

Single use restriction

A new trend in transformations with atoms

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Joint work with Mikołaj Bojańczyk

Work in progress

Deterministic register automata

Fix a countably infinite set of *atoms*.

They can only be compared for equality.

$$A = \{1\ 2\ 3\ \dots\}$$

model by Kaminski, Francez 1994

Deterministic register automata

The first letter appears again

1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{\text{start}}(\perp)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{scan}()$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{\text{scan}}(1)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{\text{scan}}(1)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{\text{scan}}(1)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

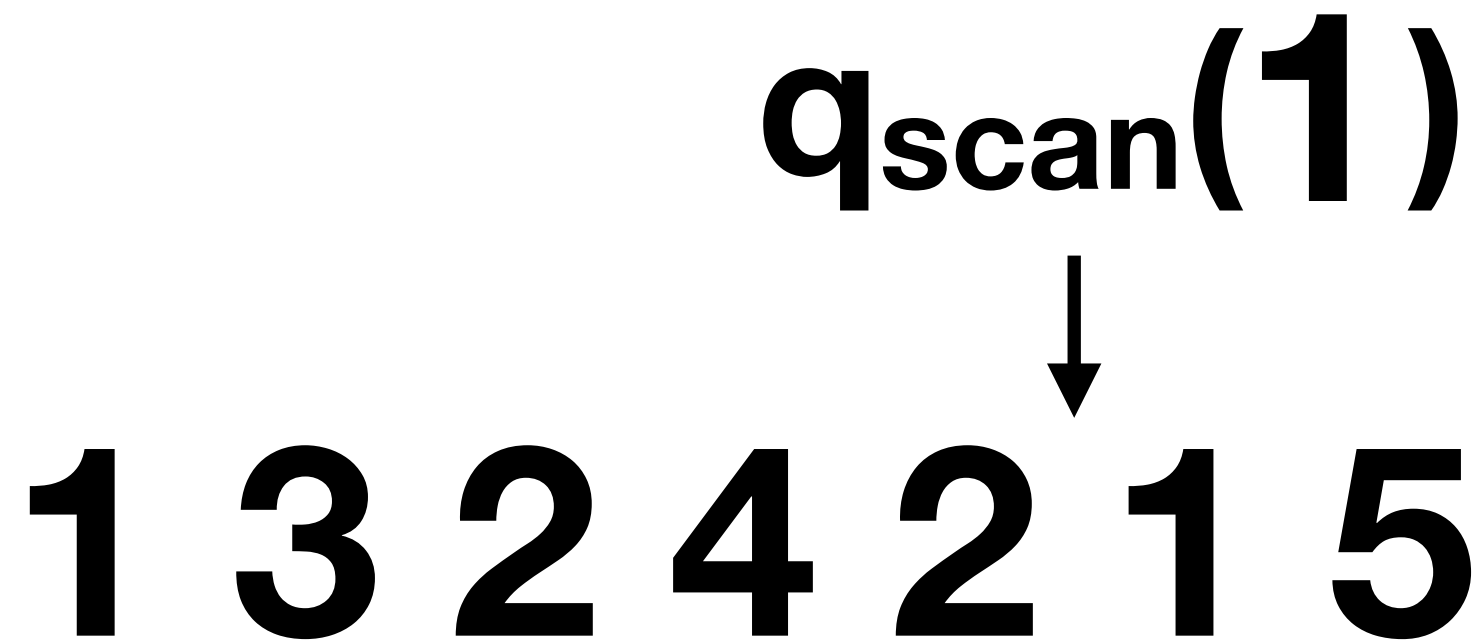
$q_{\text{scan}}(1)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again



Deterministic register automata

The first letter appears again

$q_{\text{found}}(\perp)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

$q_{\text{found}}(\perp)$



1 3 2 4 2 1 5

Deterministic register automata

The first letter appears again

1 3 2 4 2 1 5 

Single use restriction

Every read access destroys the contents of the register.

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$Q_{\text{start}}(\perp)$



1 3 2 4 2 1 5

Single use restriction

Every read access destroys the contents of the register.

$q_{\text{scan}}(1)$



1 3 2 4 2 1 5

Single use restriction

Every read access destroys the contents of the register.

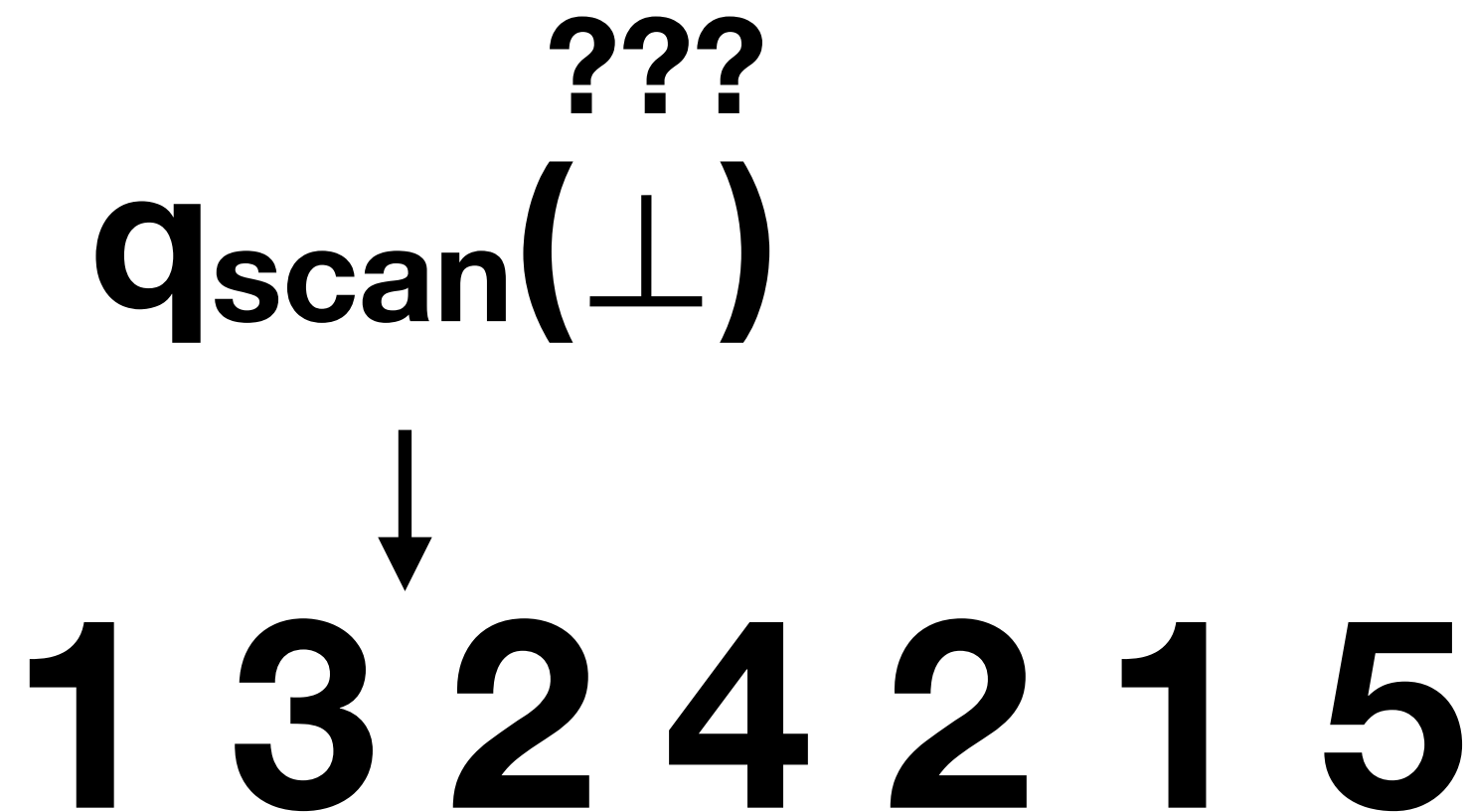
$q_{\text{scan}}(\perp)$



1 3 2 4 2 1 5

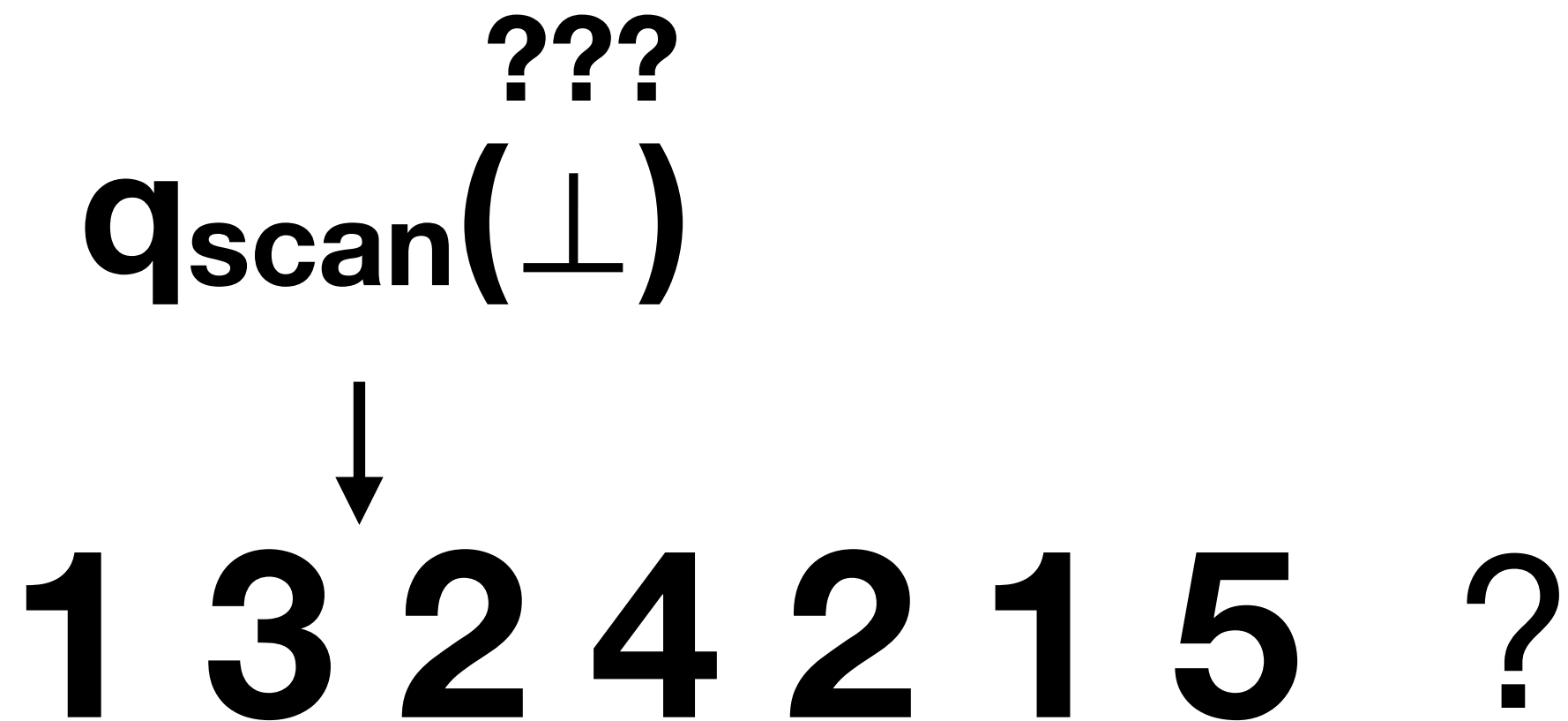
Single use restriction

Every read access destroys the contents of the register.

