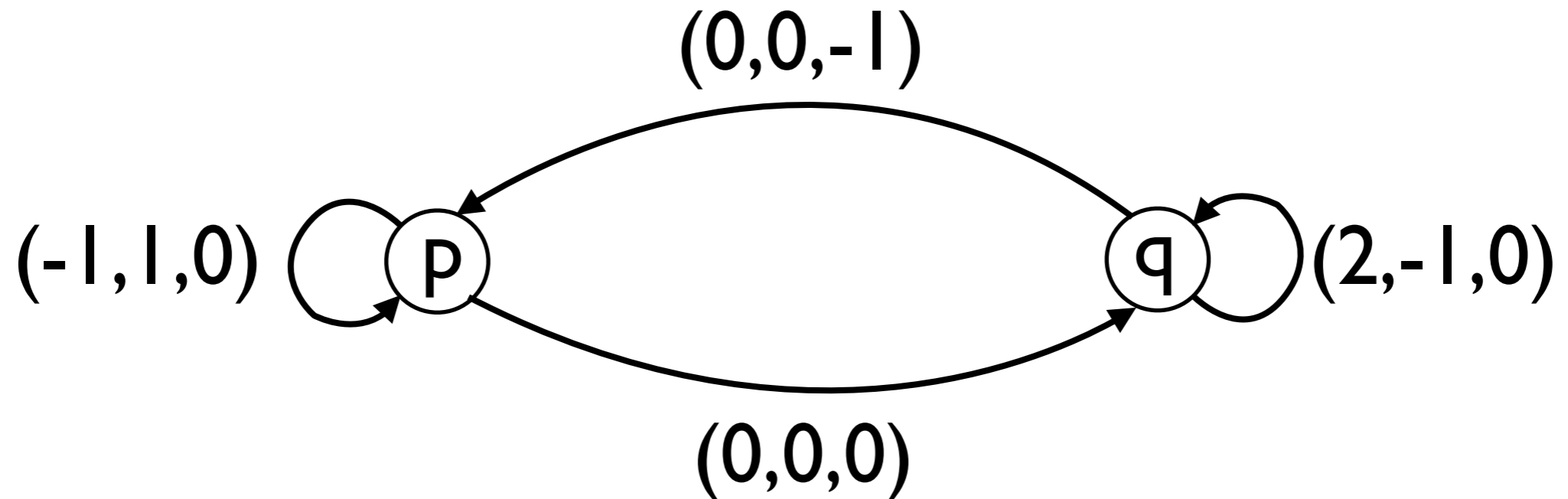
Hopcroft-Pansiot example



Each $p(x, y, 0)$ for $x+y=2^n$ is reachable

Non-semilinear reachability set

Hopcroft-Pansiot example



Each $p(x, y, 0)$ for $x+y=2^n$ is reachable

Non-semilinear reachability set

$$\{(x, y, n) : x+y \leq 2^n\}$$

Motivation

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Reachability problem for finite reachability sets

Motivation

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KLM decomposition useless?

Motivation

Reachability problem for finite reachability sets

KLM decomposition useless?

Reachability for finite 3-VASSes (PSpace? Tower?)

Conjectures

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Conjecture I

In each dimension configuration reachable by a LPS is reachable by a polynomial length path

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$$u_0 (v_1)^* u_1 (v_2)^* \dots u_{k-1} (v_k)^* u_k$$

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Conjecture II

In each dimension reachability by a LPS is in PTime

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Conjecture II

In each dimension reachability by a LPS is in PTime

Conjecture III

In each dimension shortest path is at most exponential (in binary)

Results

Results

Theorem 1

There is a 3-VASS with configuration reachable by a LPS
and with exponential shortest path

Results

Theorem I

There is a 3-VASS with configuration reachable by a LPS
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Theorem II

Reachability by a LPS in dimension 7 is NP-complete

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Theorem II

Reachability by a LPS in dimension 7 is NP-complete

Theorem III

There is a 4-VASS with shortest path of doubly
exponential length

3D example

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There is a 3-VASS with configuration reachable by a LPS
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$$(2 / 1) \cdot (3 / 2) \cdot \dots \cdot (k / k-1) = k$$

3D example

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Idea: enforce that run visits a configuration $p(N, N, 0)$ with N divisible by all primes in $\{1, 2, \dots, k\}$.

3D example

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$$(1,1,0) (1,1,0)^* (-1,0,1)^*(k,0,-(k-1)) \dots (-1,0,1)^*(3,0,-2) \dots$$

3D example

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3D example

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$$\dots (-1,0,1)^*(2,0,-1) (-k,-1,0)$$

$(N,N,0)$ after green part with exponential N

4D example

4D example

Theorem III

There is a 4-VASS with shortest path of doubly exponential length

4D example

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There is a 4-VASS with shortest path of doubly exponential length

Idea: enforce exact multiplications by **fraction** **powers**

4D example

Theorem III

There is a 4-VASS with shortest path of doubly exponential length

Idea: enforce exact multiplications by **fraction** **powers**

Needed: lemma about **fraction** **powers**

Fractional equation

Fractional equation

Lemma

Fractional equation

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For every k there are k fractions

$$1 < a_1 / b_1 \leq \dots \leq a_k / b_k$$

such that

Fractional equation

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$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

Fractional equation

Lemma

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such that

$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

and all a_i , b_i , a and b are at most exponential in k .

VASS building blocks

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$$s(0, 0, 0, 0) \longrightarrow s(K, K, 0, 0)$$

VASS building blocks

$$s(0, 0, 0, 0) \longrightarrow s(K, K, 0, 0) \quad (+1, 0, 0, +1) \text{ in } s$$

VASS building blocks

$$s(0, 0, 0, 0) \longrightarrow s(K, K, 0, 0) \quad (+1, 0, 0, +1) \text{ in } s$$

