

Multi-Stack Automata

Wojciech Czerwiński

Multi-Stack Automata

(Partially-Commutative Context-Free Graphs)

Wojciech Czerwiński

Plan

Plan

- what are Multi-Stack Automata

Plan

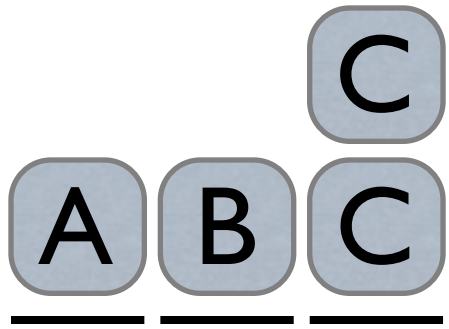
- what are Multi-Stack Automata
- why I investigate them

Plan

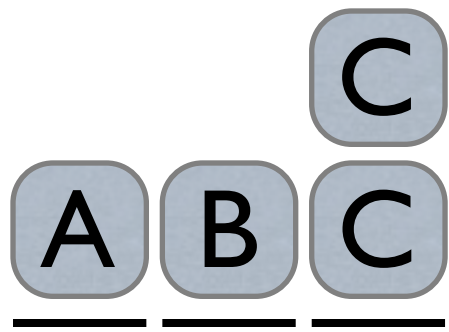
- what are Multi-Stack Automata
- why I investigate them
- what I have solved

Multi-Stack Automaton

Multi-Stack Automaton

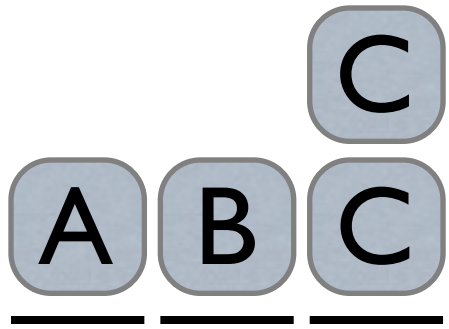


Multi-Stack Automaton

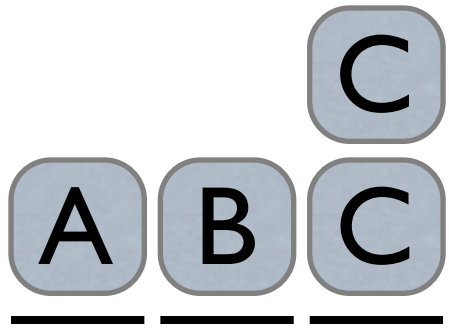


Configuration
(3 stacks)

Multi-Stack Automaton

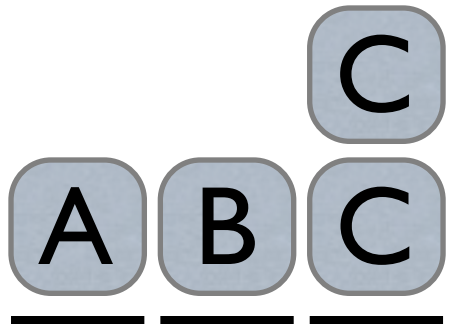


Multi-Stack Automaton



$C \rightarrow (A, BD, C)$

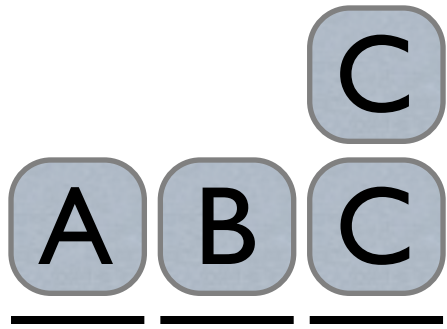
Multi-Stack Automaton



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Multi-Stack Automaton

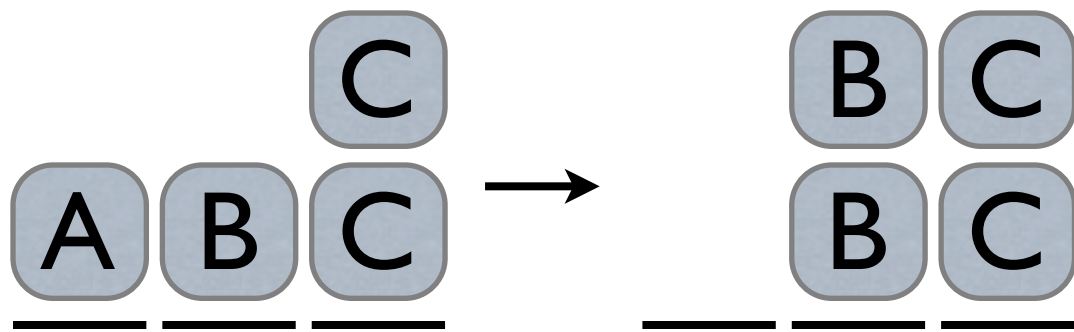


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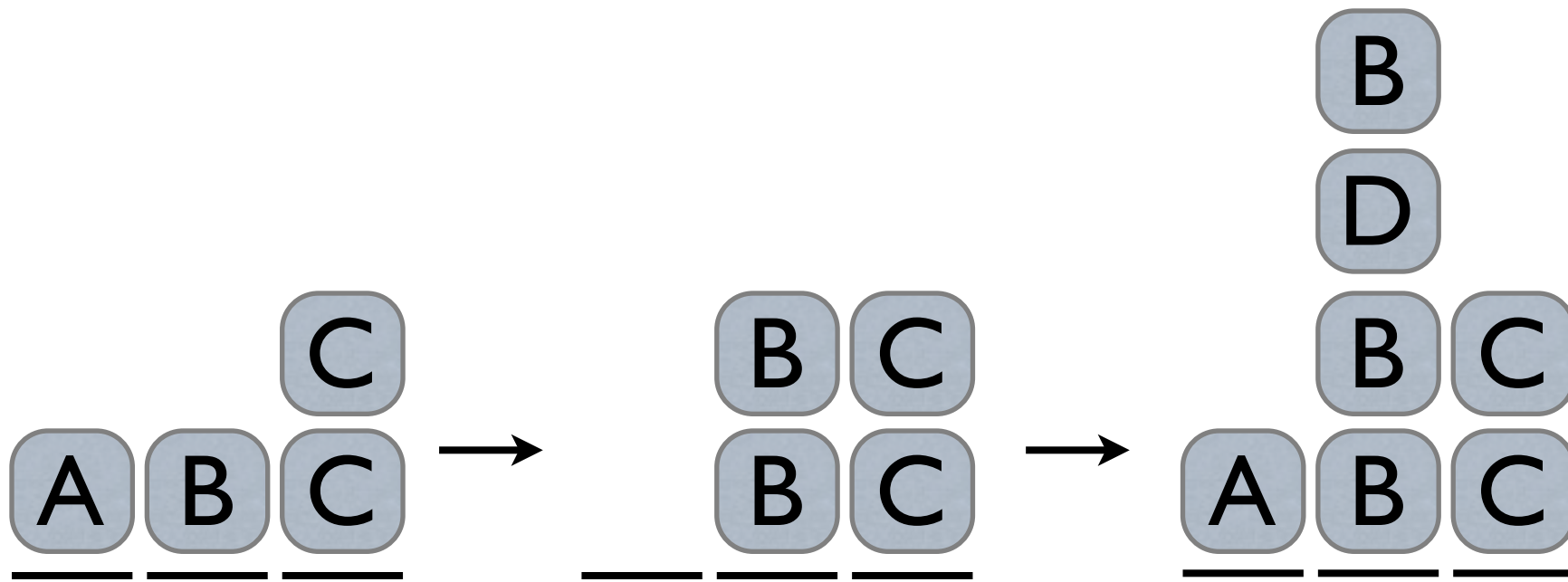