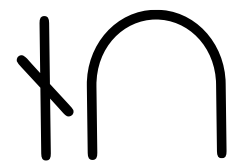


Single use restriction

Every read access destroys the contents of the register.



Single use register automata

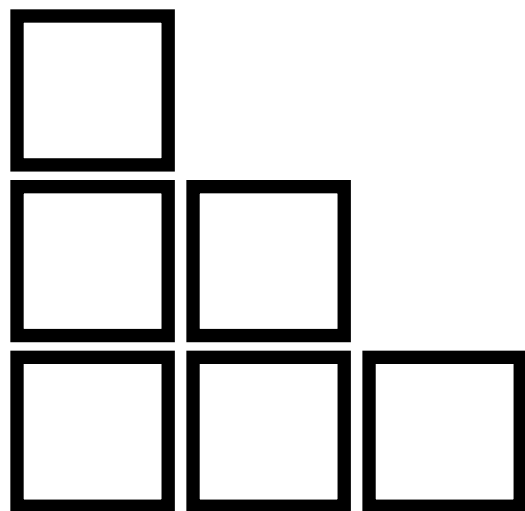


The first letter
appears again

Register automata

Single use register automata

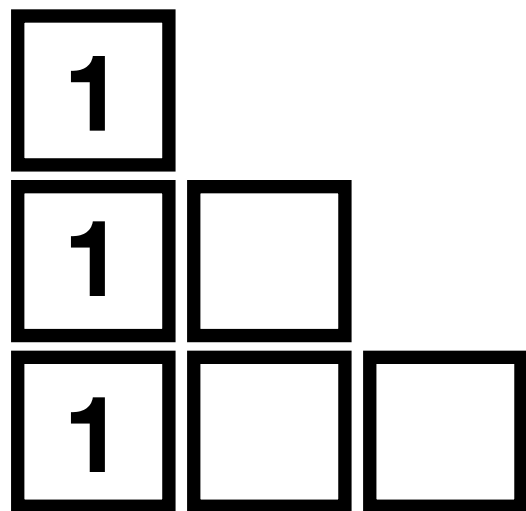
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

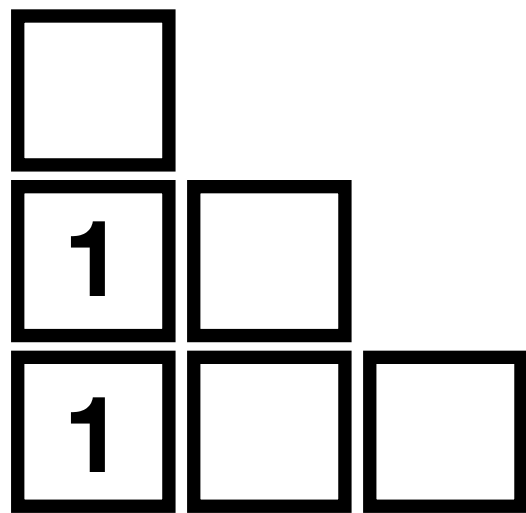
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1 2 3 1 3 4

Single use register automata

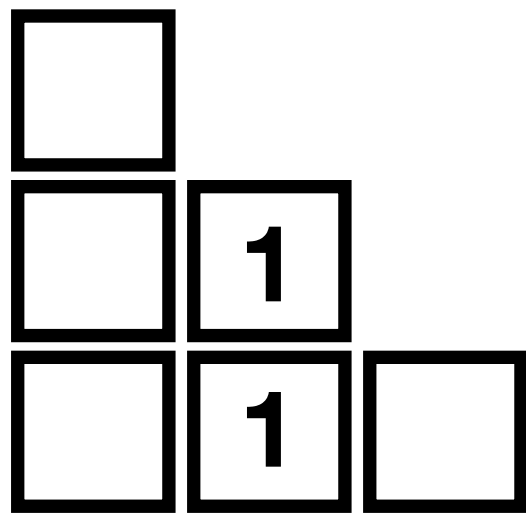
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

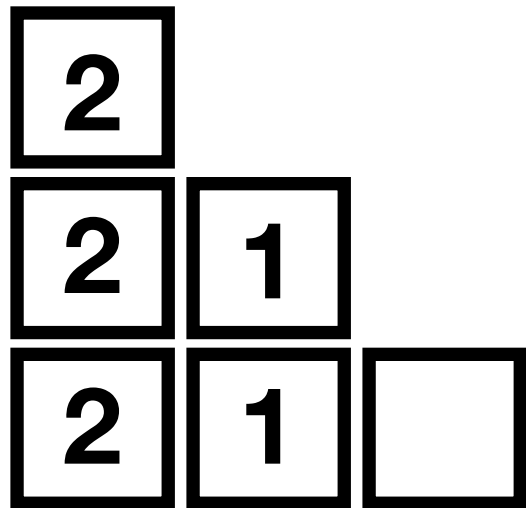
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