$$t(K, Ka/b, 0, 0) \longrightarrow t(0, 0, 0, 0)$$

VASS building blocks

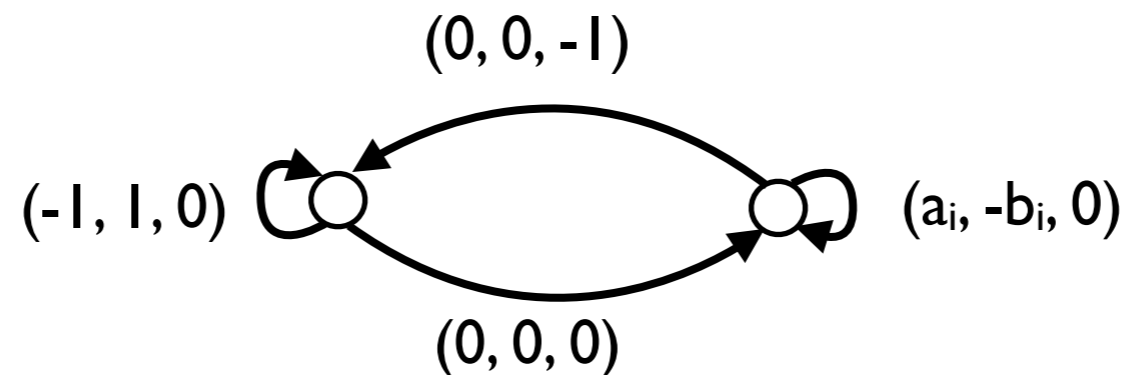
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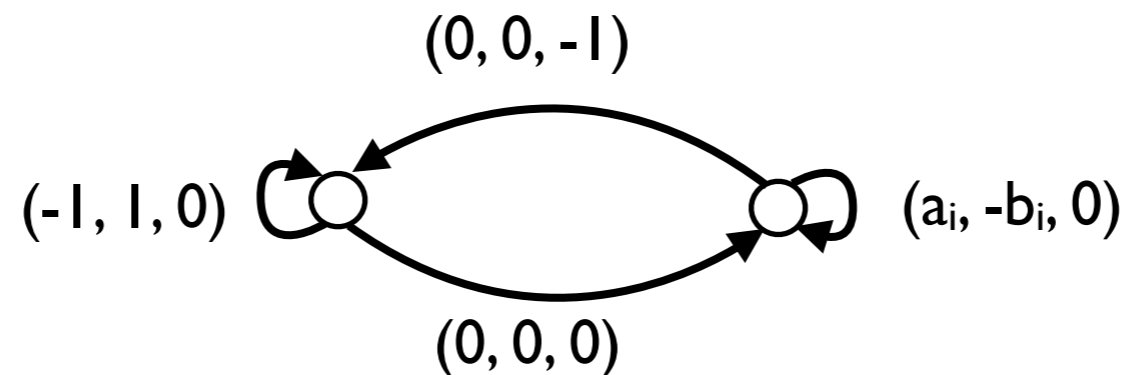
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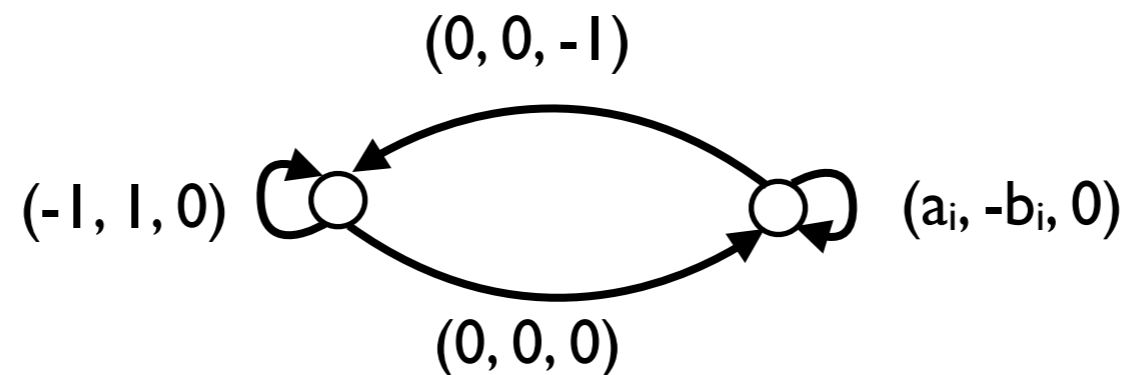


$$(Cb_i^n, 0, n) \longrightarrow (Ca_i b_i^{n-1}, 0, n-1) \longrightarrow \dots \longrightarrow (Ca_i^n, 0, 0)$$

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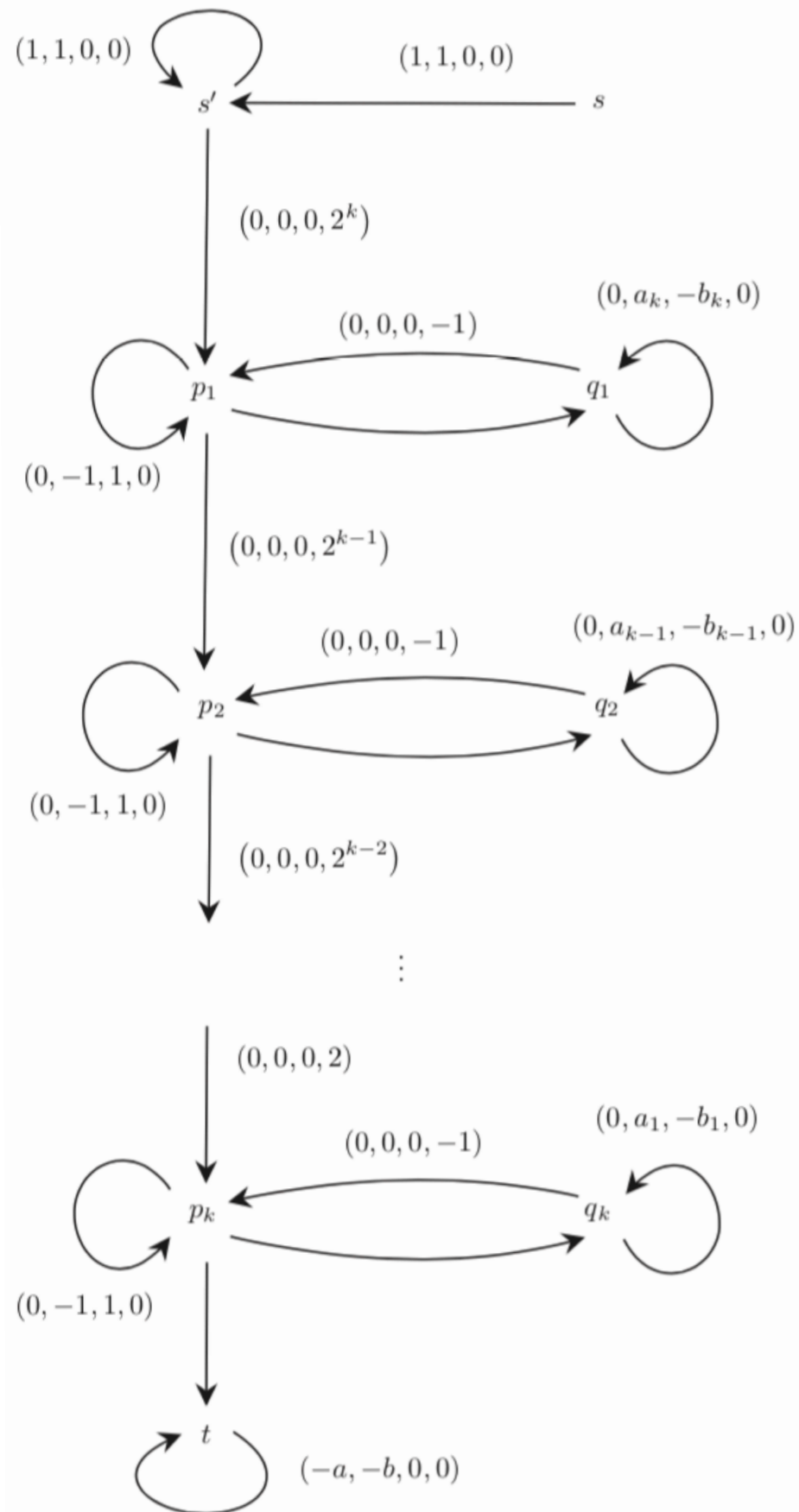