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Multi-Stack Automaton

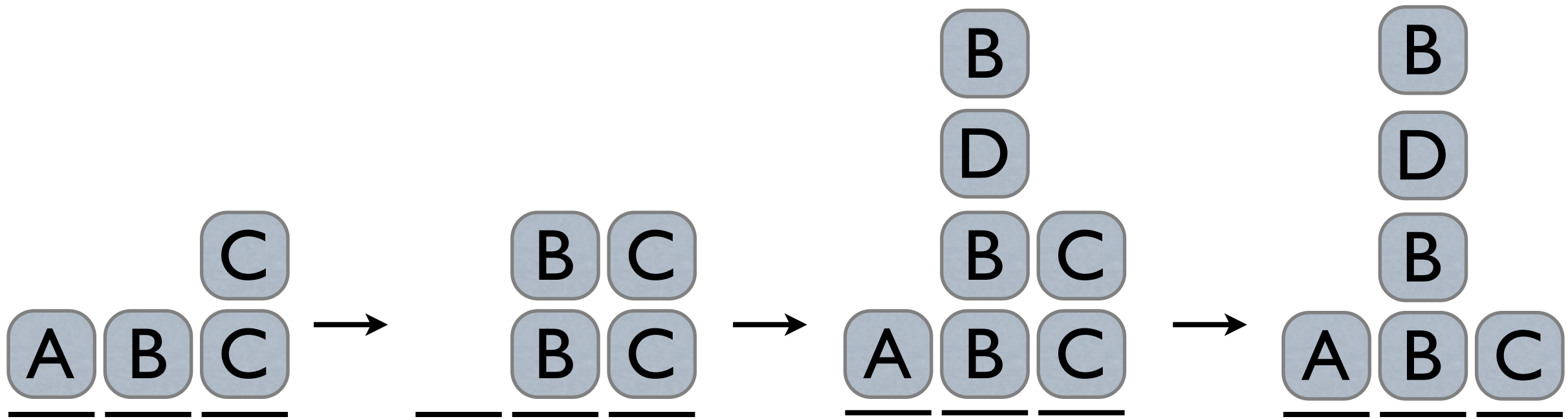


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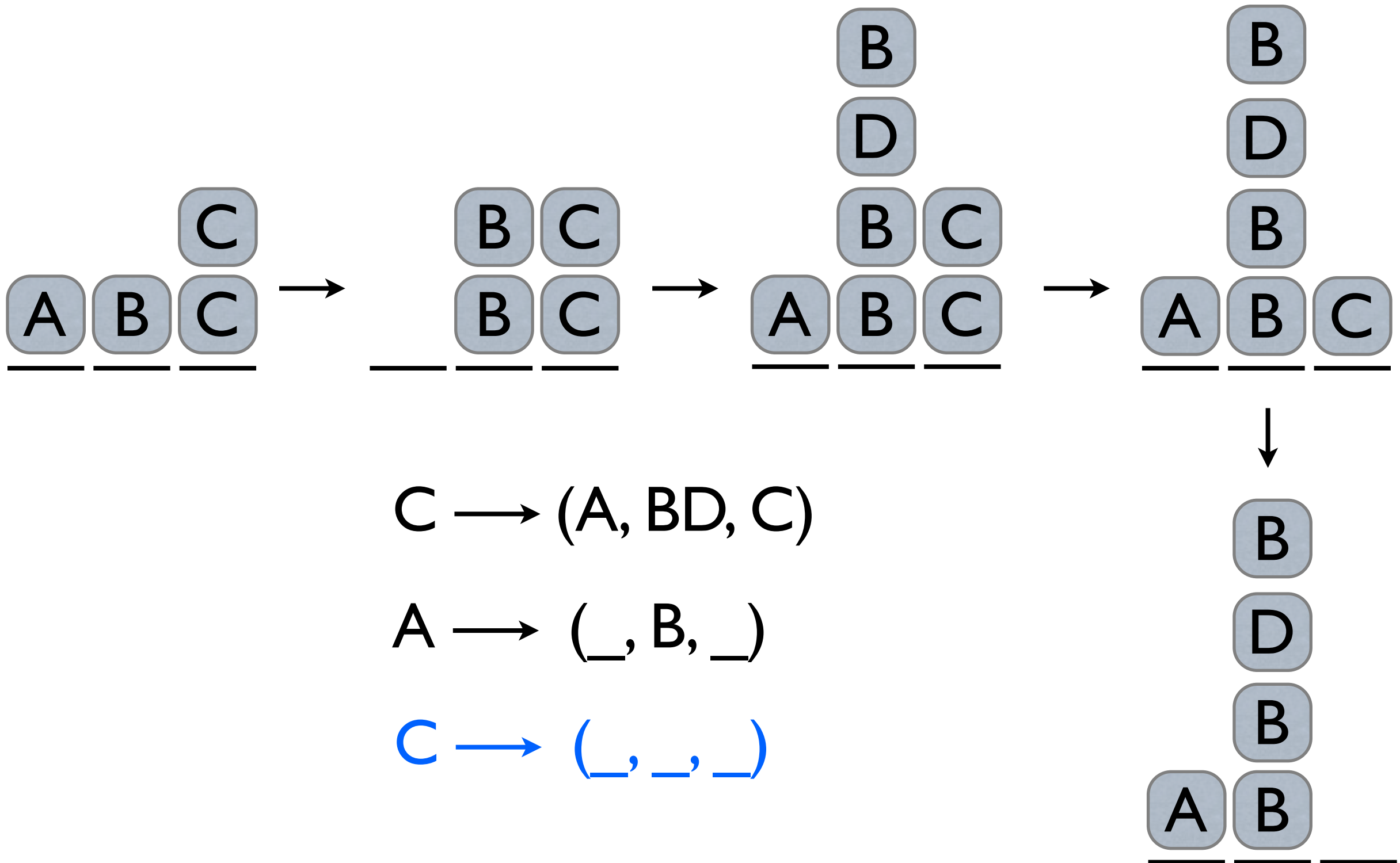


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Multi-Stack Automaton



Motivation

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- errors in programs

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- automatic detection of errors

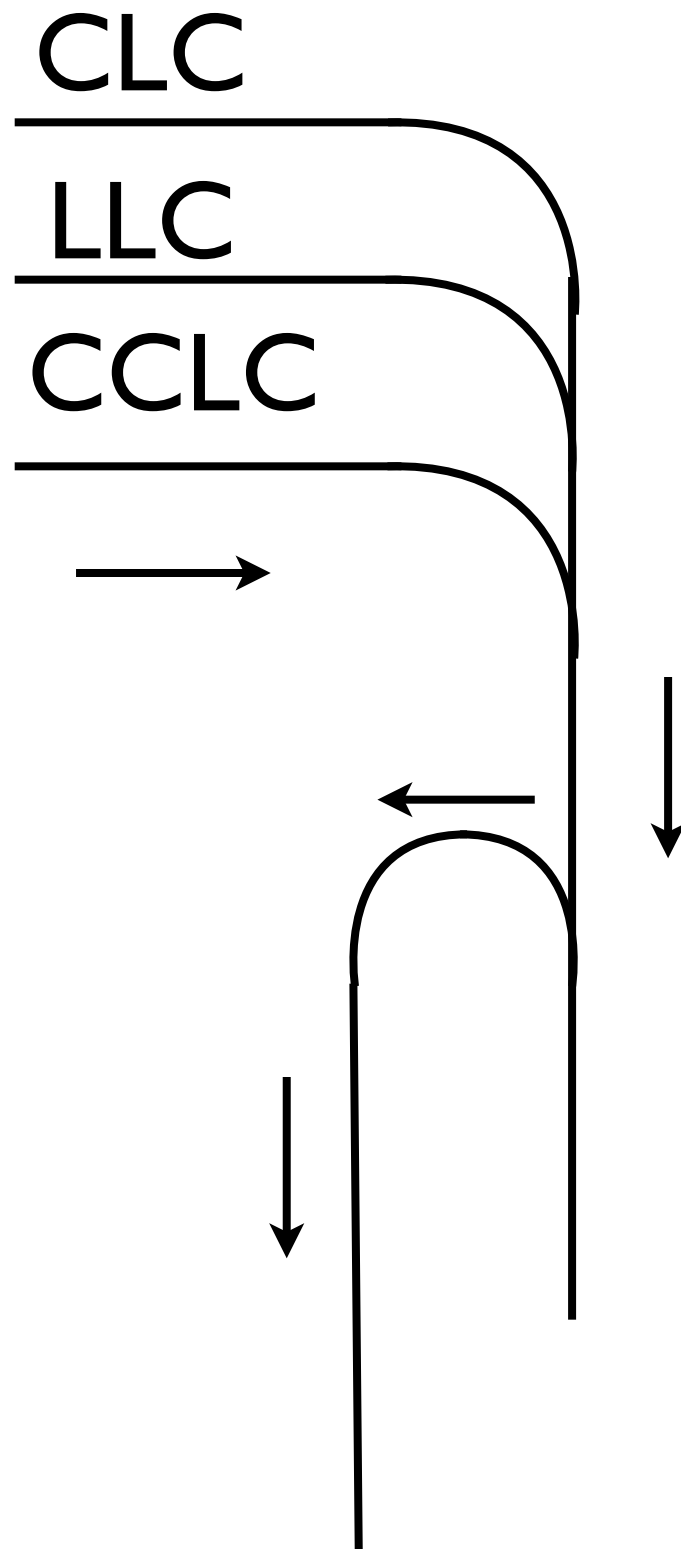
Motivation

- errors in programs
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- model the program

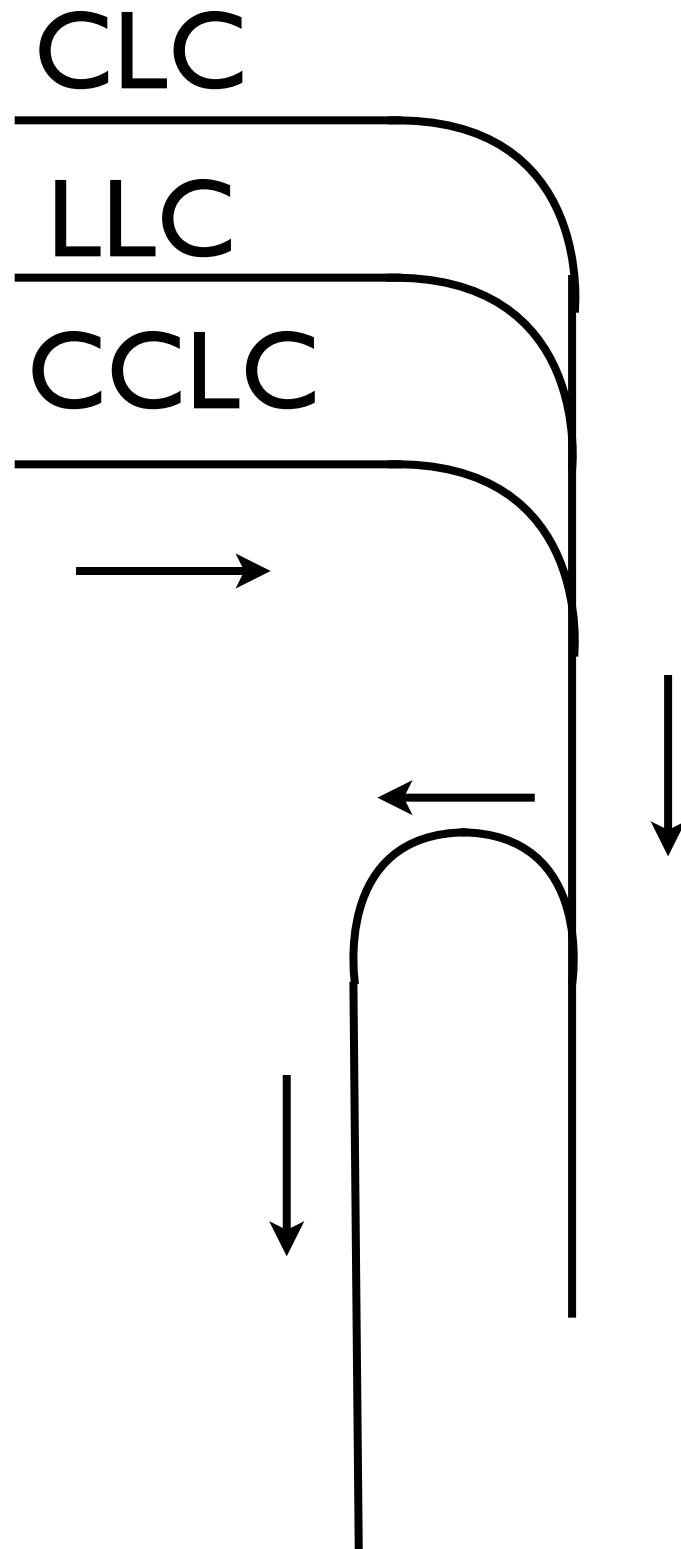
Motivation

- errors in programs
- automatic detection of errors
- model the program
- Multi-Stack Automata as a model