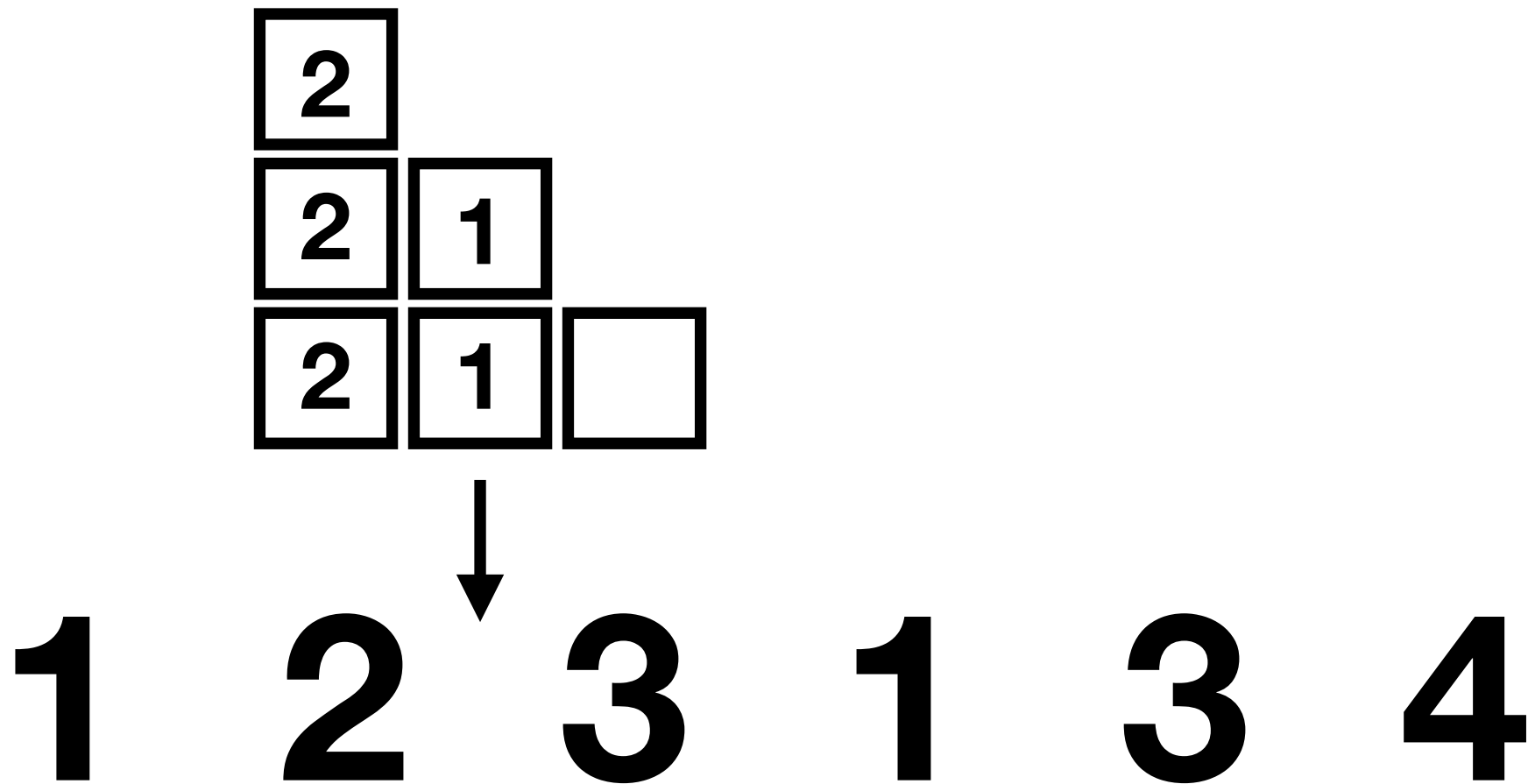
There are at most 3 distinct letters



1 2 3 1 3 4

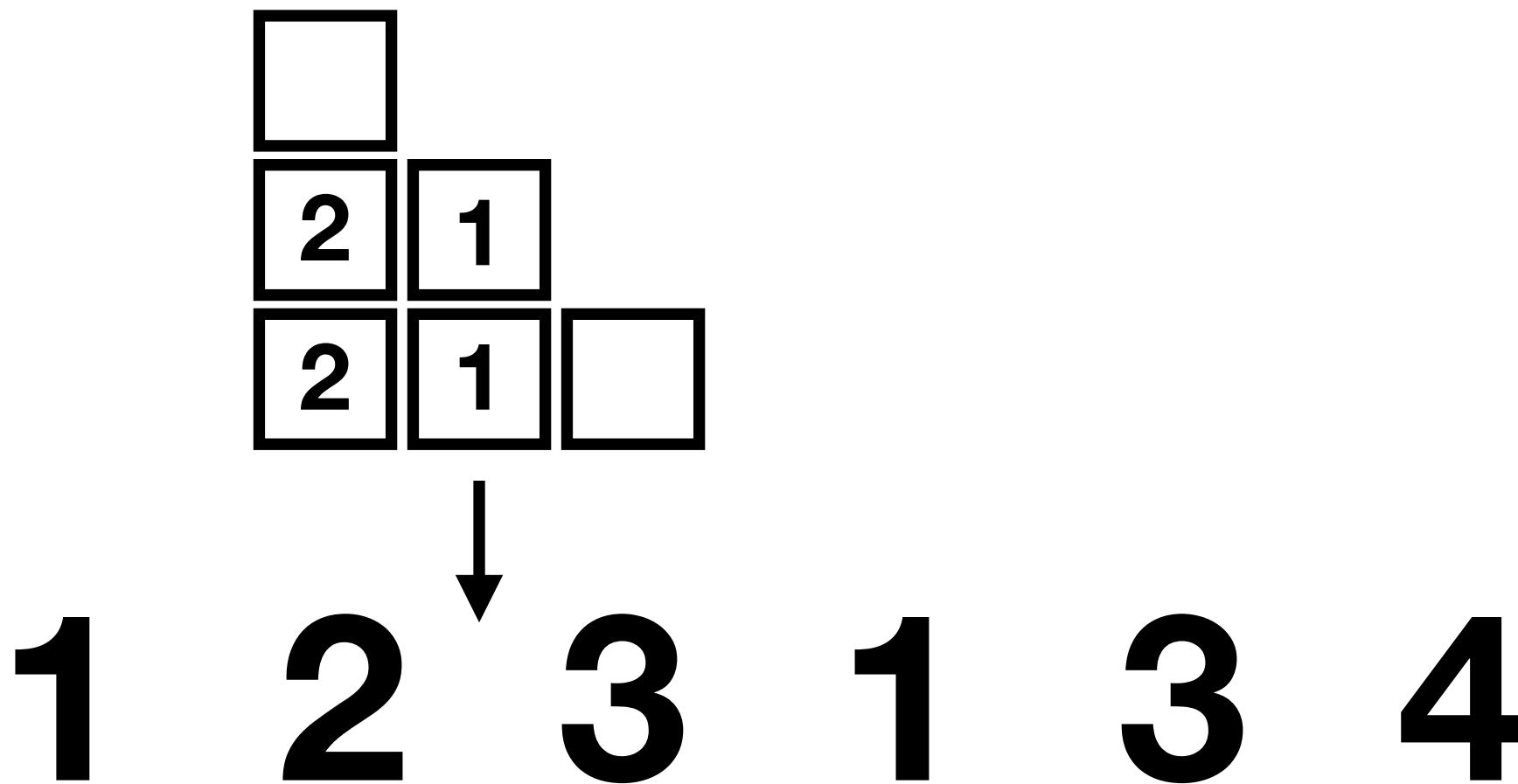
Single use register automata

There are at most 3 distinct letters



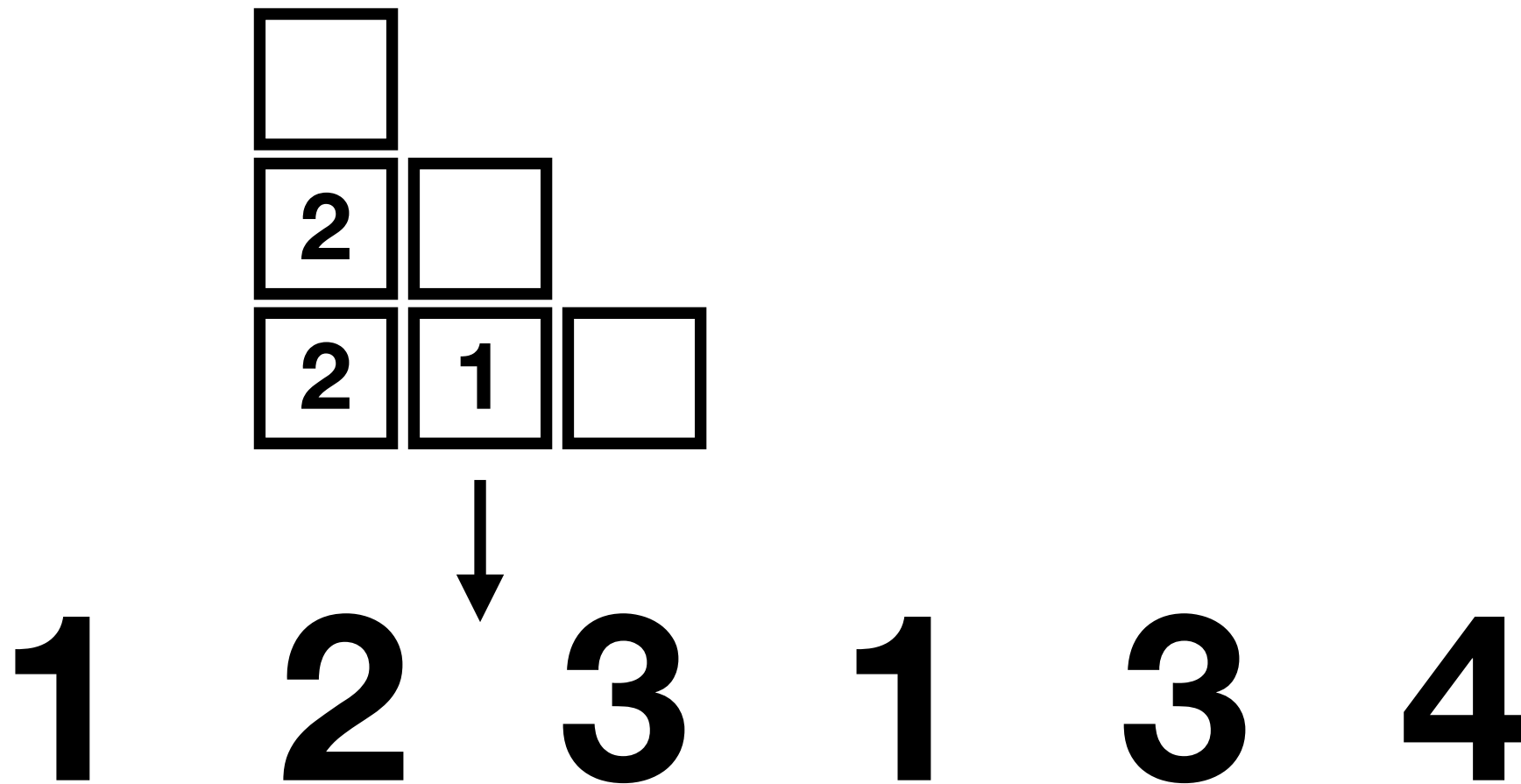
Single use register automata

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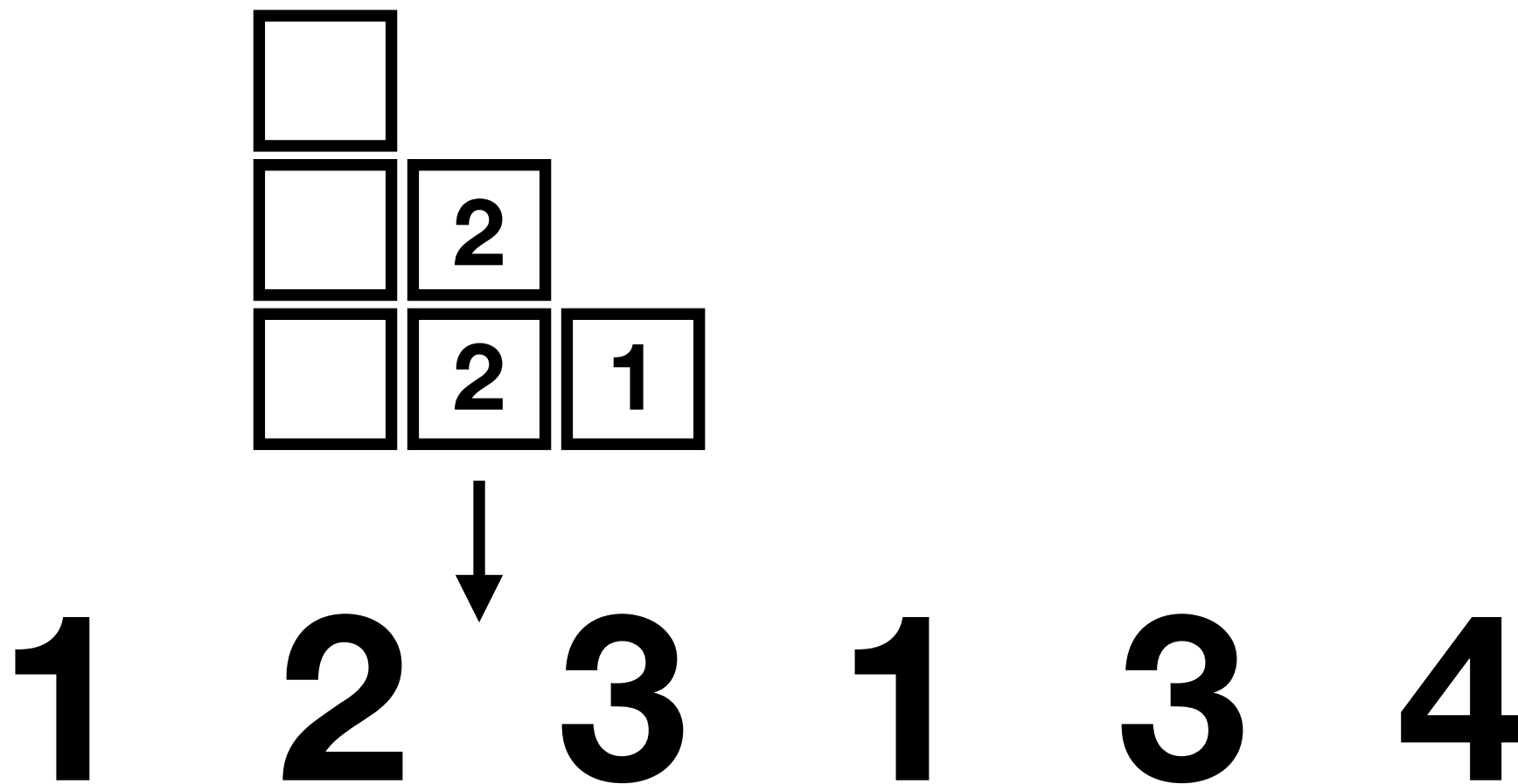
Single use register automata