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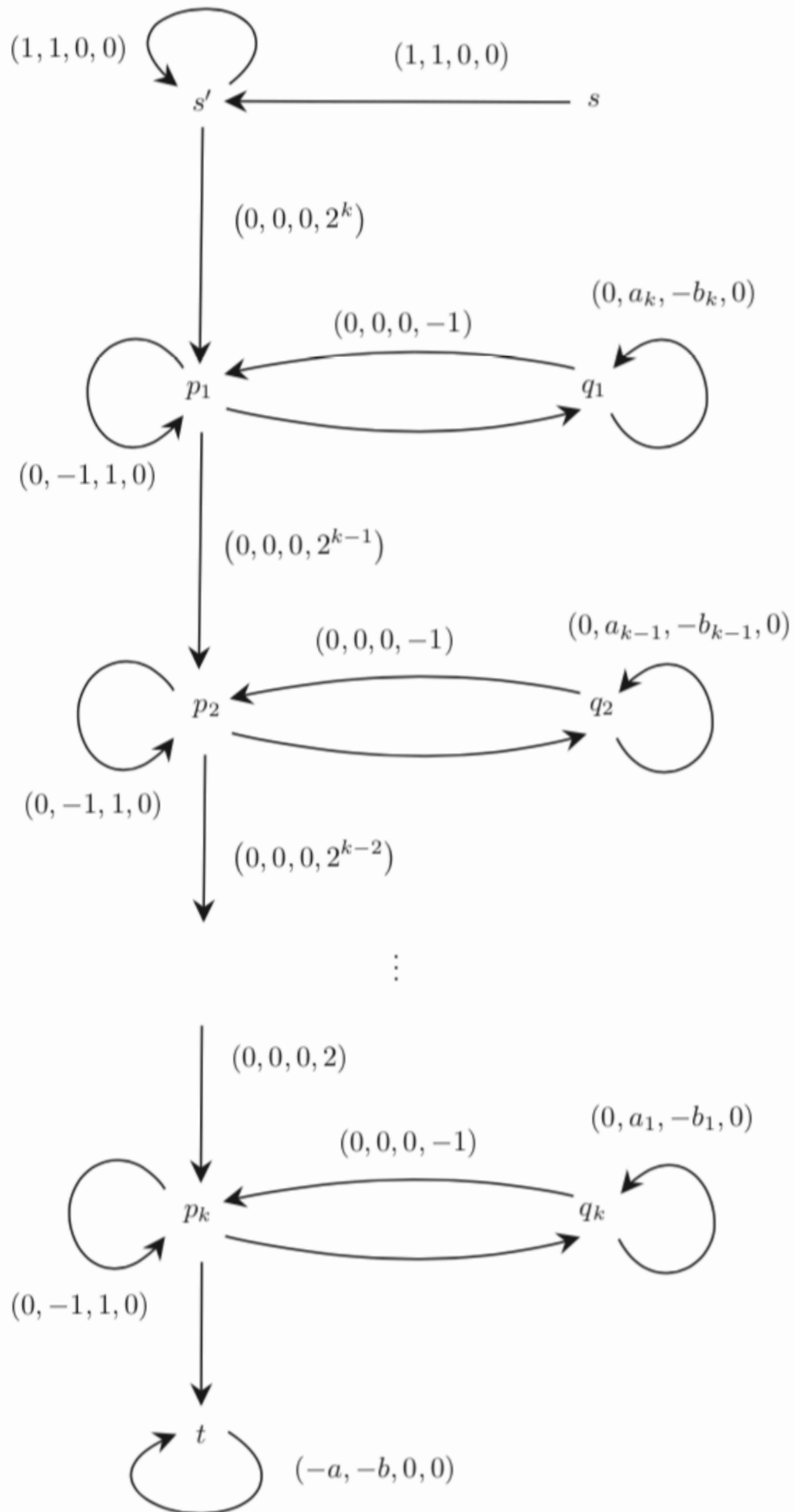
one can multiply by **at most** n times by **at most** a_i / b_i