Real life example

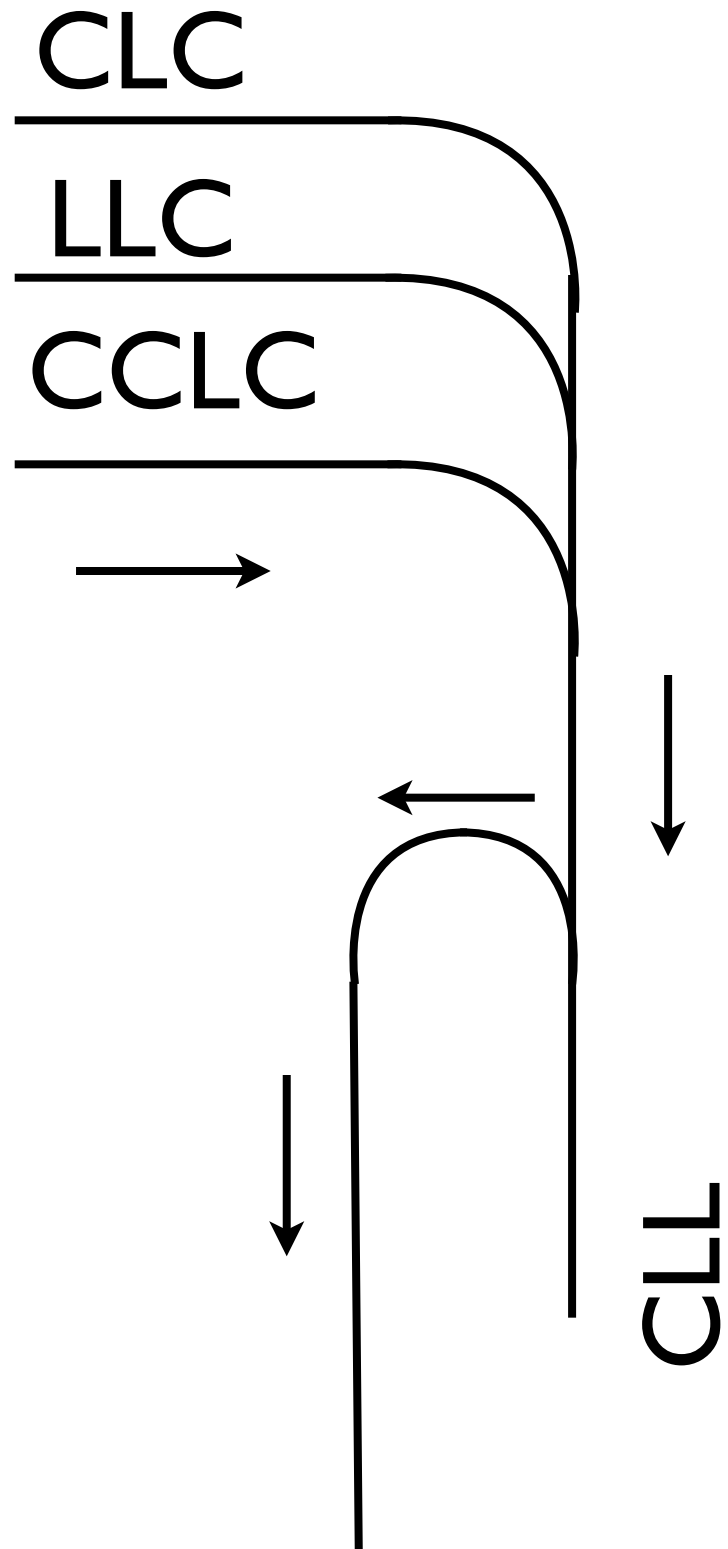


Real life example



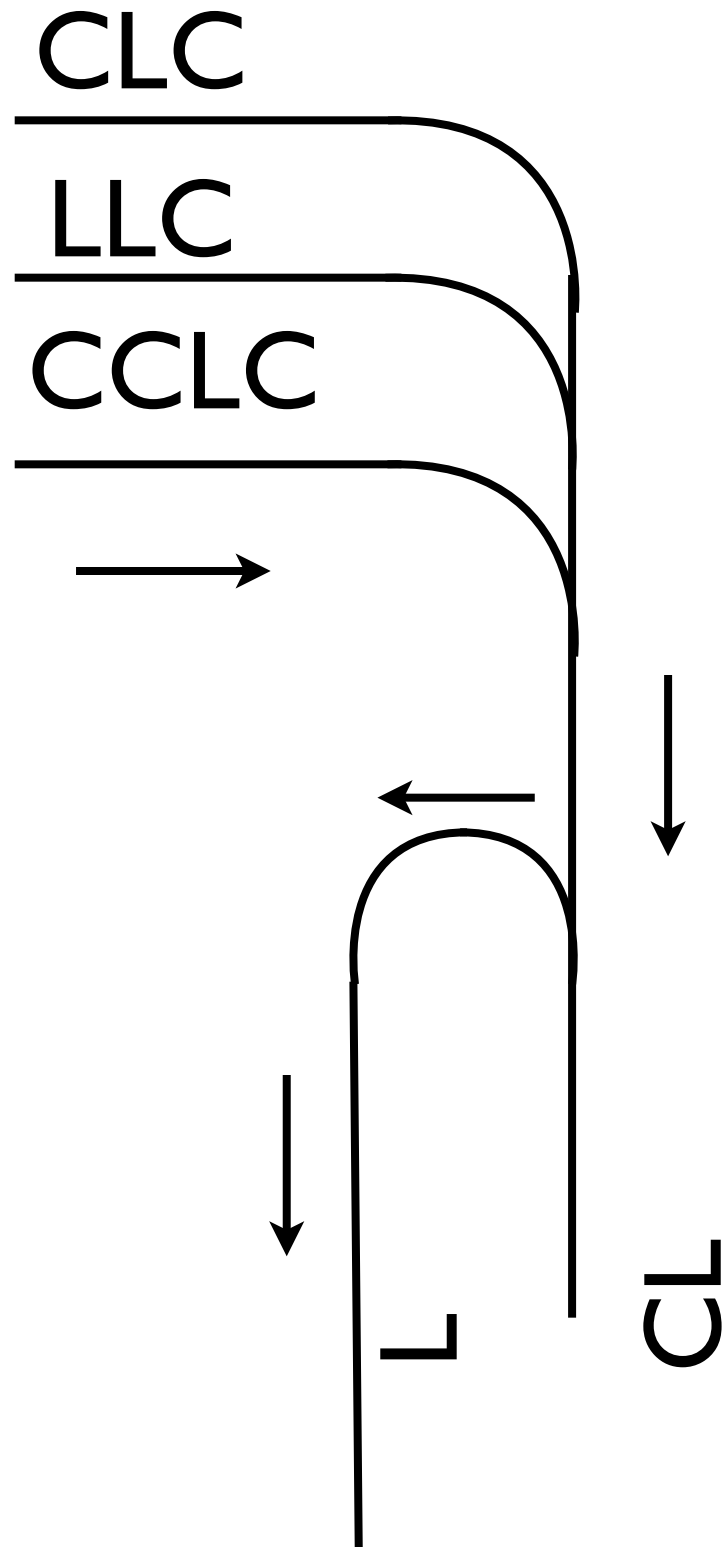
Can we produce trains with
2 cars and 1 locomotive?