There are at most 3 distinct letters



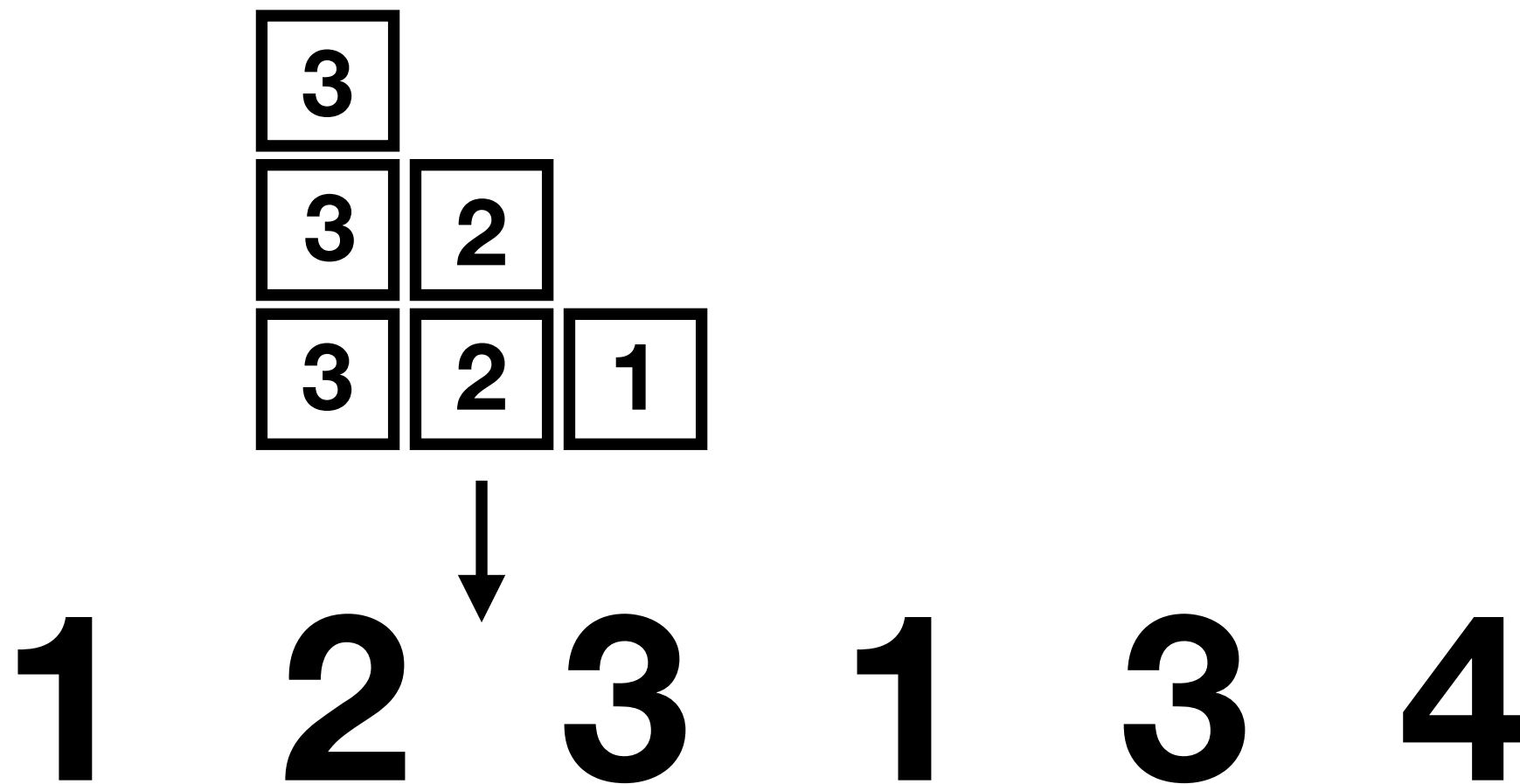
Single use register automata

There are at most 3 distinct letters



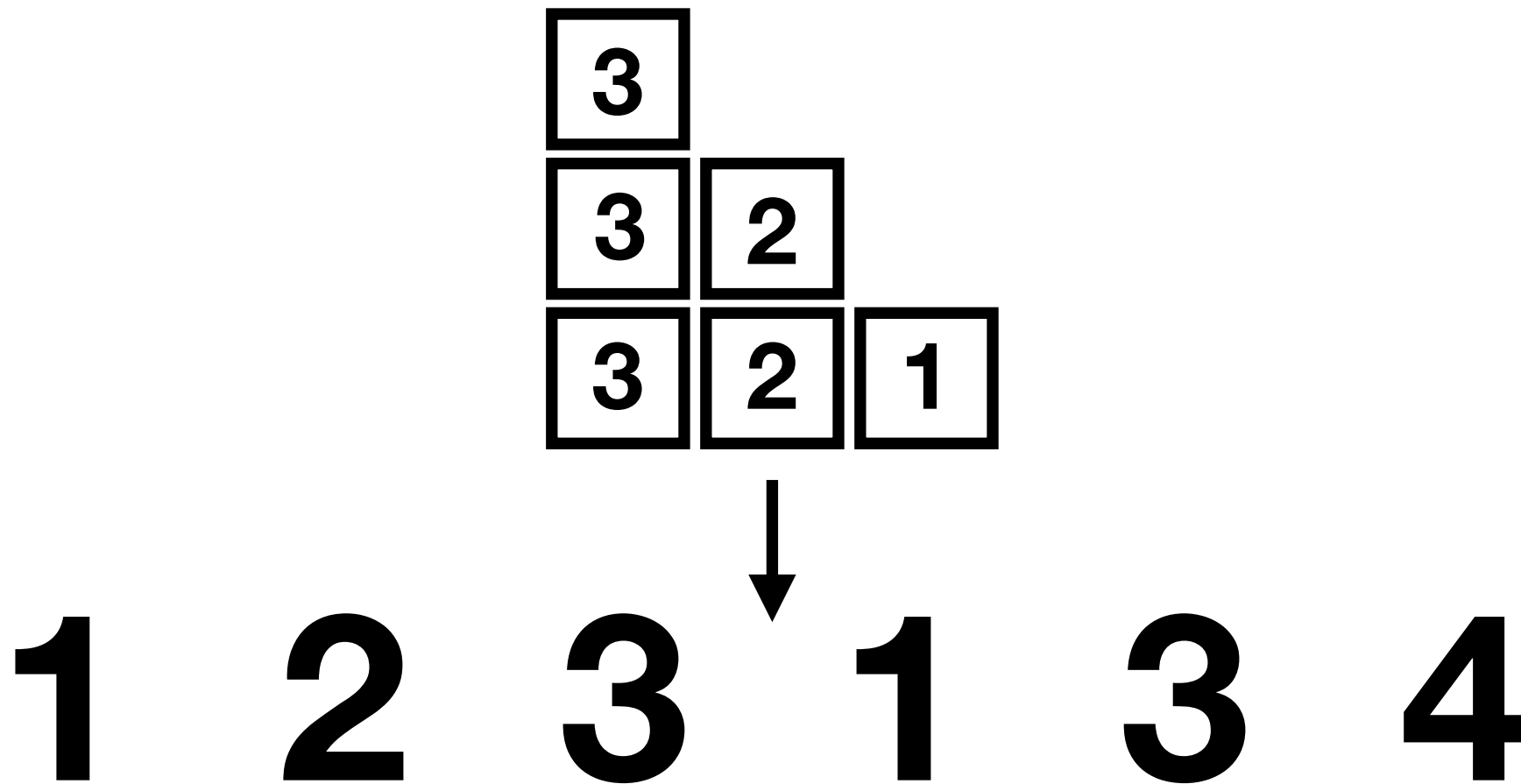
Single use register automata

There are at most 3 distinct letters



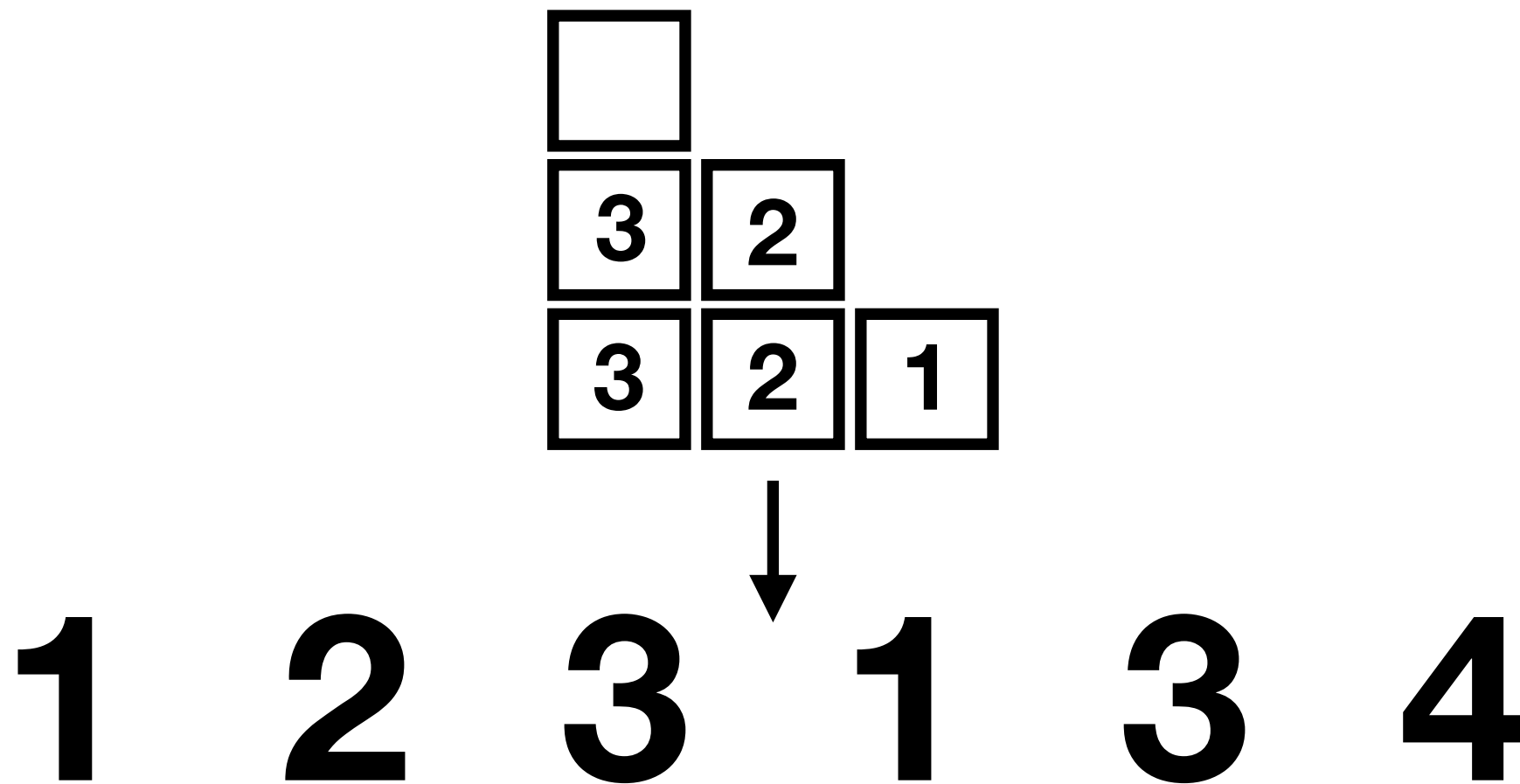
Single use register automata

There are at most 3 distinct letters