VASS implementation

VASS implementation

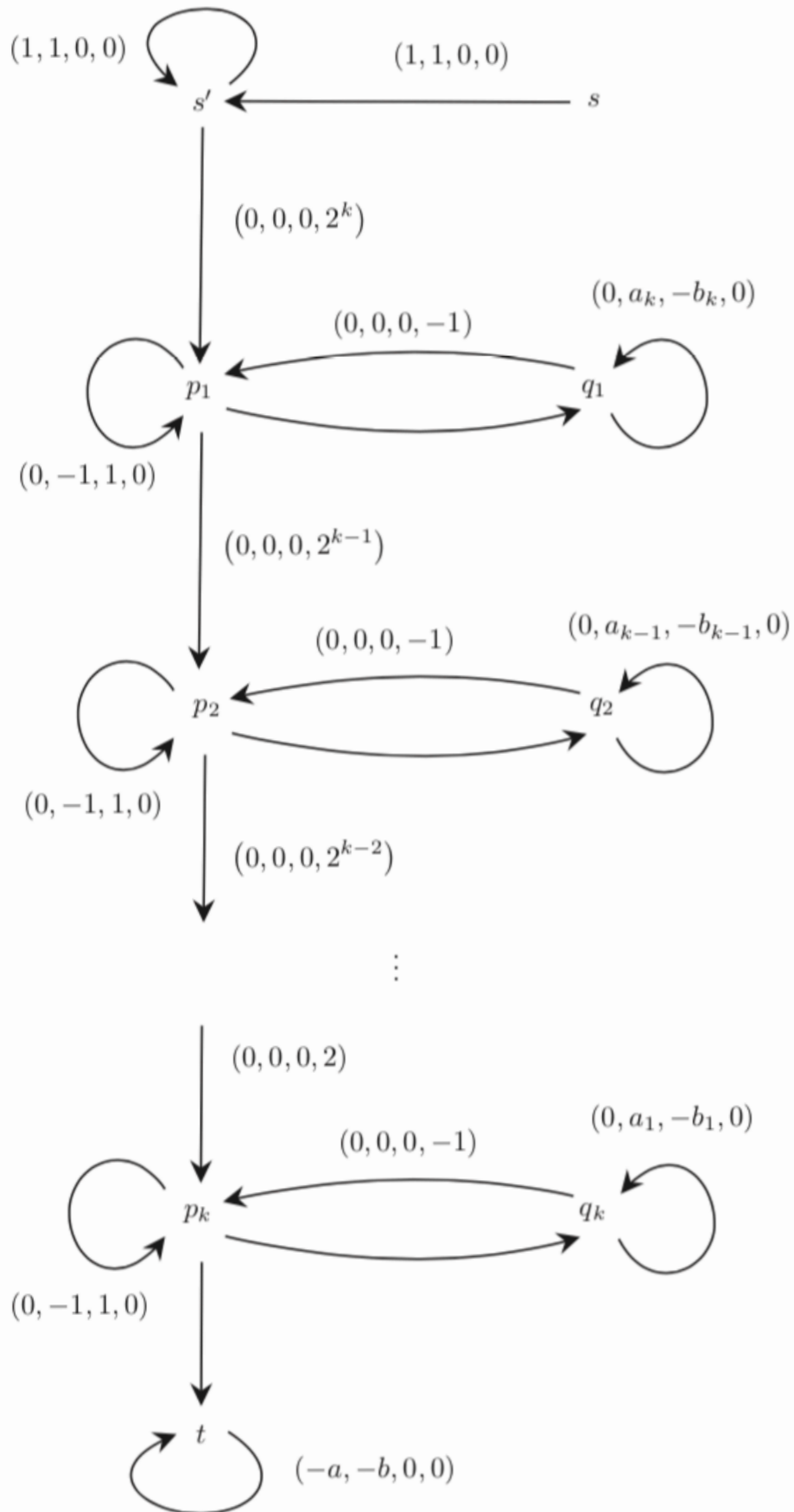


VASS implementation

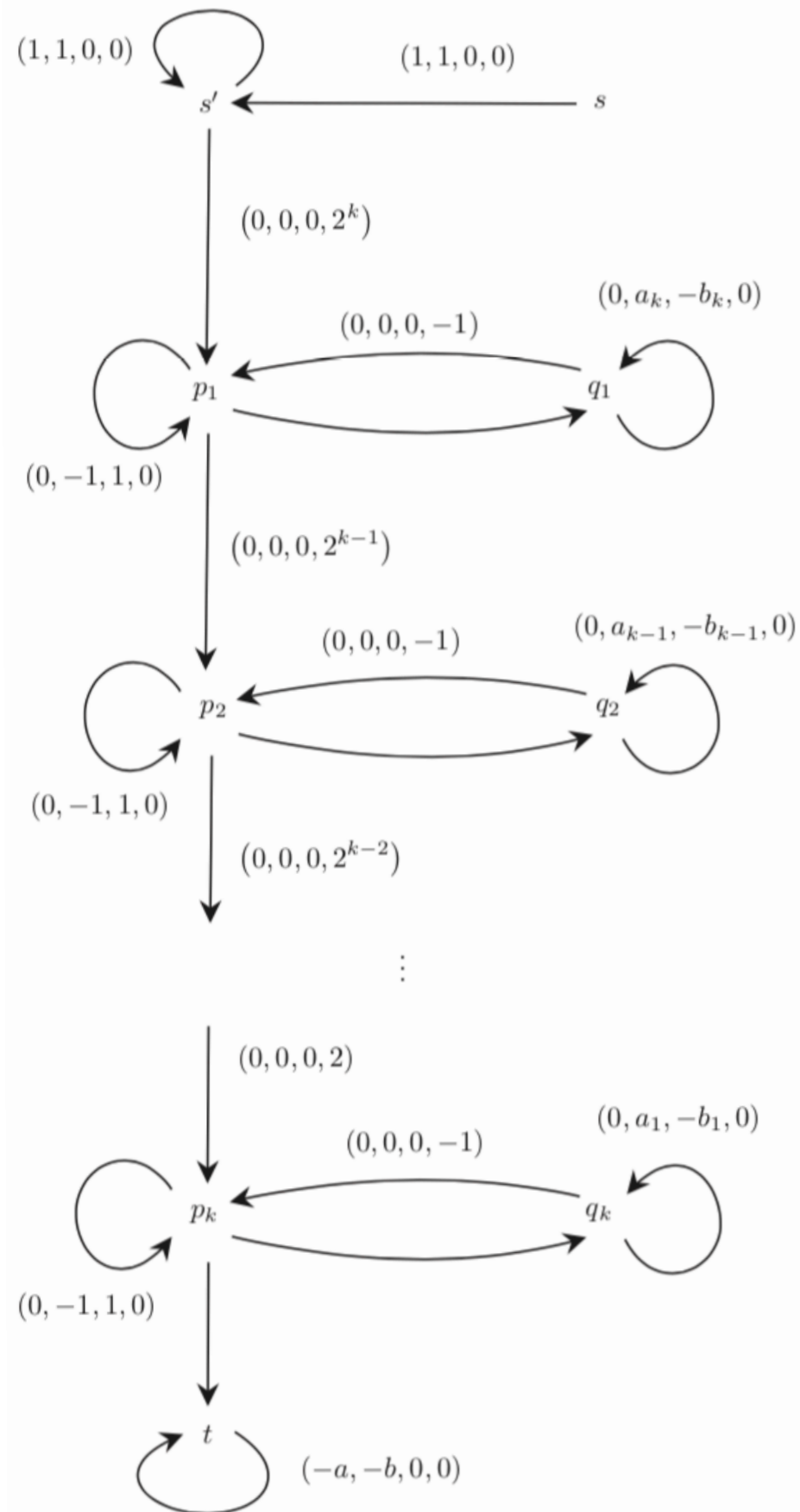


can I reach $t(0,0,0,0)$ from $s(0,0,0,0)$?