Real life example



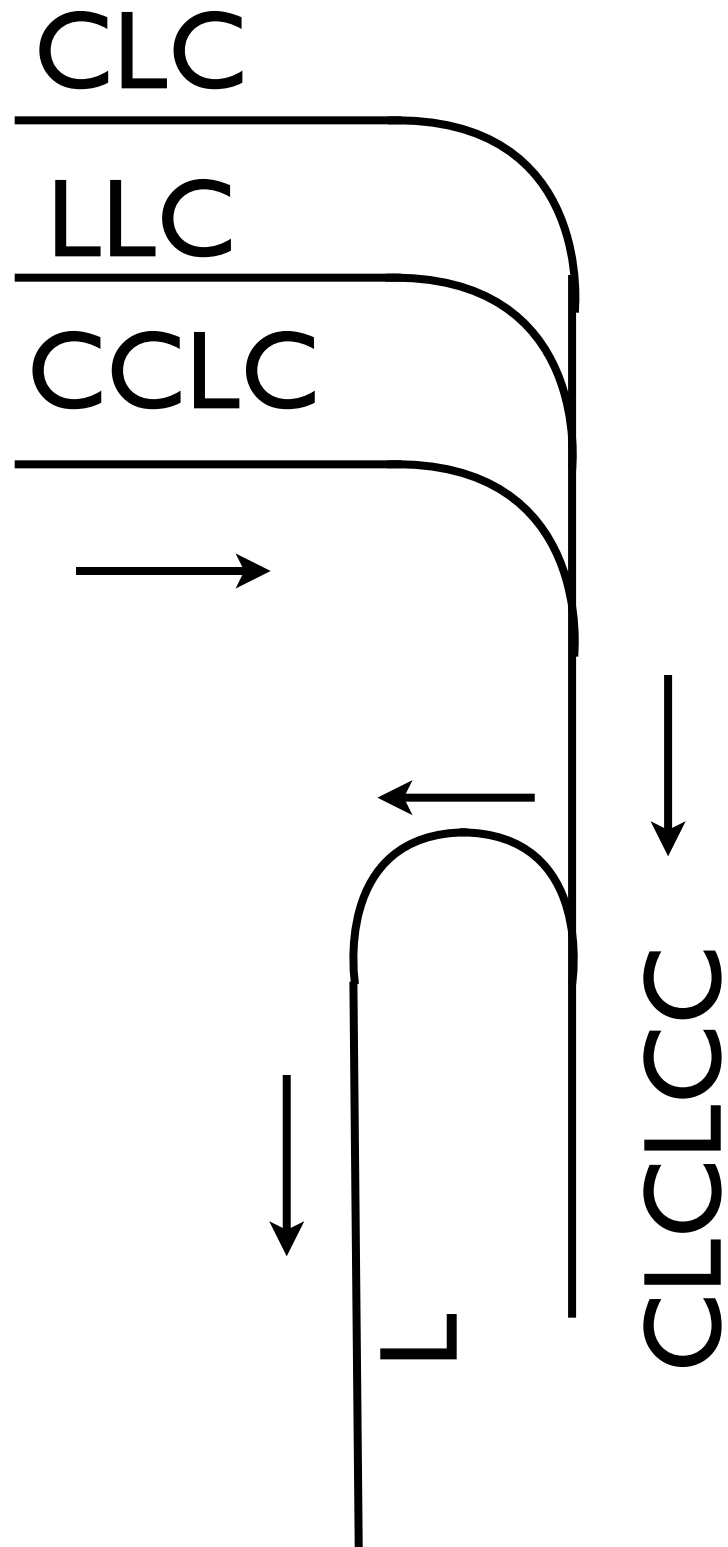
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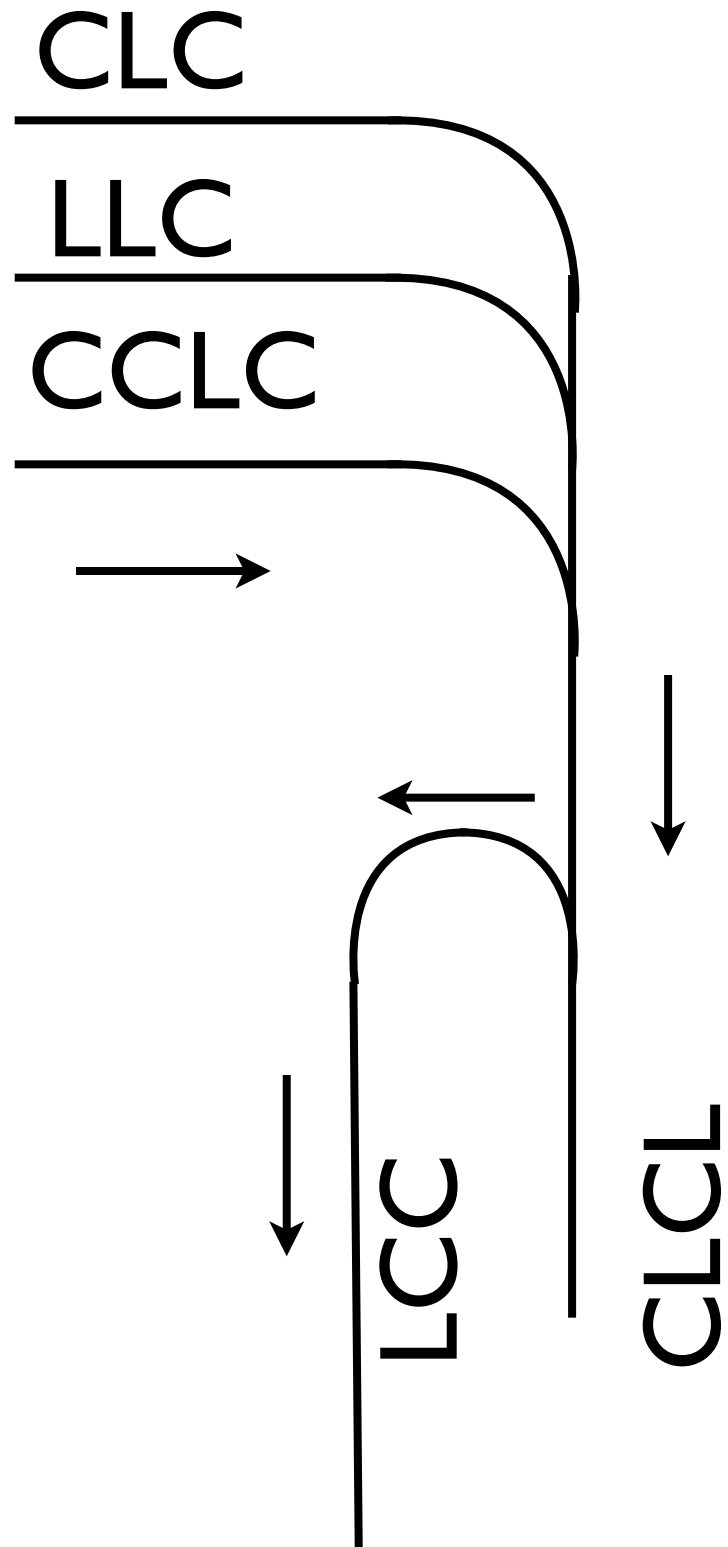
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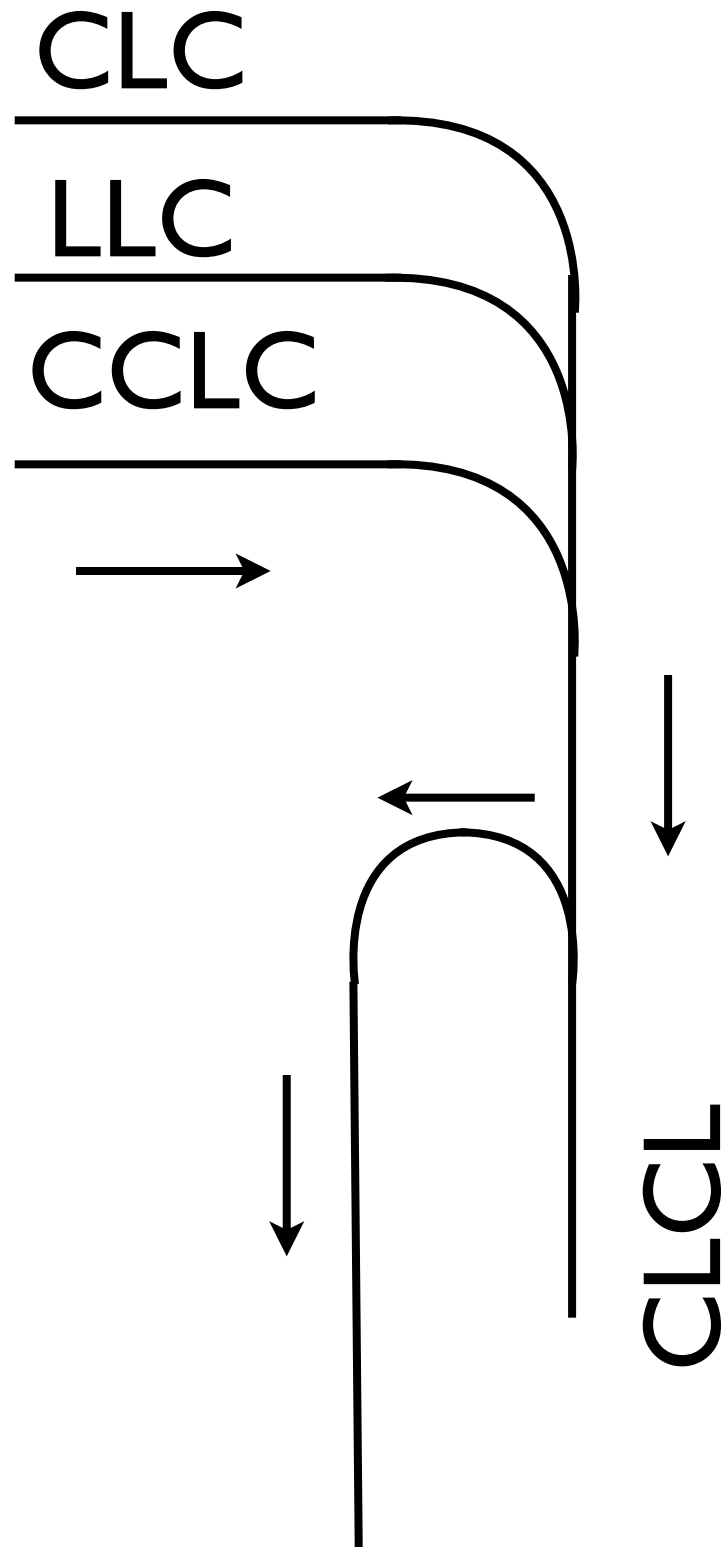
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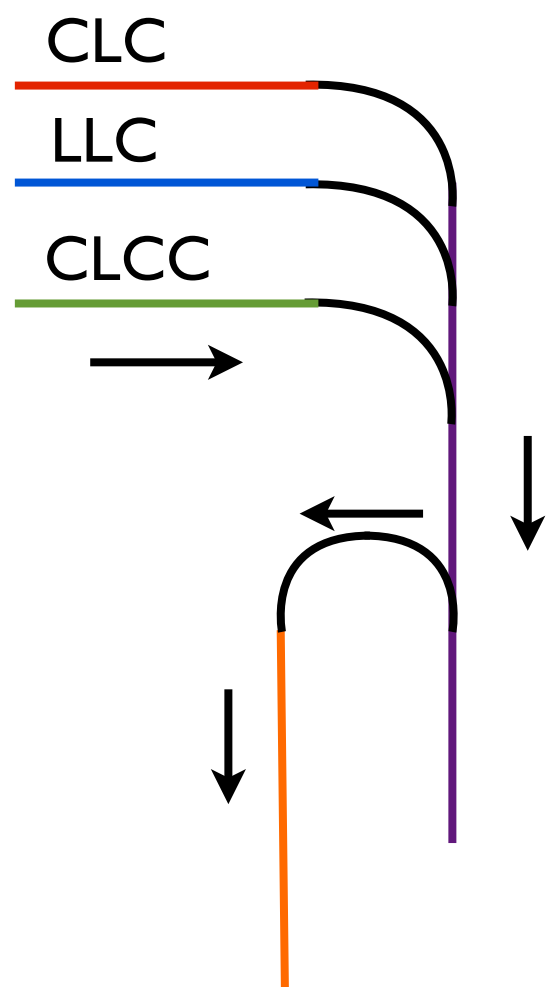
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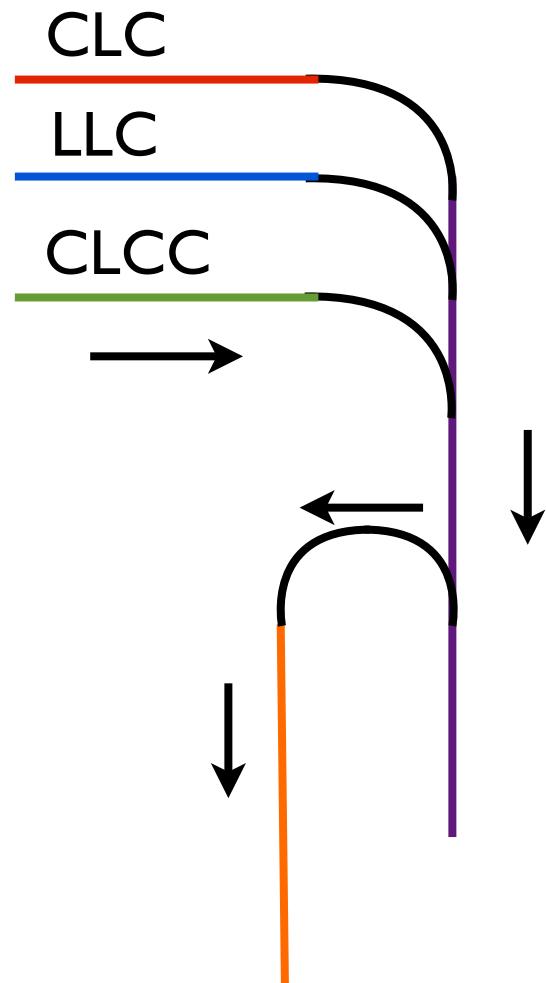


Can we produce trains with
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Modelling