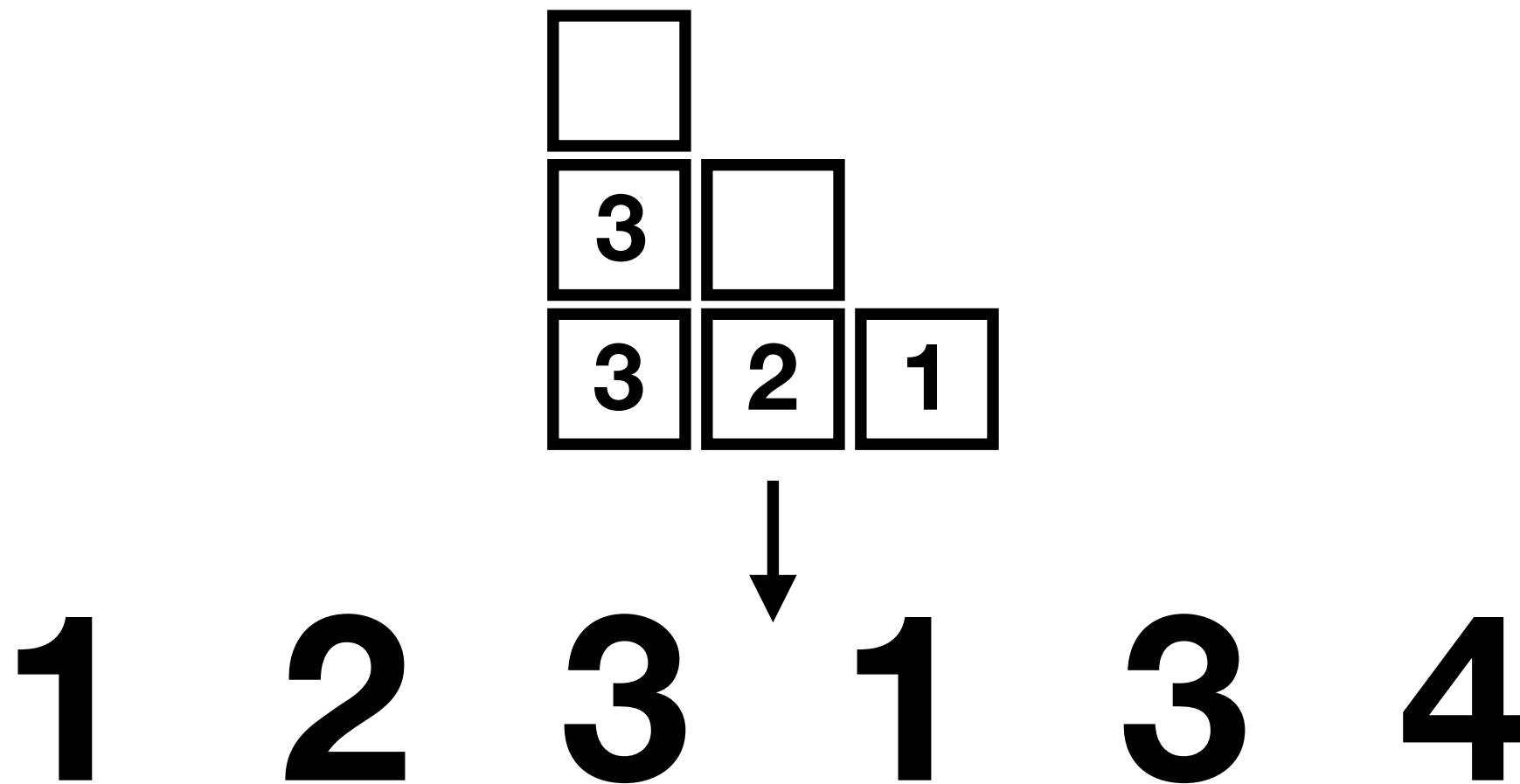
Single use register automata

There are at most 3 distinct letters



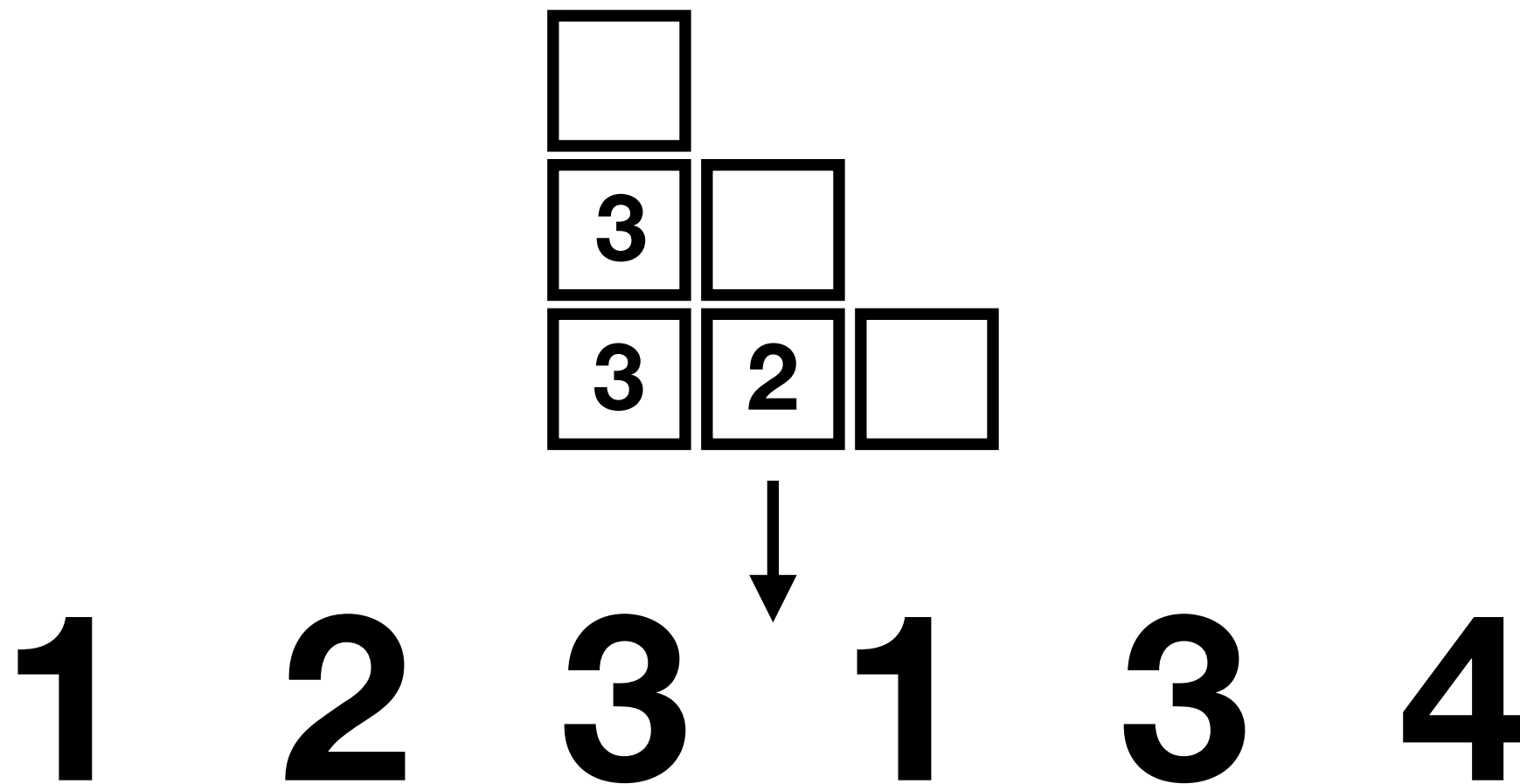
Single use register automata

There are at most 3 distinct letters



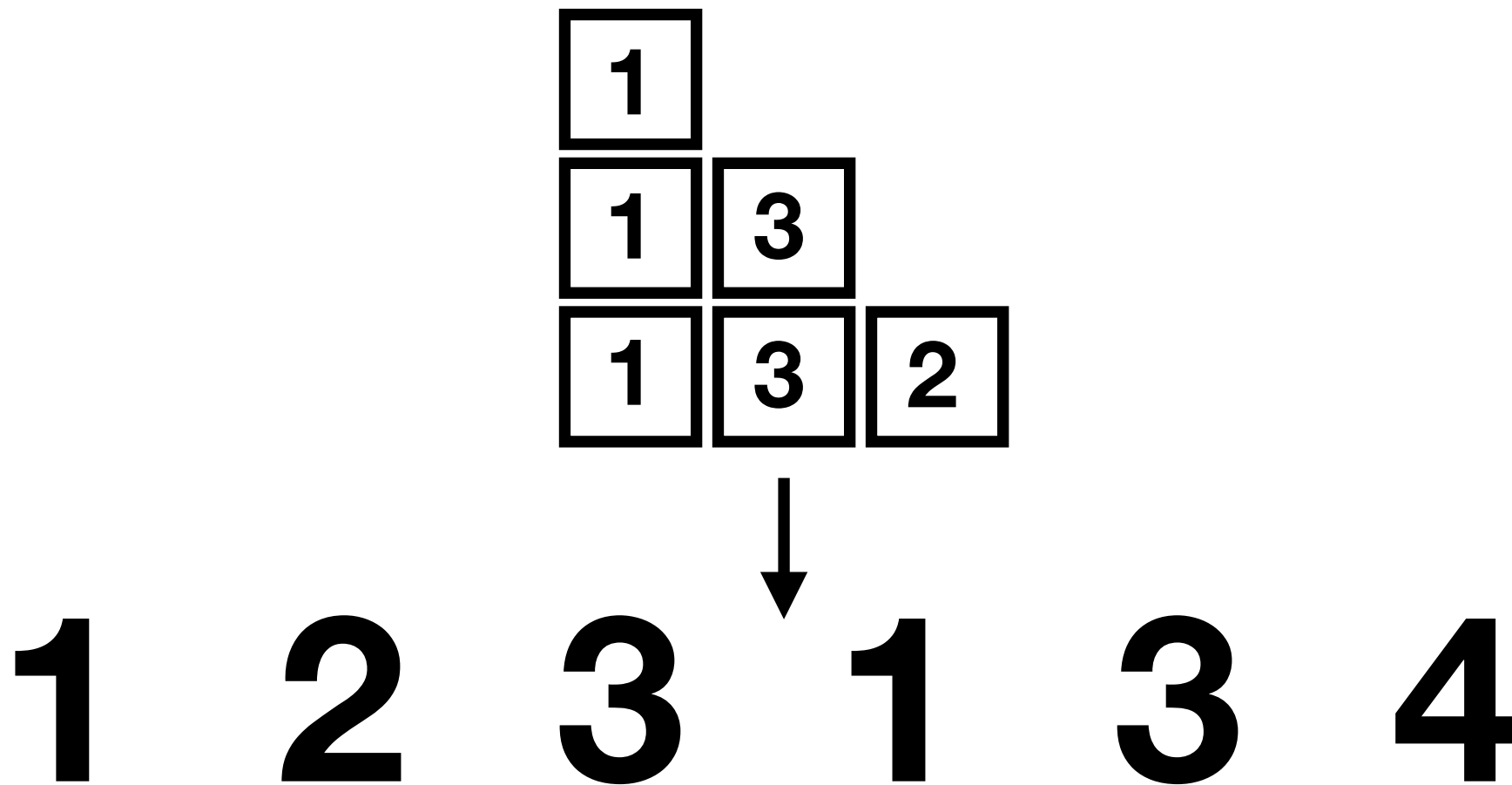
Single use register automata

There are at most 3 distinct letters



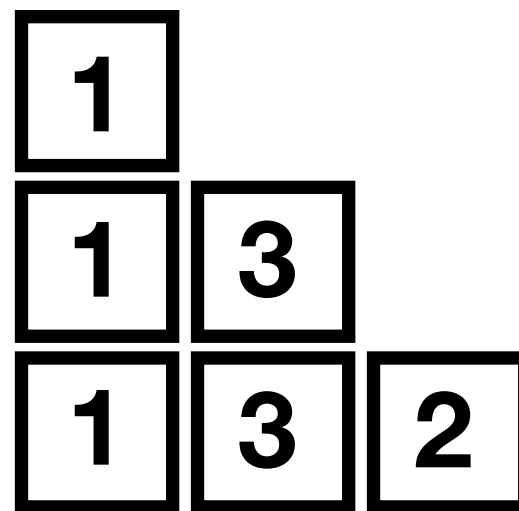