VASS implementation

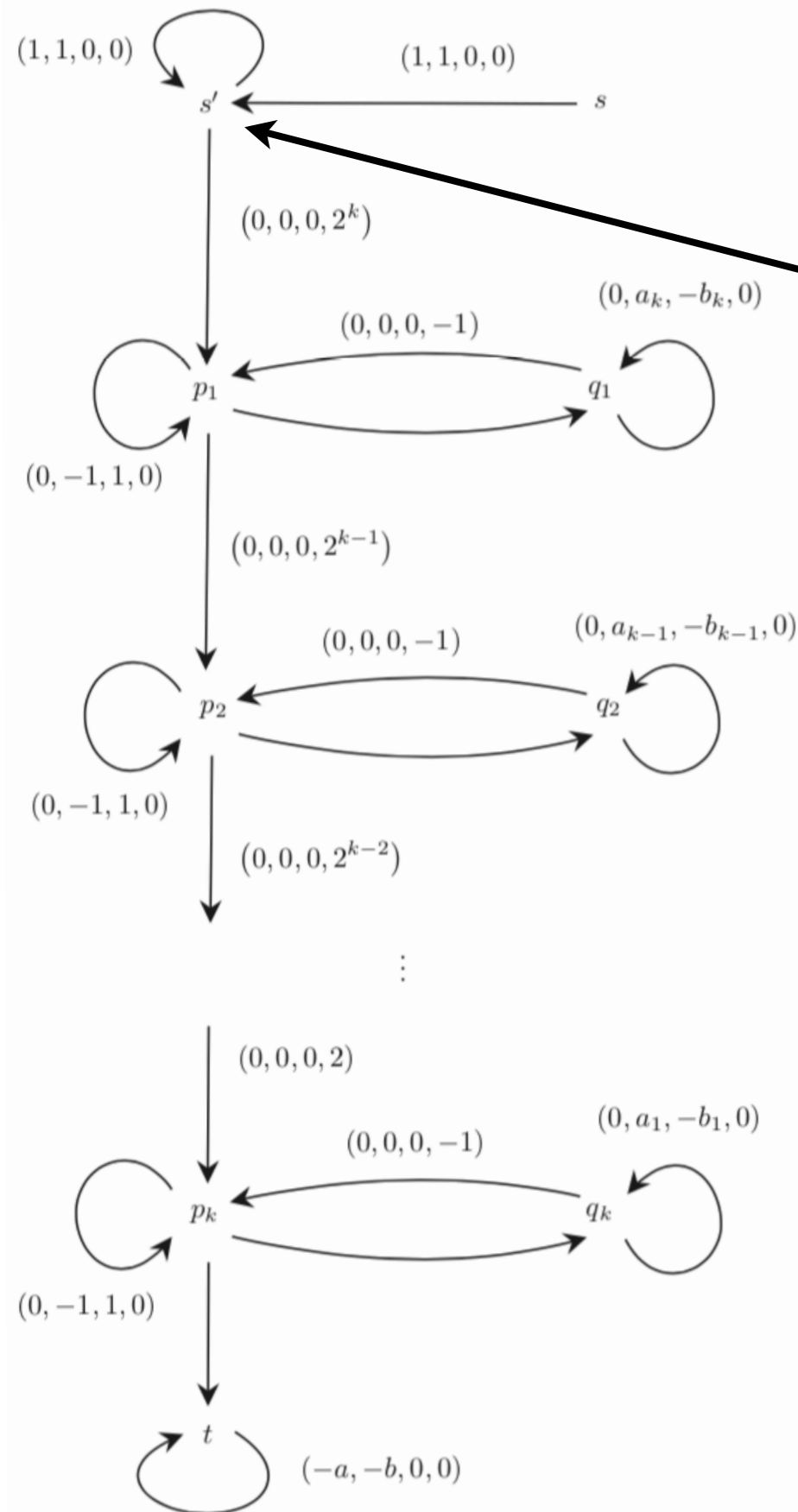


VASS implementation



$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

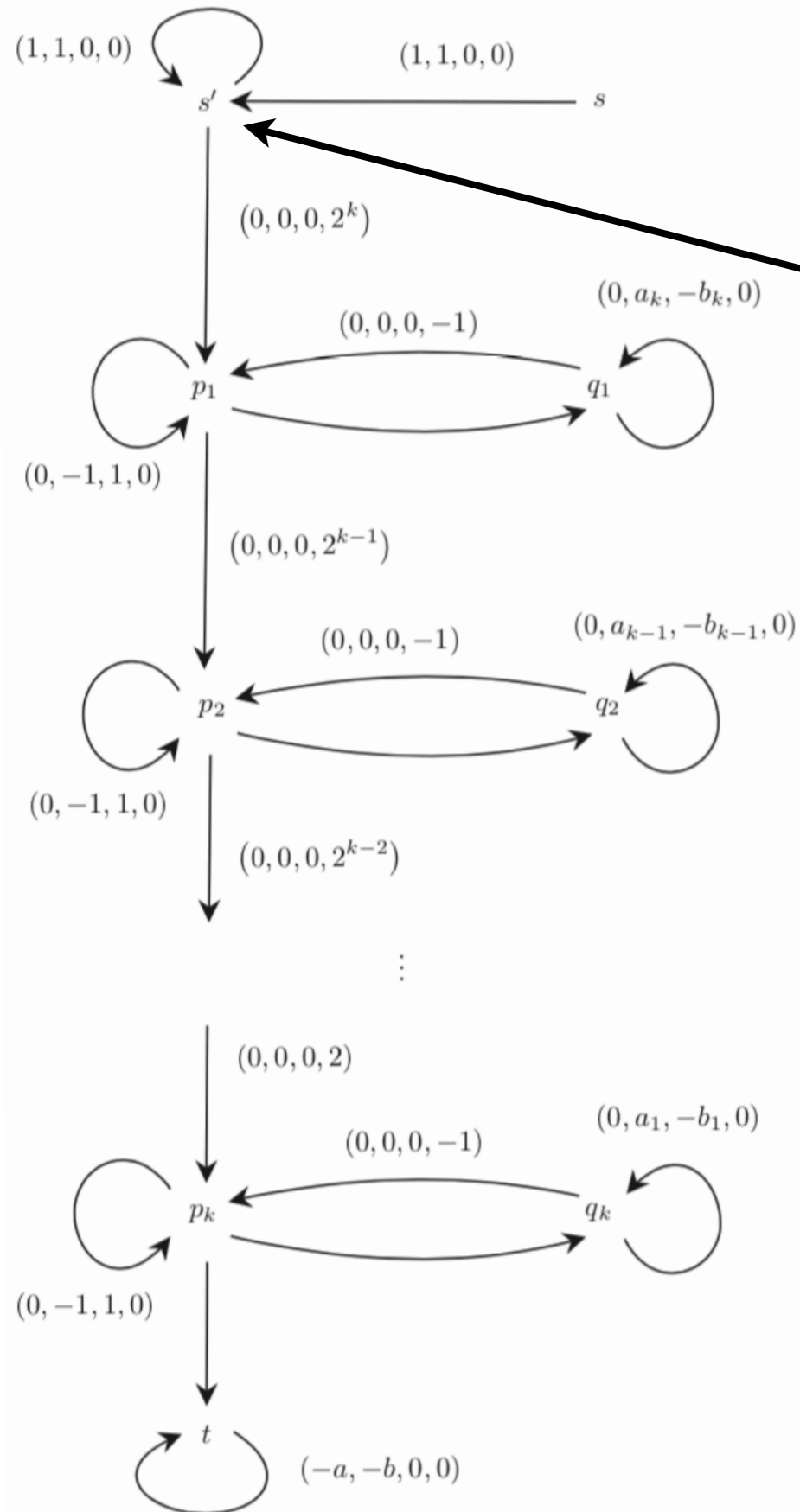
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$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

VASS implementation

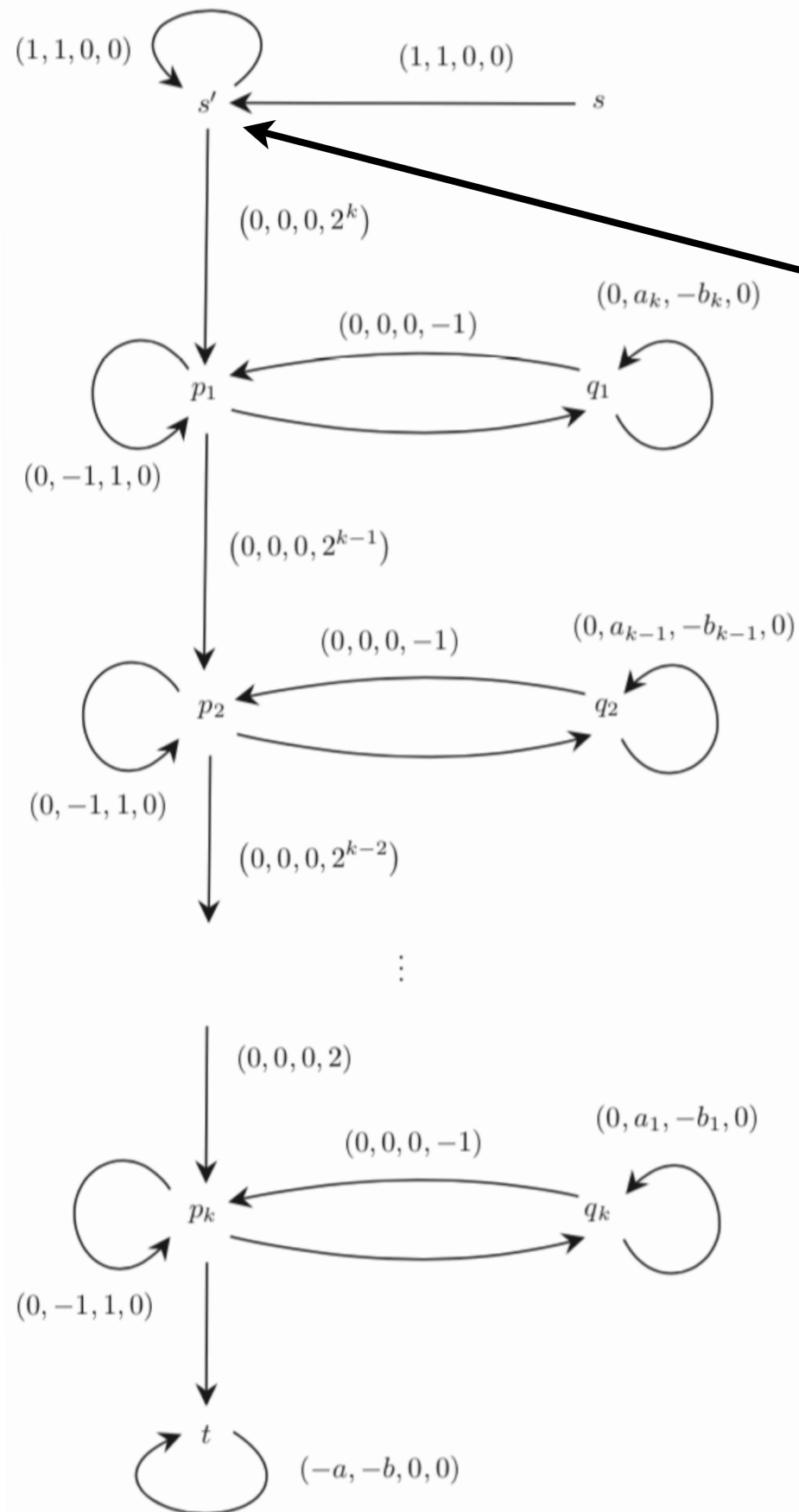


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$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