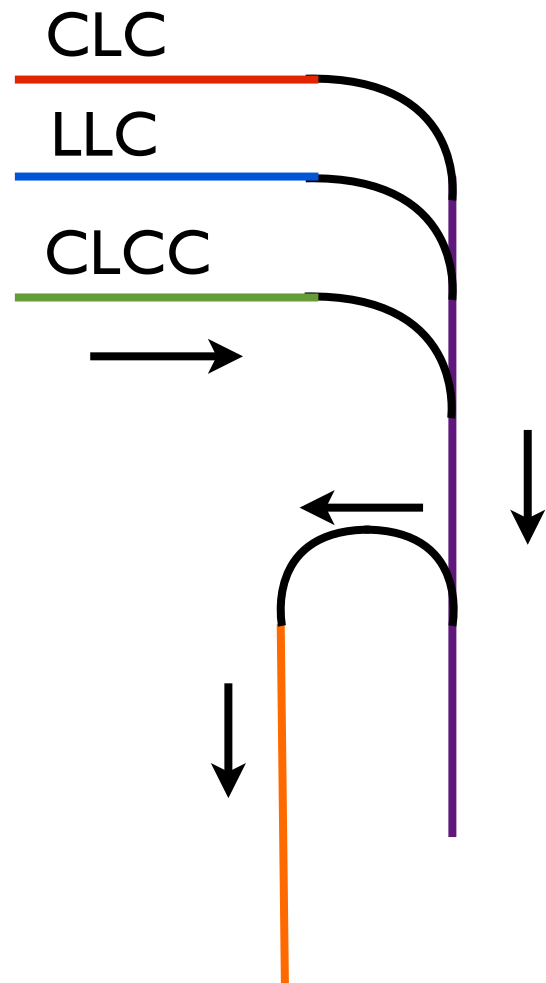


Modelling



$$T \longrightarrow (T, _, _, \text{CLC}, _)$$

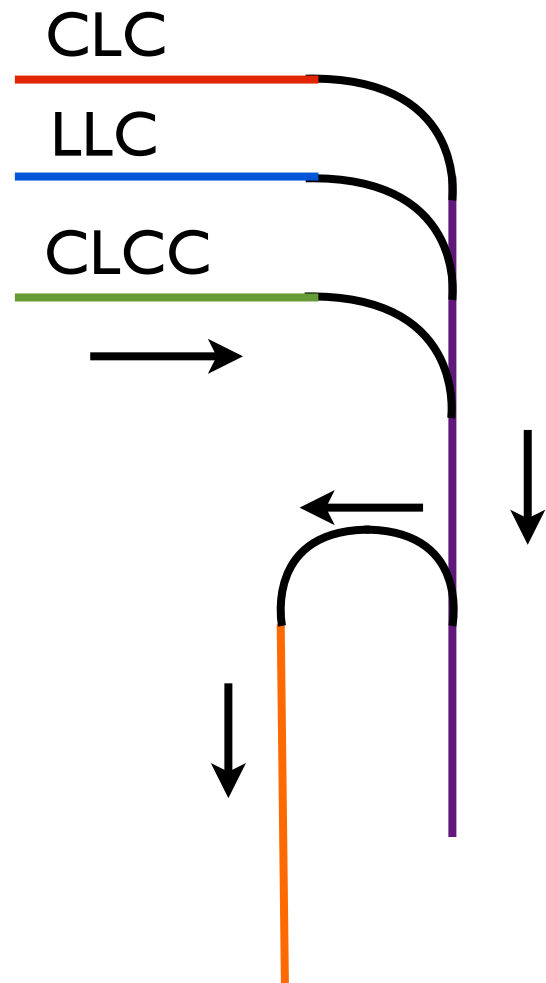
Modelling



$$T \longrightarrow (T, _, _, \text{CLC}, _)$$

$$T \longrightarrow (_, T, _, \text{LLC}, _)$$

Modelling

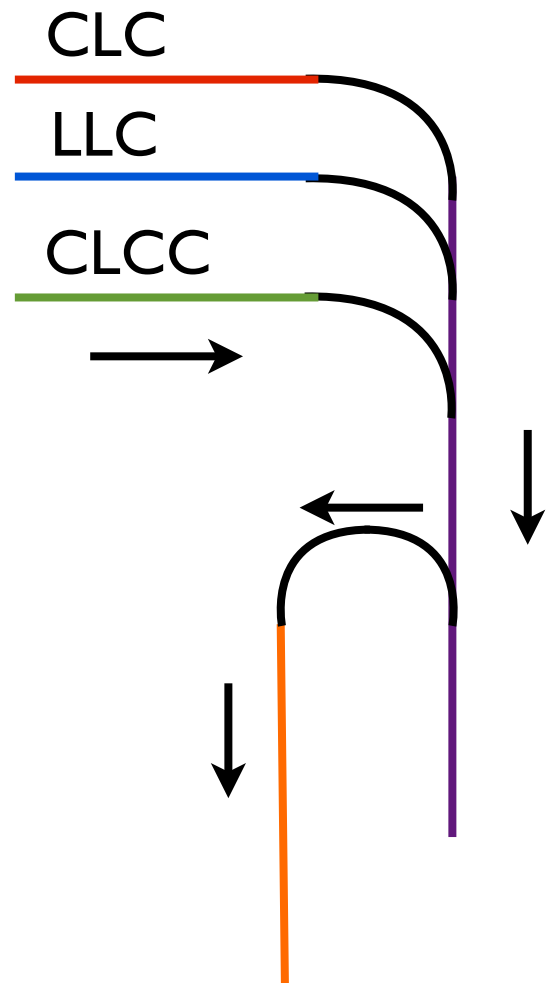


T \longrightarrow (T, _, _, CLC, _)

T \longrightarrow (_, T, _, LLC, _)

T \longrightarrow (_, _, T, CLCC, _)

Modelling



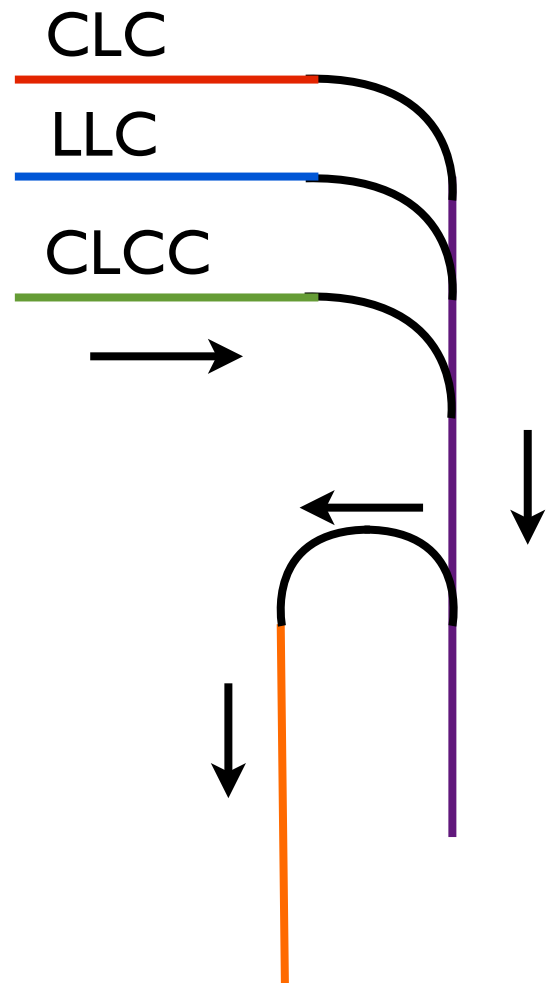
T \longrightarrow (T, _, _, CLC, _)

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C \longrightarrow (_, _, _, _, C)

Modelling



T \longrightarrow (T, _, _, CLC, _)

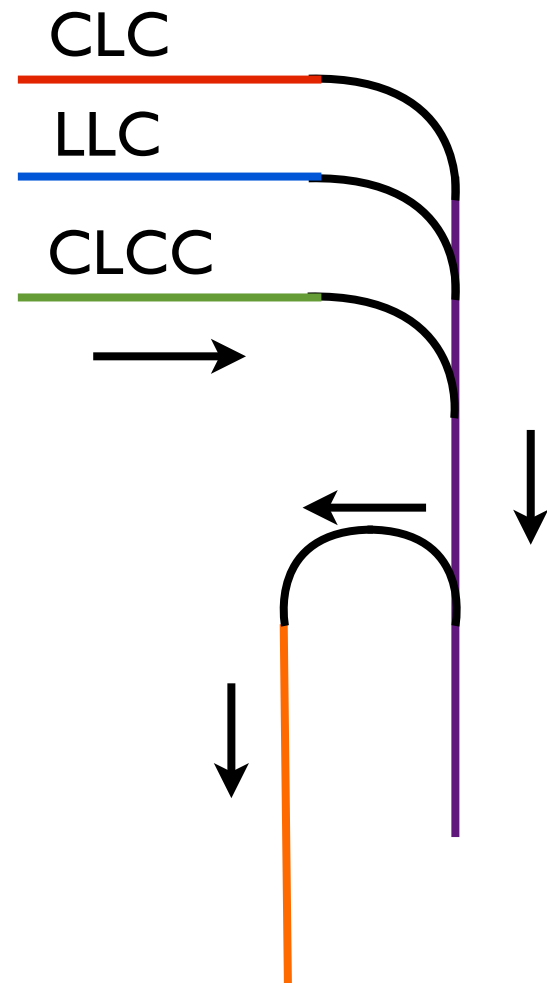
T \longrightarrow (_, T, _, LLC, _)

T \longrightarrow (_, _, T, CLCC, _)

C \longrightarrow (_, _, _, _, C)

L \longrightarrow (_, _, _, _, L)

Modelling



(T, T, T, _, _)

T → (T, _, _, CLC, _)

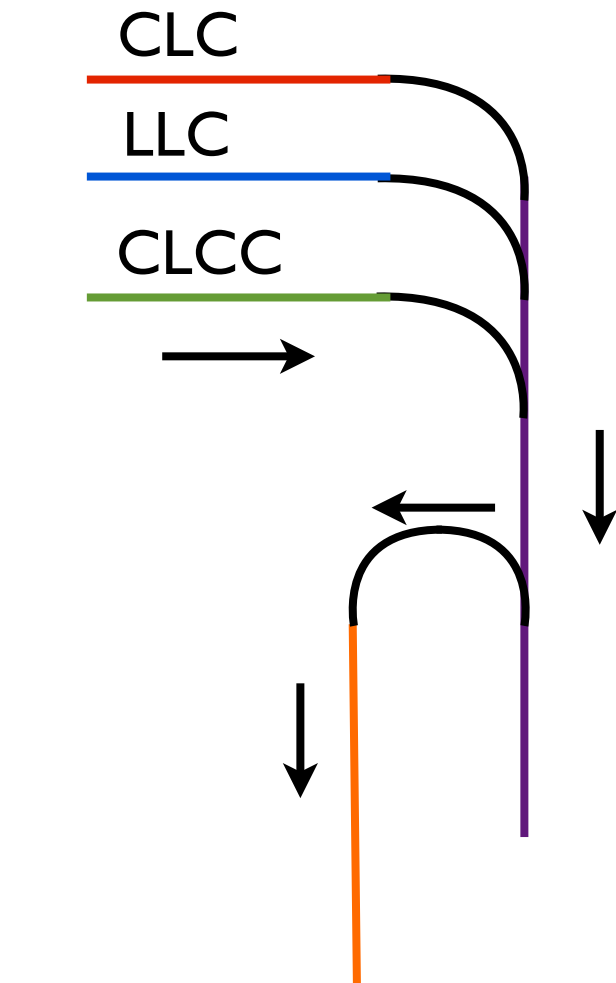
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Modelling



$(T, T, T, _, _)$



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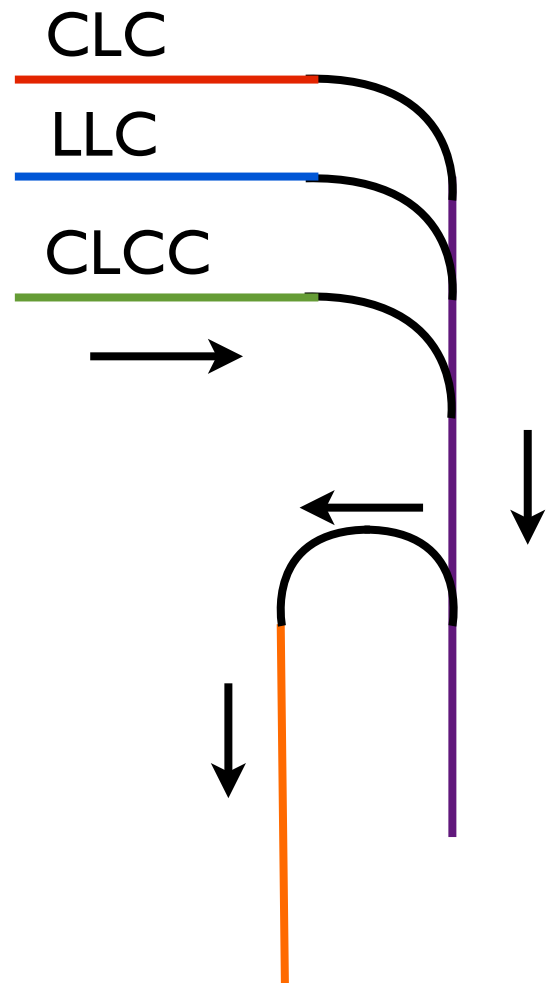
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$$T \longrightarrow (T, _, _, CLC, _)$$