Single use register automata

There are at most 3 distinct letters



Single use register automata

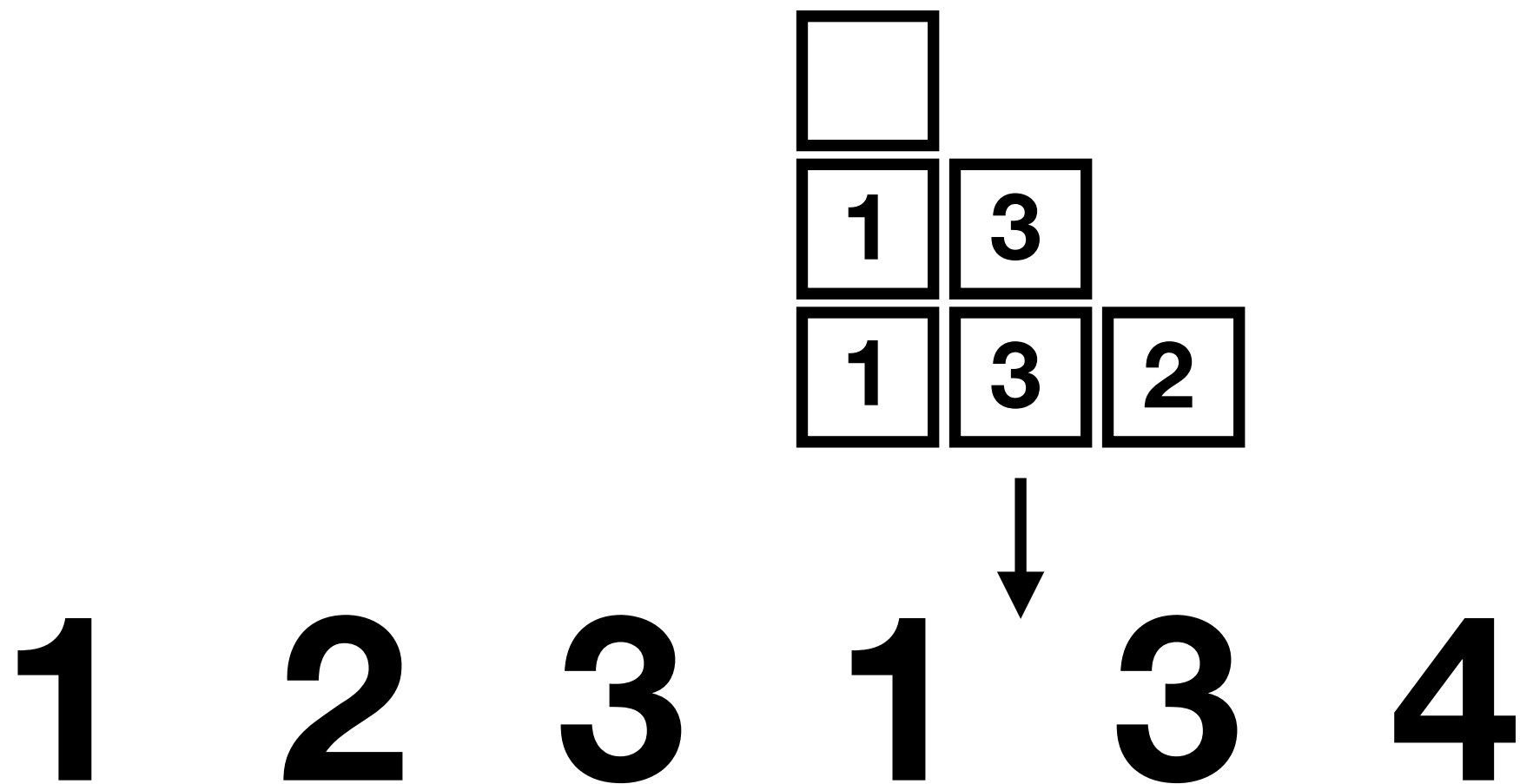
There are at most 3 distinct letters



1 2 3 1 3 4

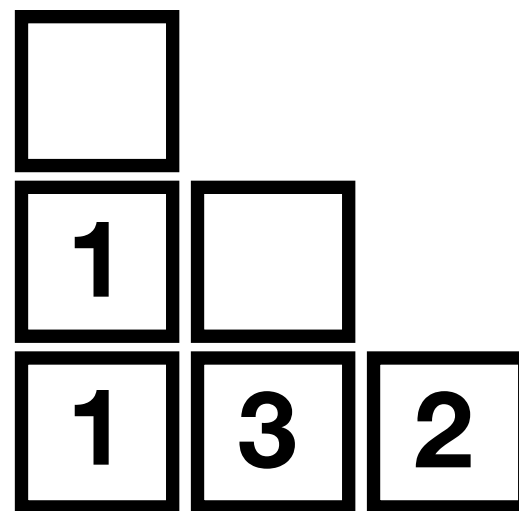
Single use register automata

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Single use register automata