multiplies \mathbf{K} at most 2^k times by $\leq a_k / b_k$

VASS implementation



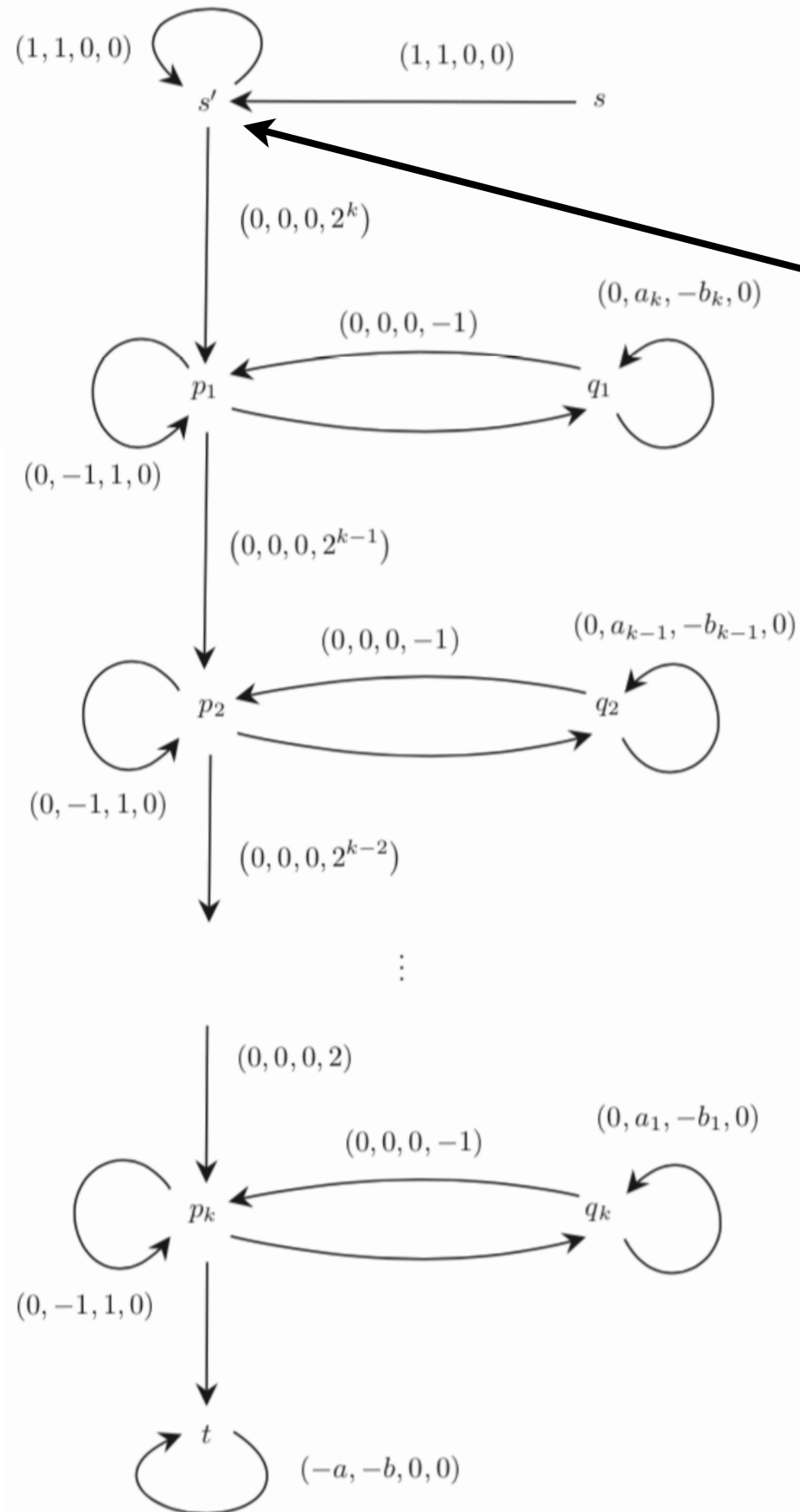
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VASS implementation



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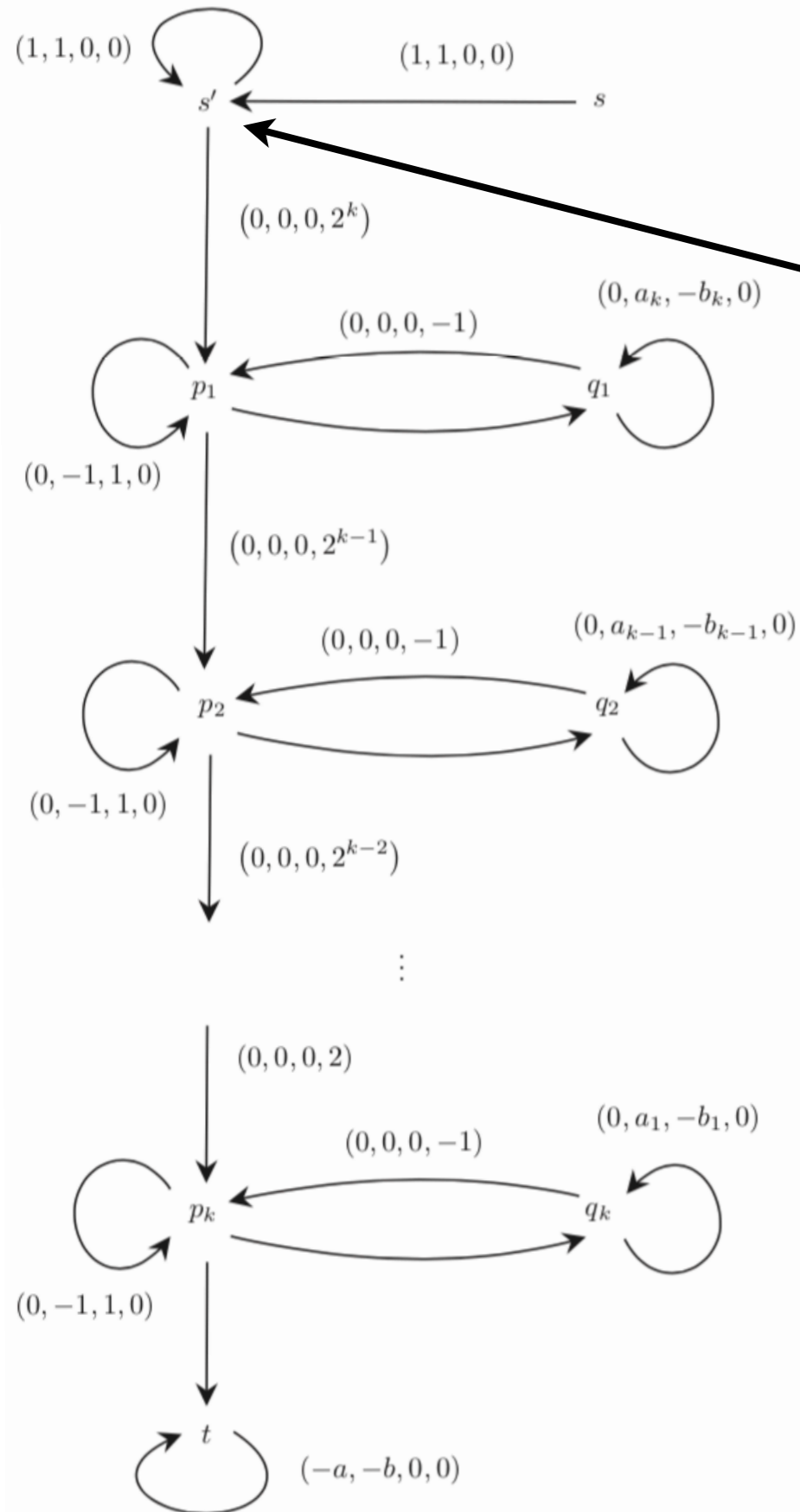
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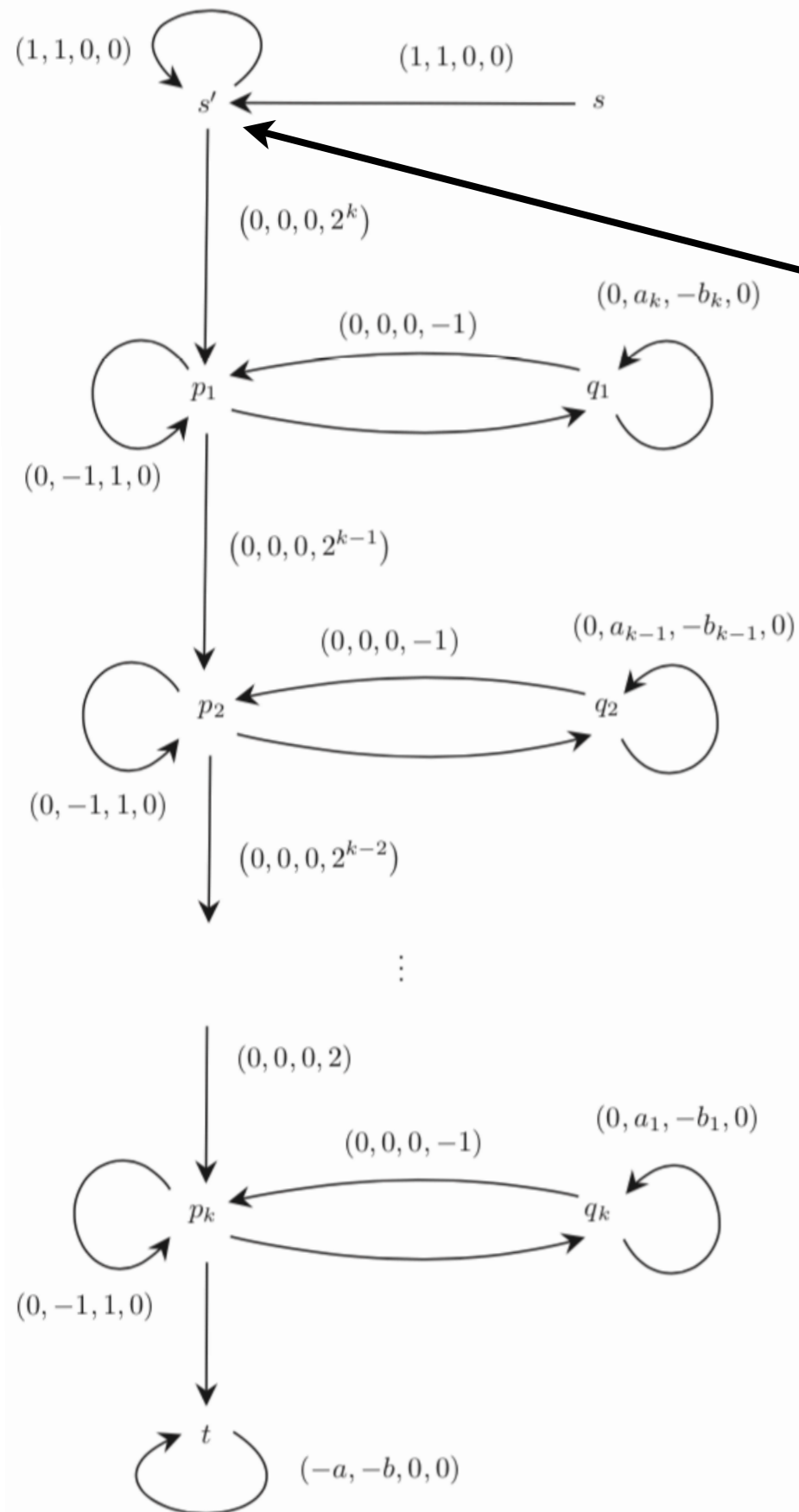
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...