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$$(T, T, T, _, _) \longrightarrow \dots \xrightarrow{?} (T, T, T, _, (CCL)^+)$$

Problems

Problems

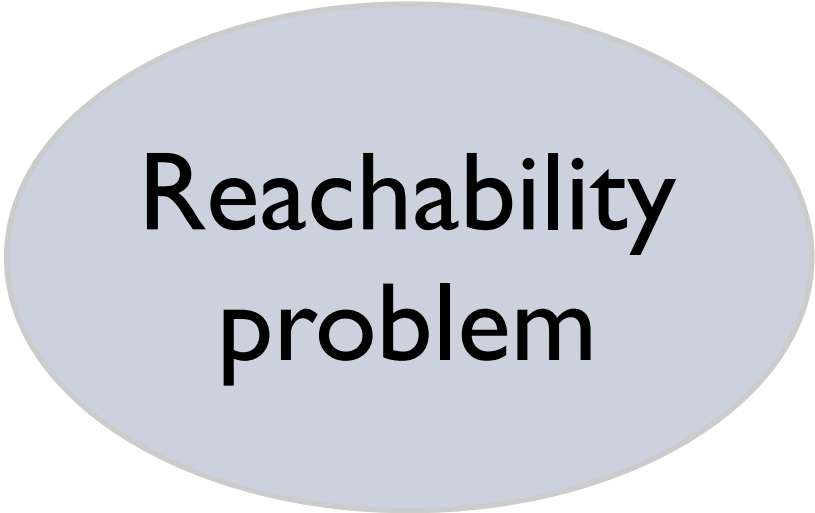


Expressivity

Problems



Expressivity



Reachability
problem

Problems



Expressivity

Reachability
problem

Equivalence
problem

Main result (expressivity)

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**Multi-Stack
Automata**

Main result (expressivity)

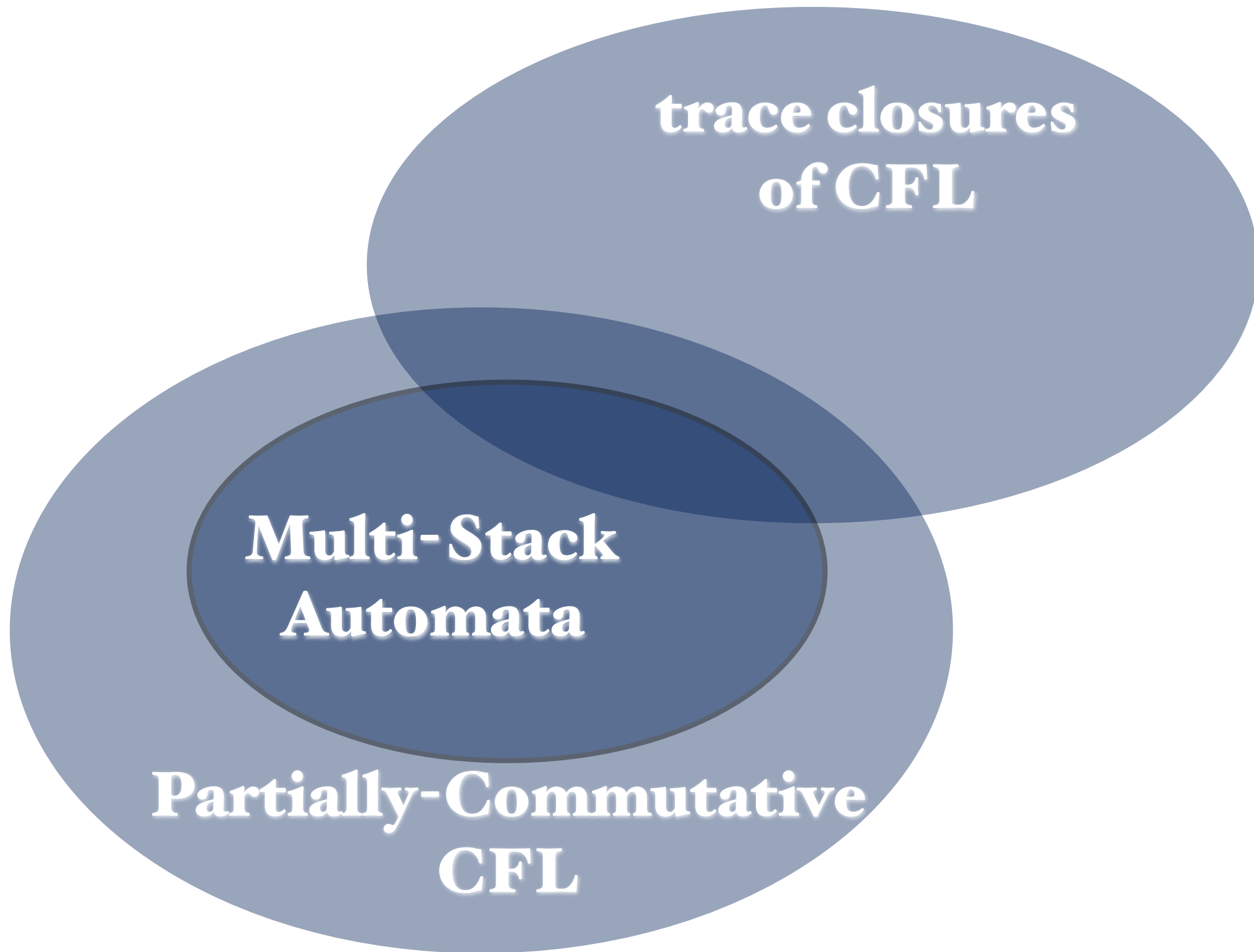


**Multi-Stack
Automata**

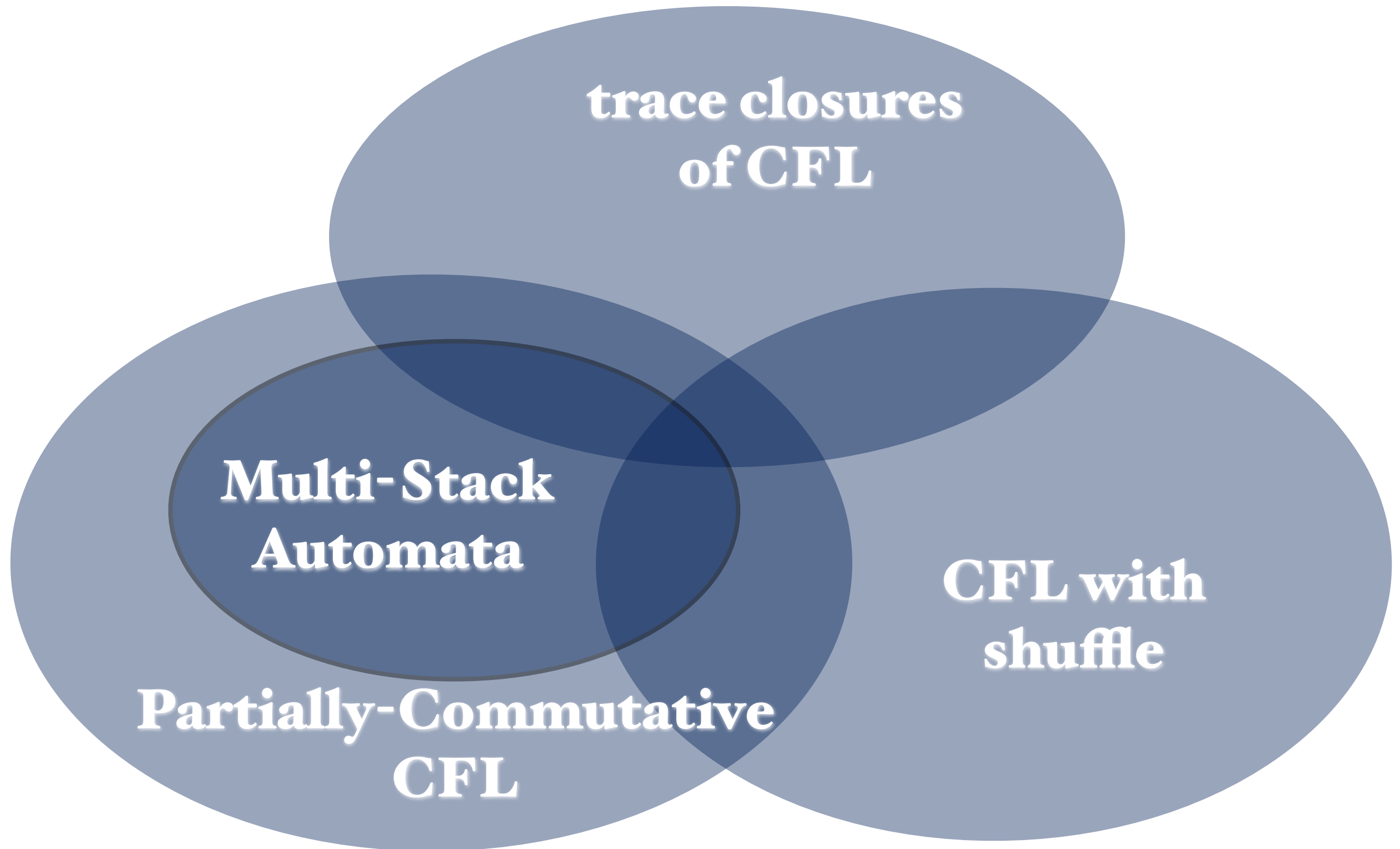
The diagram consists of two concentric ovals. The inner oval is dark blue and contains the text 'Multi-Stack Automata'. The outer oval is a lighter shade of blue and contains the text 'Partially-Commutative CFL'. This visualizes that Multi-Stack Automata are a subset of Partially-Commutative CFL.