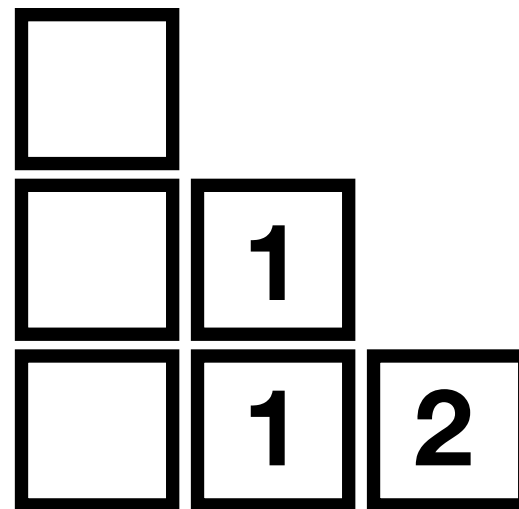
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

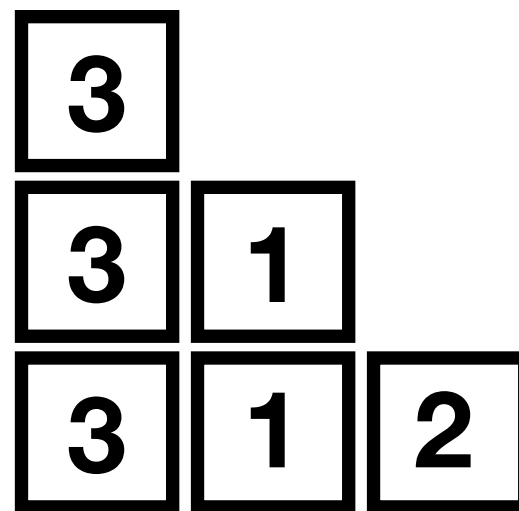
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

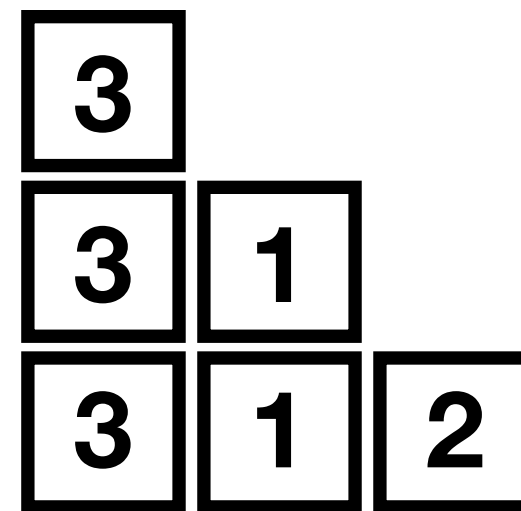
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

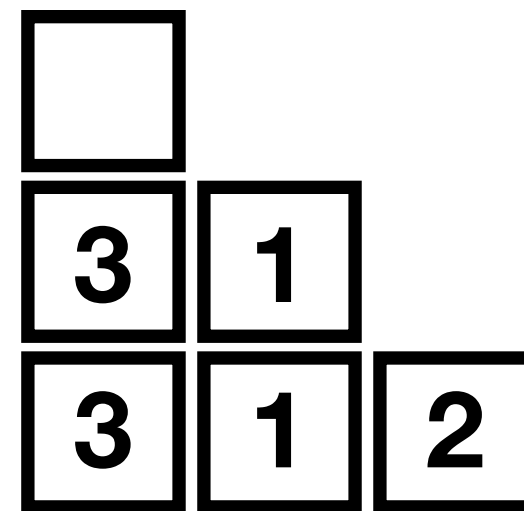
There are at most 3 distinct letters



1 2 3 1 3 4

Single use register automata

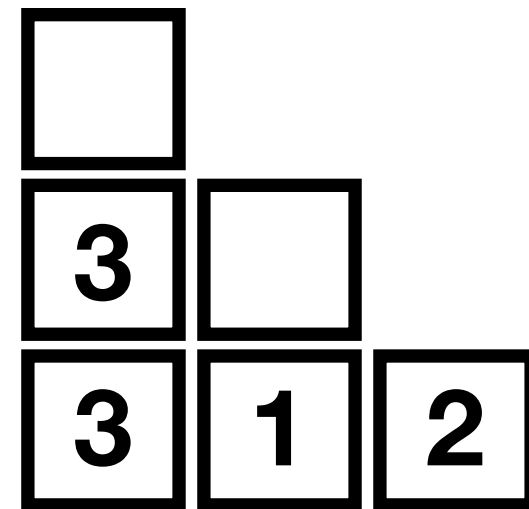
There are at most 3 distinct letters



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Single use register automata

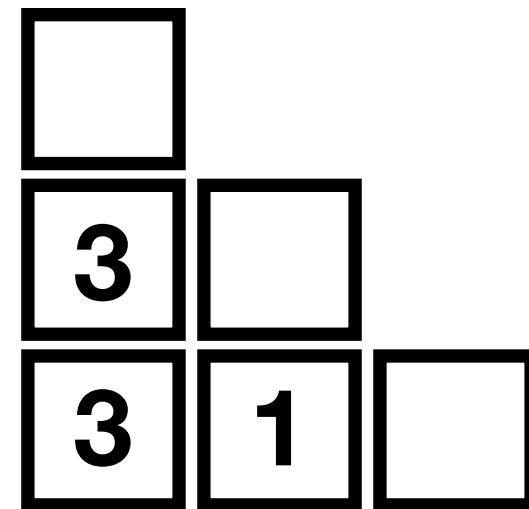
There are at most 3 distinct letters



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Single use register automata