multiplies \mathbf{it} at most 2^1 times by $\leq a_1 / b_1$

$t(\mathbf{K}, \leq \mathbf{K}a/b, \geq 0, \geq 0)$

VASS implementation



$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

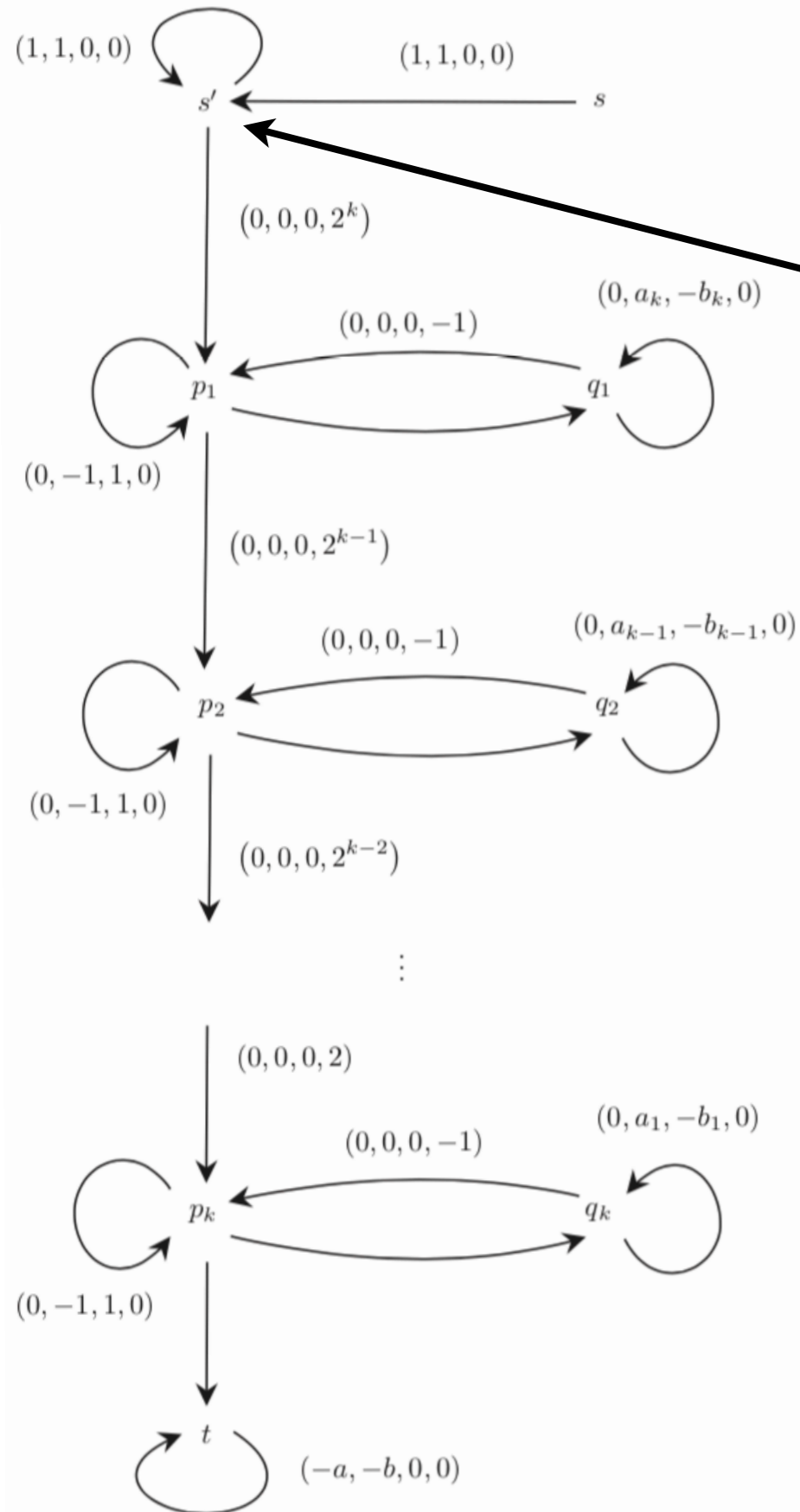
multiplies \mathbf{K} at most 2^k times by $\leq a_k / b_k$

...

multiplies \mathbf{it} at most 2^i times by $\leq a_i / b_i$

subtracts from \mathbf{it} \mathbf{K} a/b $t(\mathbf{K}, \leq \mathbf{K}a/b, \geq 0, \geq 0)$