**Partially-Commutative
CFL**

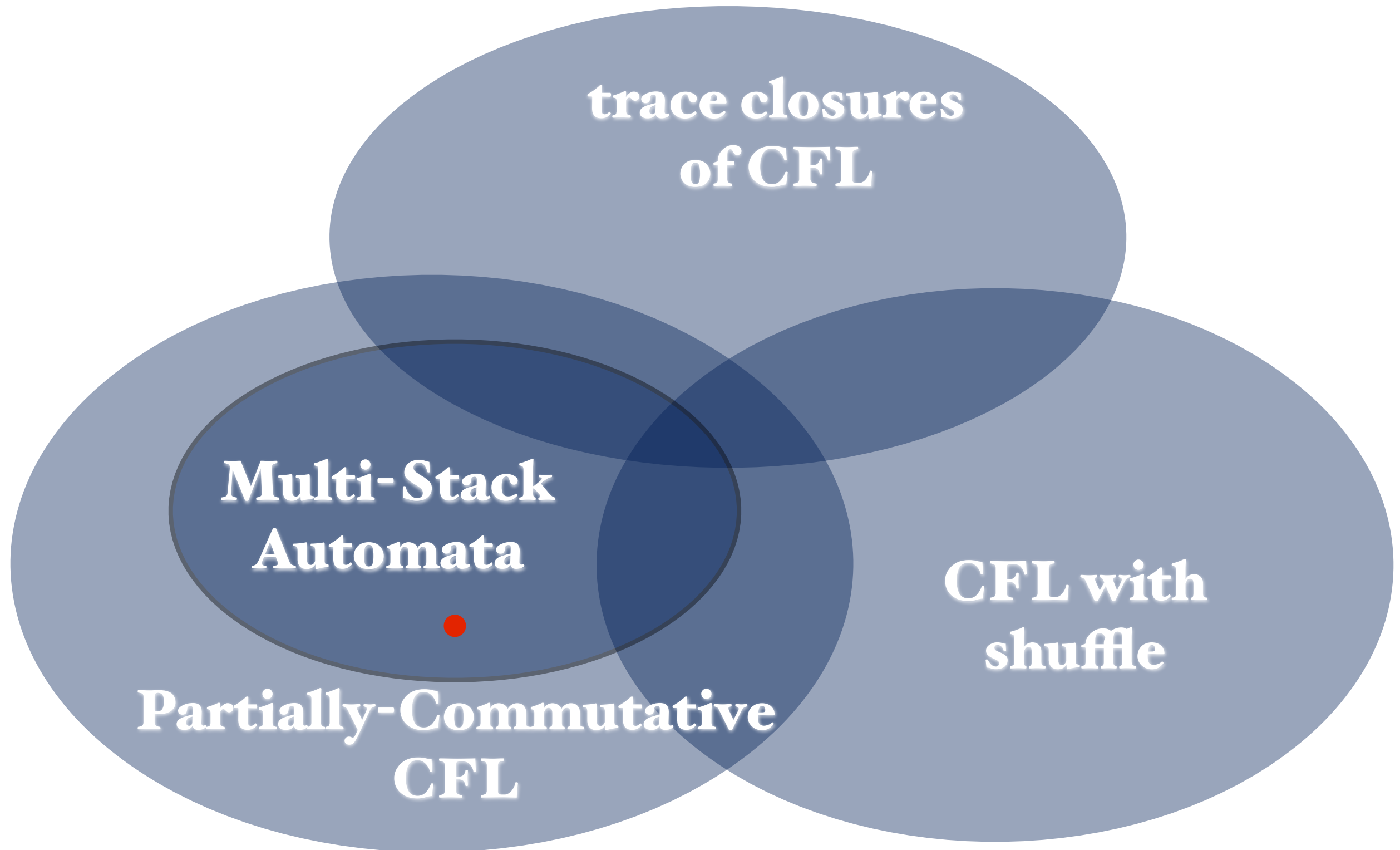
Main result (expressivity)



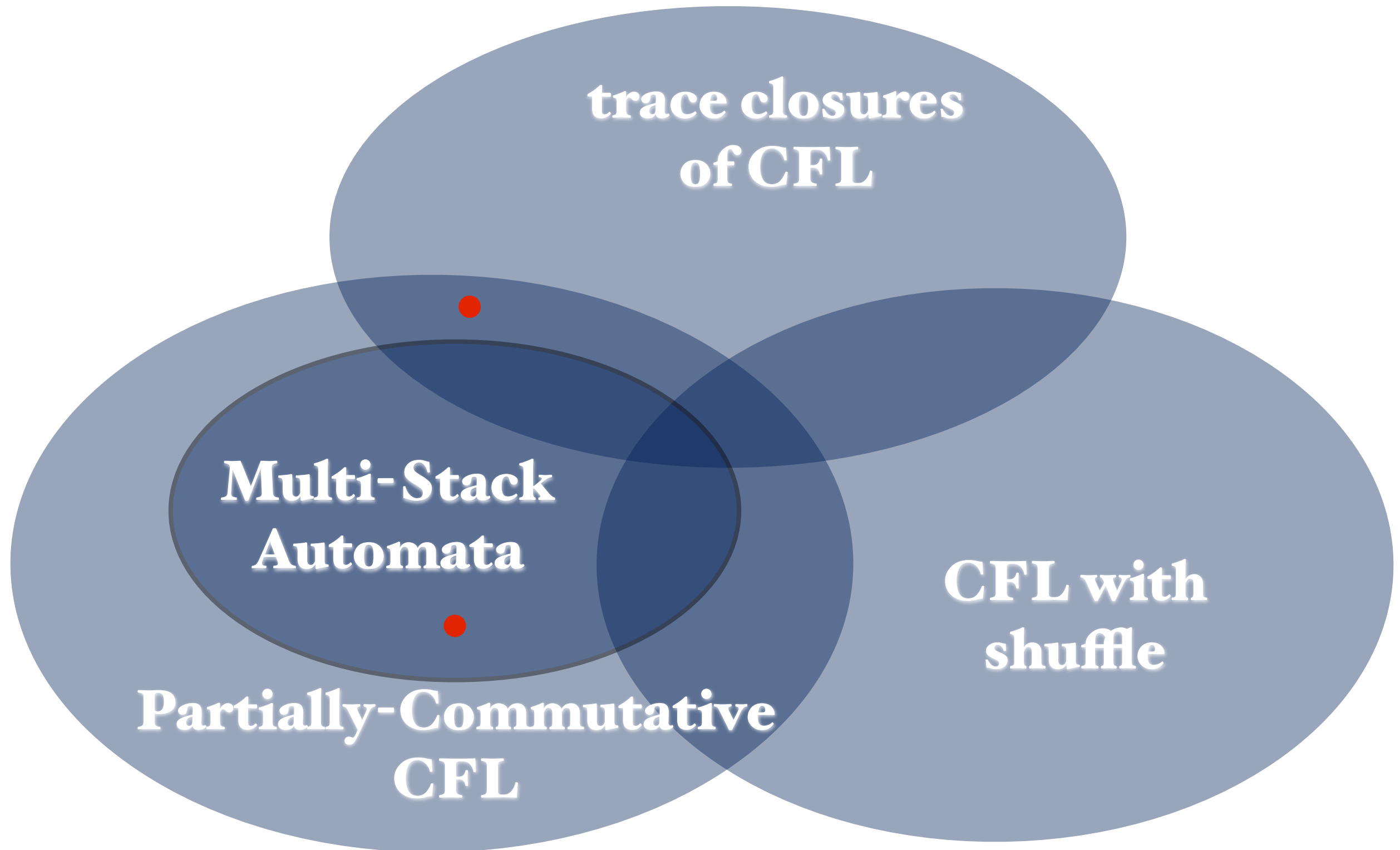
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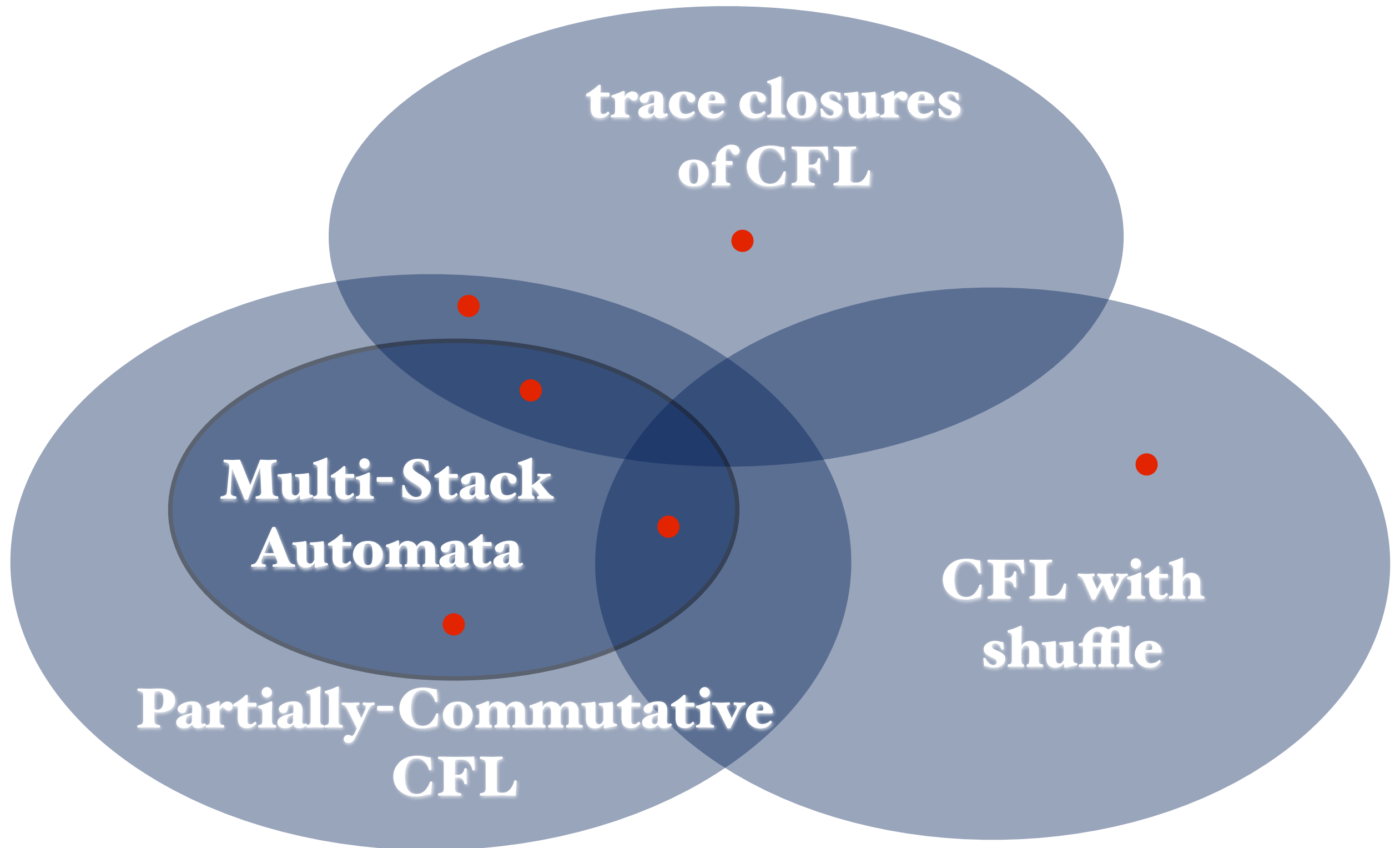
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Reachability problem

Reachability problem

Given: A Multi-Stack Automaton and two its configurations: s (source) and t (target)