VASS implementation



$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

multiplies \mathbf{K} at most 2^k times by $\leq a_k / b_k$

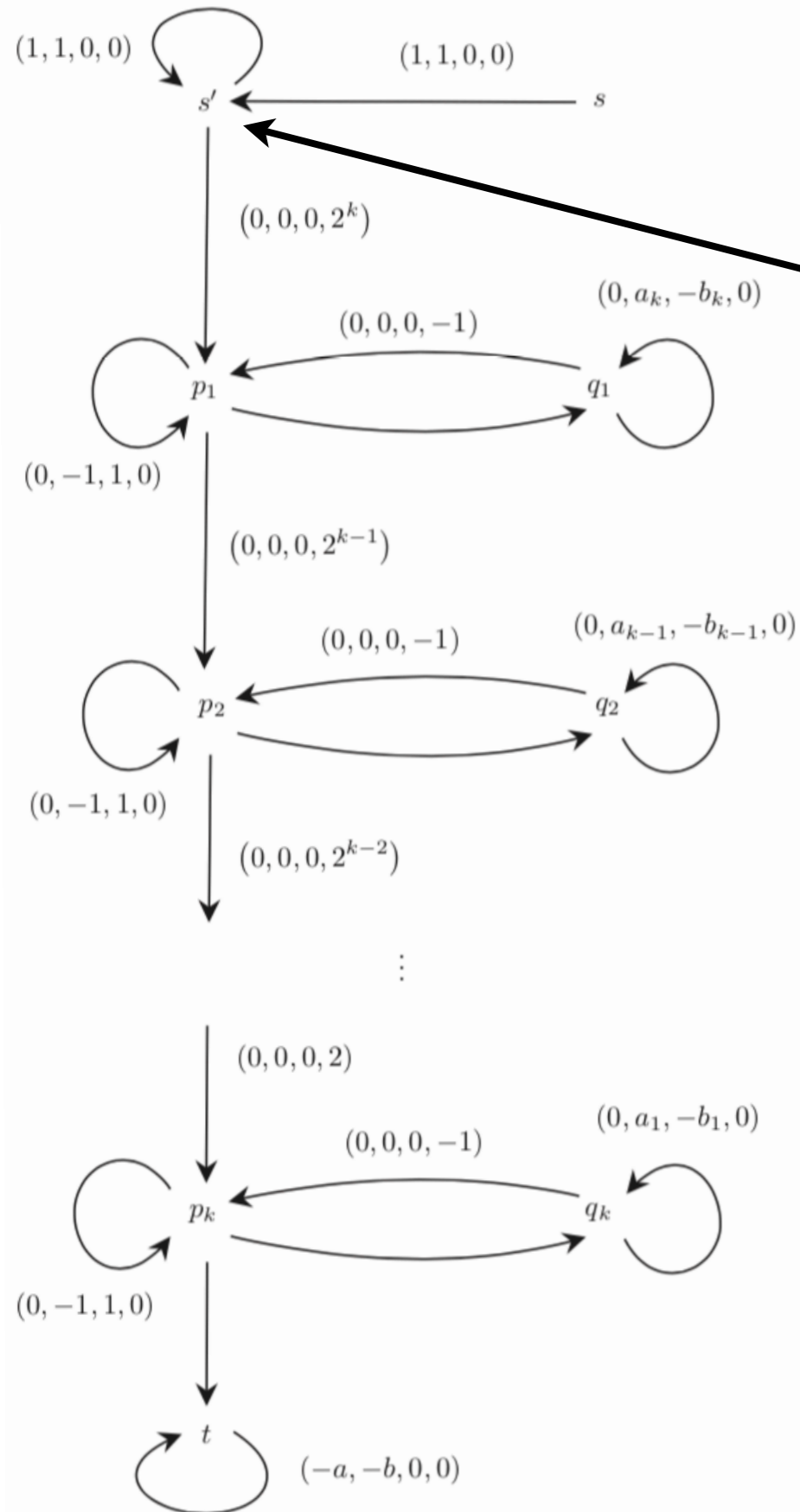
...

multiplies it at most 2^l times by $\leq a_l / b_l$

subtracts from it $\mathbf{K} a/b$ $t(\mathbf{K}, \leq \mathbf{K}a/b, \geq 0, \geq 0)$

\mathbf{K} divisible by $b_k 2^k$

VASS implementation



$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

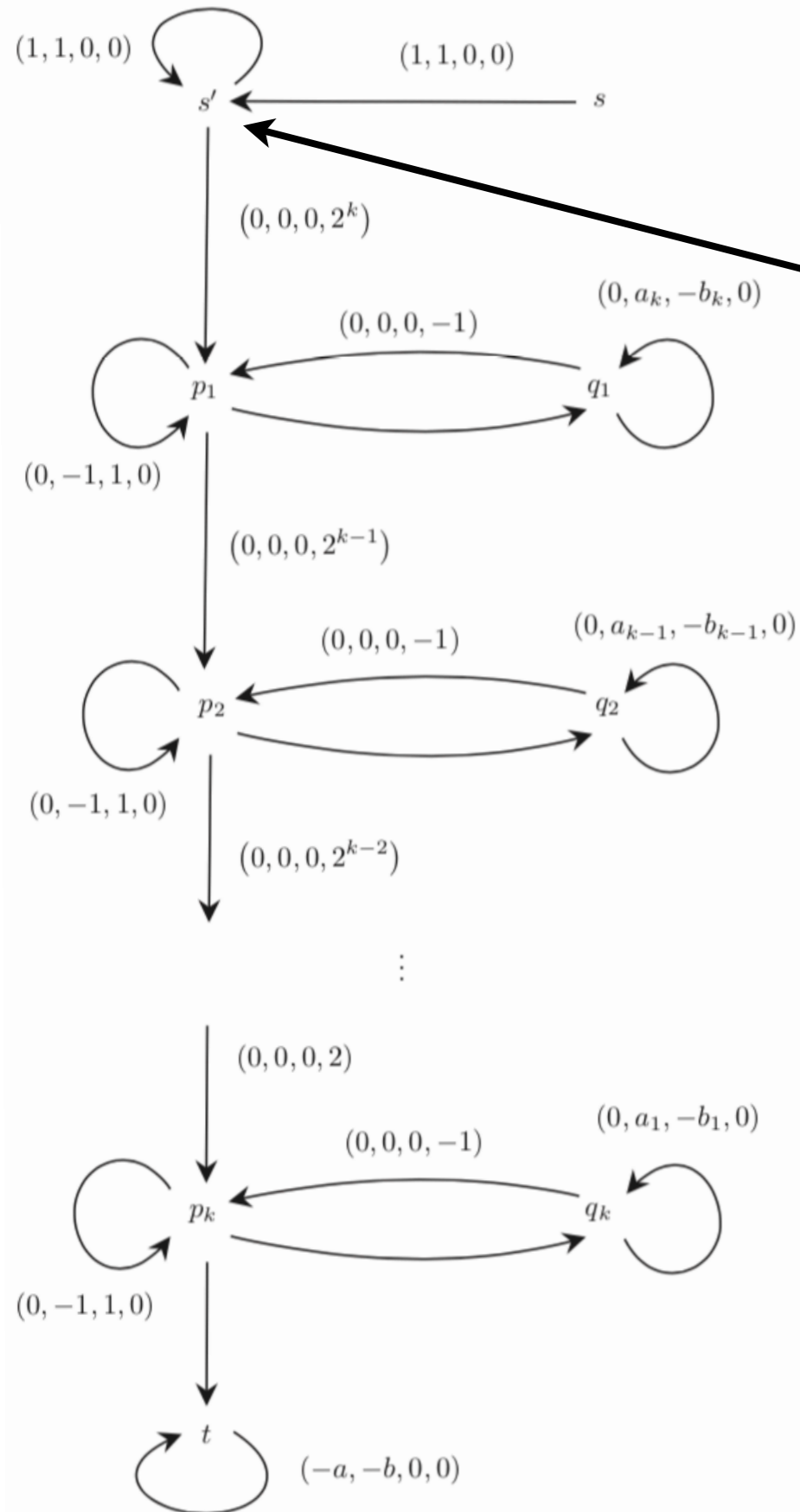
multiplies \mathbf{K} at most 2^k times by $\leq a_k / b_k$

...

multiplies \mathbf{it} at most 2^1 times by $\leq a_1 / b_1$

subtracts from \mathbf{it} \mathbf{K} a/b $t(\mathbf{K}, \leq \mathbf{K}a/b, \geq 0, \geq 0)$

VASS implementation



$$(a_1 / b_1)^{2^1} \cdot (a_2 / b_2)^{2^2} \cdot \dots \cdot (a_k / b_k)^{2^k} = a / b$$

$s'(\mathbf{K}, \mathbf{K}, 0, 0)$

multiplies \mathbf{K} at most 2^k times by $\leq a_k / b_k$

...

multiplies \mathbf{it} at most 2^l times by $\leq a_l / b_l$

subtracts from $\mathbf{it} \mathbf{K} a/b$ $t(\mathbf{K}, \leq \mathbf{K}a/b, \geq 0, \geq 0)$

\mathbf{K} doubly exponential

Open problem

Open problem

Reachability for finite 3-VASSes

Open problem

Reachability for finite 3-VASSes

PSpace?

Open problem

Reachability for finite 3-VASSes

PSpace?

Tower?

Open problem

Reachability for finite 3-VASSes

PSpace?

Tower?

Conjecture: PSpace

Thank you!