Reachability problem

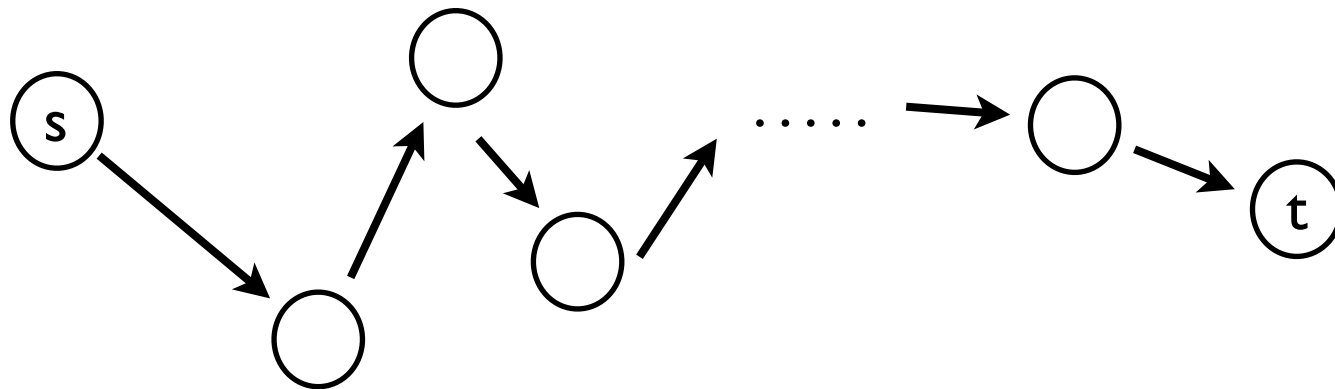
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Question: Decide whether there is a path from s to t

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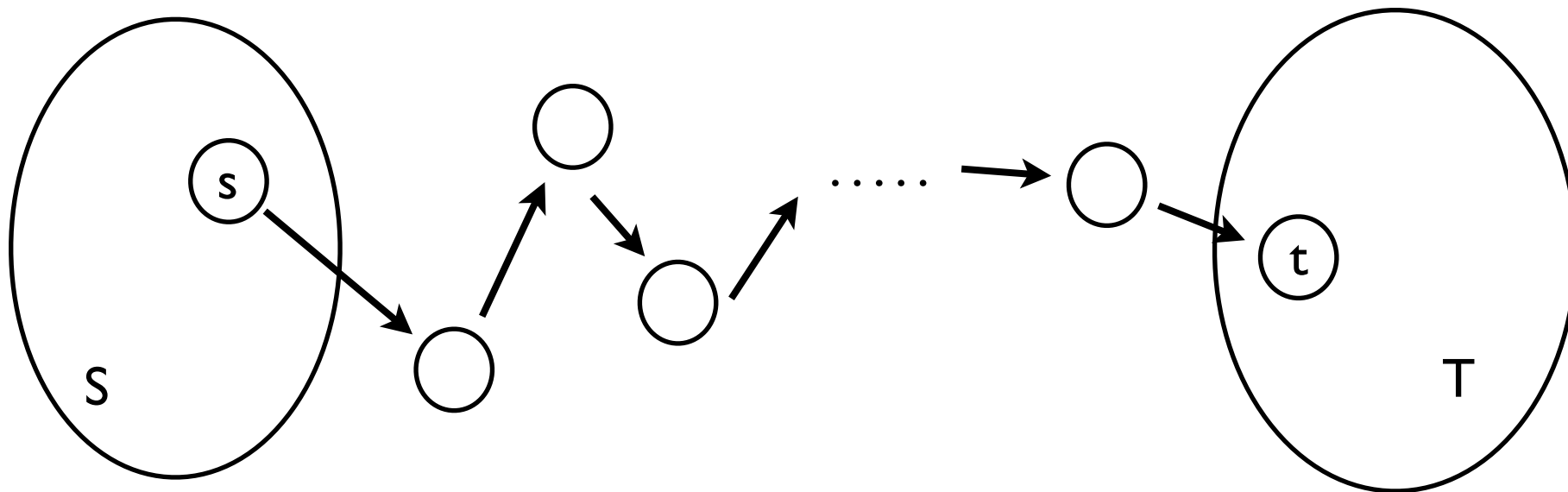
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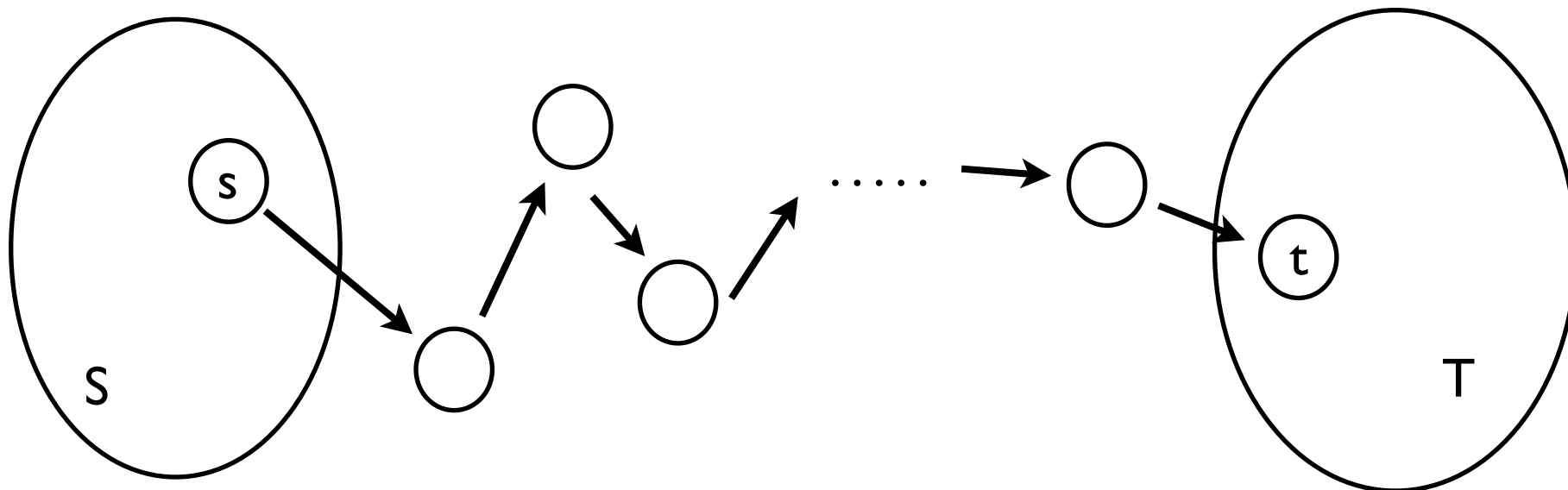
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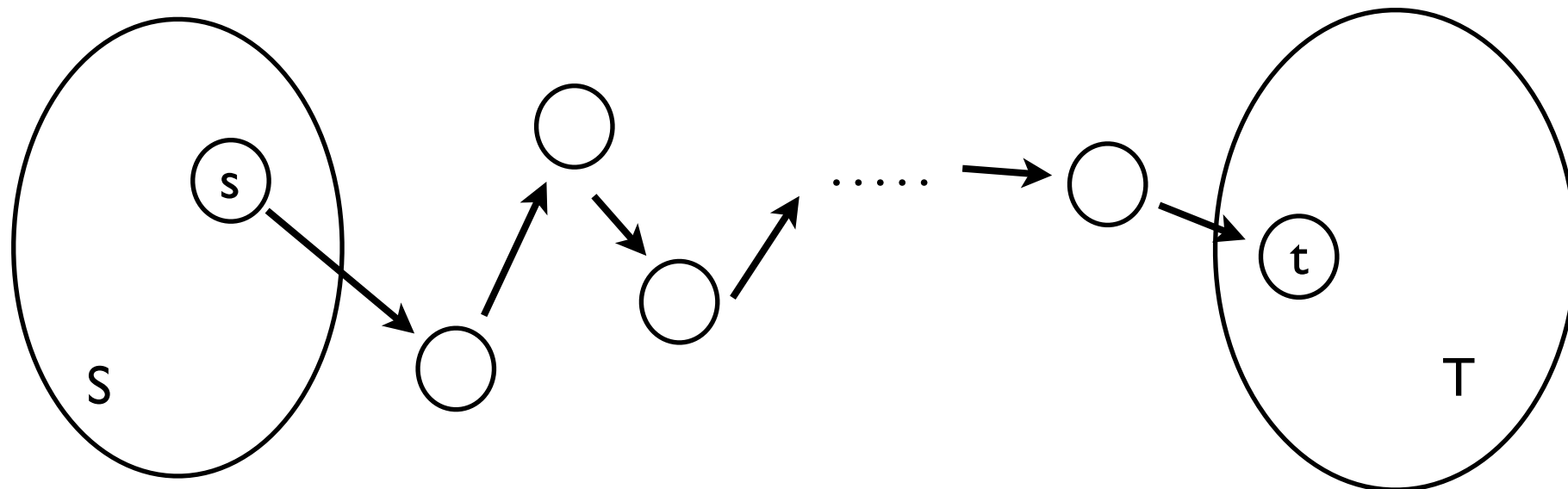


Regular: $((AB)^*, CD) \cup (A^*B, C^*D^*)$

Reachability problem

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Regular: $((AB)^*, CD) \cup (A^*B, C^*D^*)$

Not regular: (A^n, B^n)

Main result (reachability)

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Theorem:

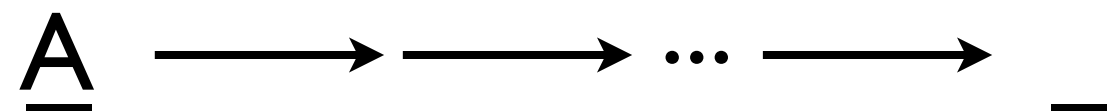
Reachability problem for *normed* MSA with regular source and target sets is NP-complete

Main result (reachability)

Theorem:

Reachability problem for *normed* MSA with regular source and target sets is NP-complete

Normedness: for every letter A



Equivalence

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Rules are labelled:

Equivalence

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The idea: type of the action (c - moving of a car)

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Bisimulation equivalence:

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Bisimulation equivalence:

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$$t \xrightarrow{a} t' \xrightarrow{b} \rightarrow$$

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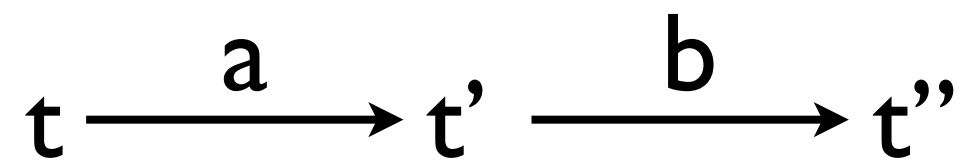
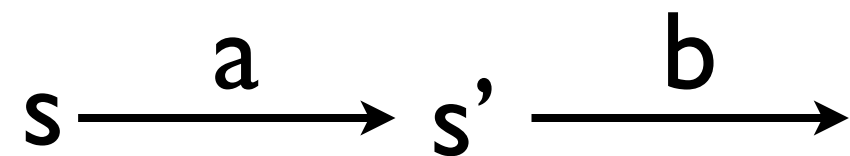
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Bisimulation equivalence is decidable

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- in time $O(n^4)$ for normed Single-Stack Automata

Main results (equivalence)

Bisimulation equivalence is decidable

- in polynomial time for normed *disjoint* Multi-Stack Automata
- in time $O(n^4)$ for normed Single-Stack Automata
- in time $O(n^3)$ for Simple Grammars

Thank you!