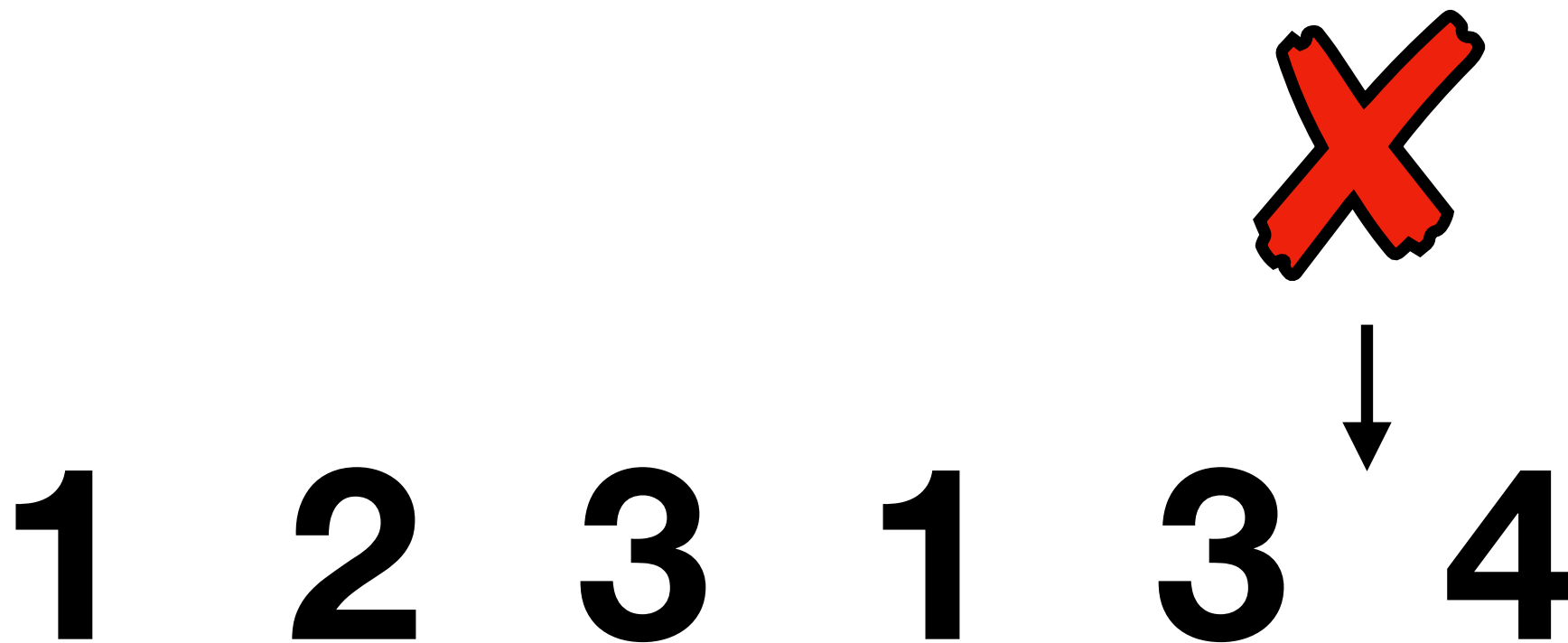
There are at most 3 distinct letters



1 2 3 1 3 4


Single use register automata

There are at most 3 distinct letters



Single use register automata

There are at most 3 distinct letters

1 2 3 1 3 4 

Rigidly guarded MSO~

Colcombet, Ley, Puppis 2015



Orbit-finite semigroups

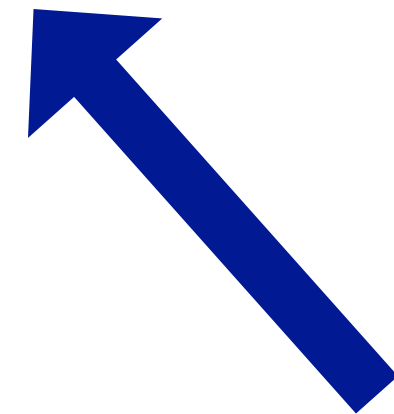
Bojańczyk 2013



Single-use register automata



Two-way single-use register automata



Robustness – a case for the single use restriction

All the following models recognise
different classes of languages

Multiple-use, deterministic register automata

Multiple-use, 2-way deterministic register automata

MSO~

Transducers

with atoms

An example

Remove repetitions from the input

1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3 2

1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3 2

1 1 2 3 3 3 2 2

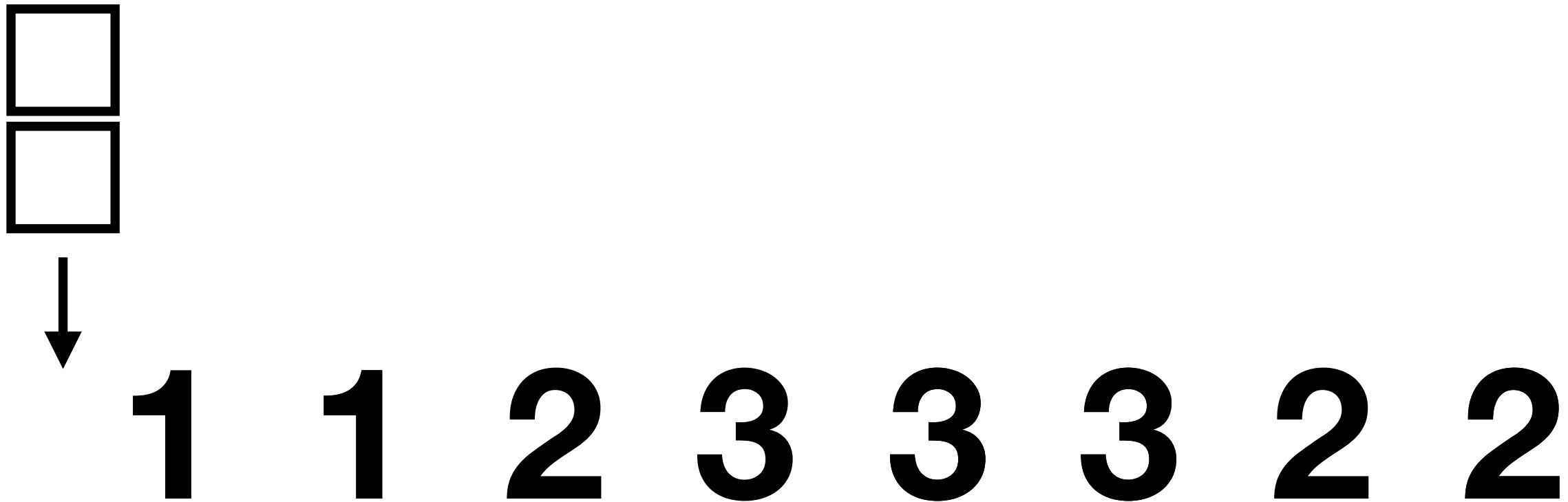
An example

Remove repetitions from the input

1 1 2 3 3 3 2 2

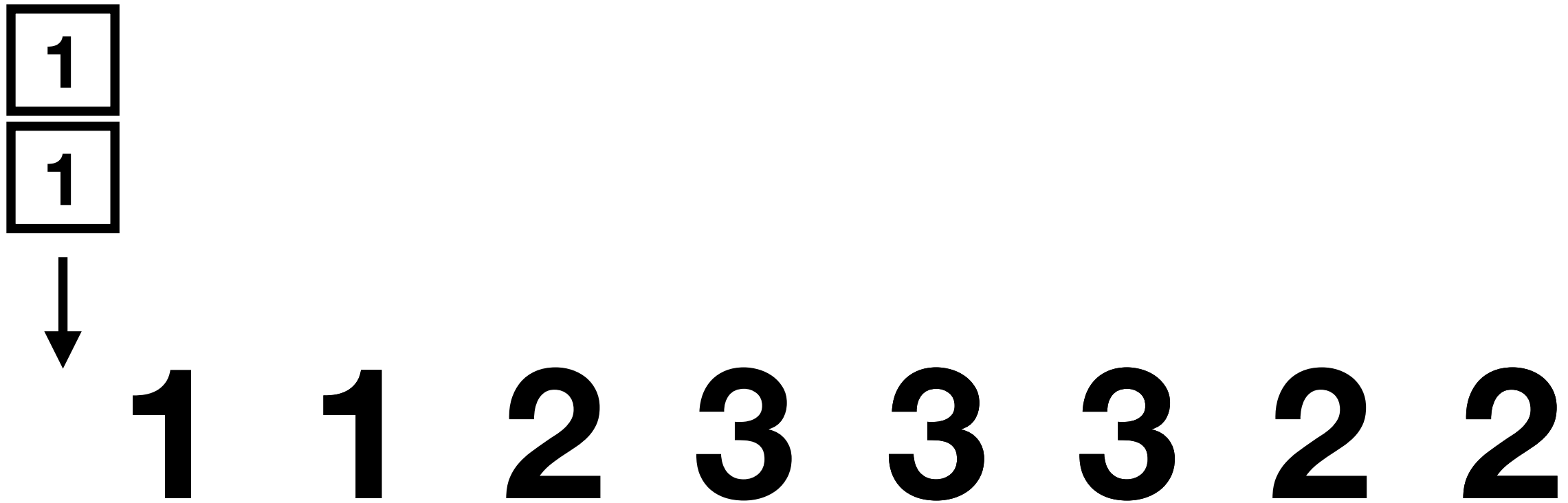
An example

Remove repetitions from the input



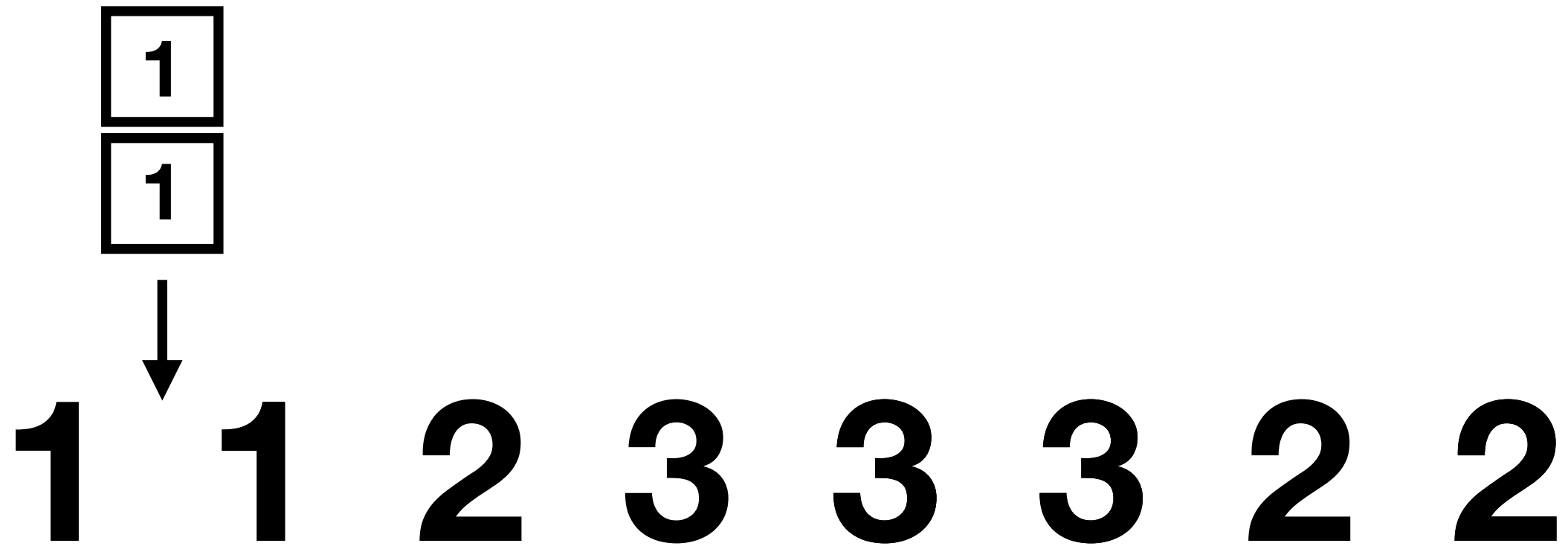
An example

Remove repetitions from the input



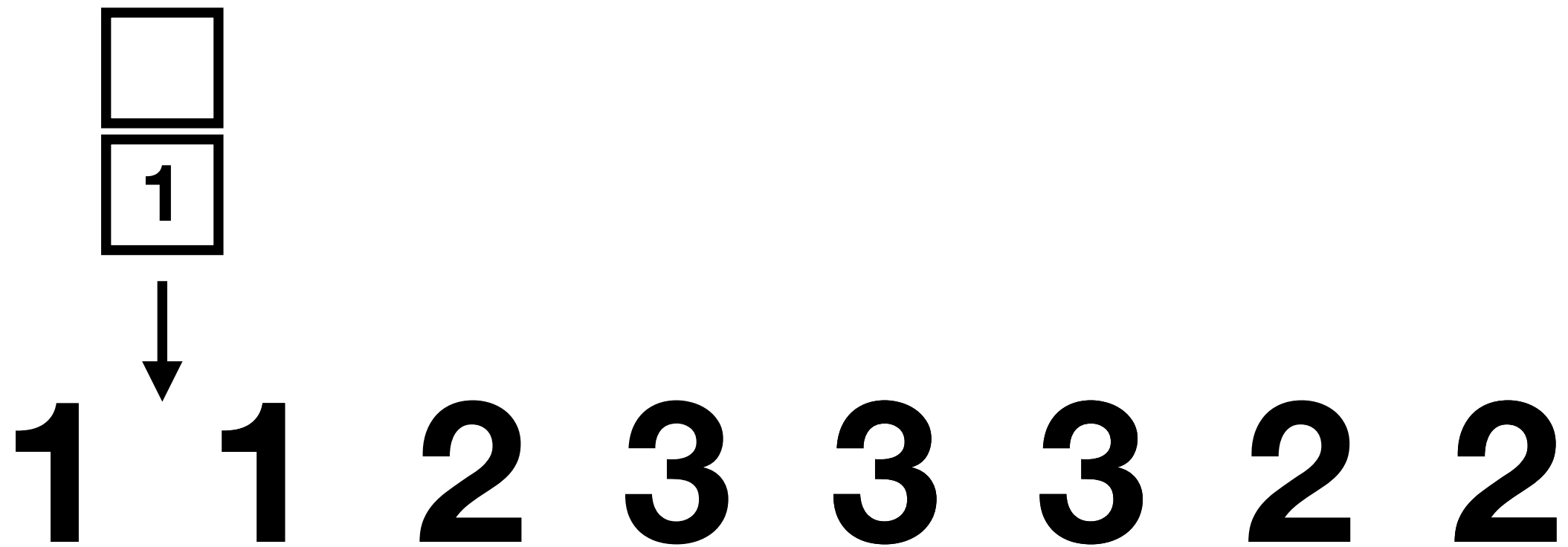
An example

Remove repetitions from the input



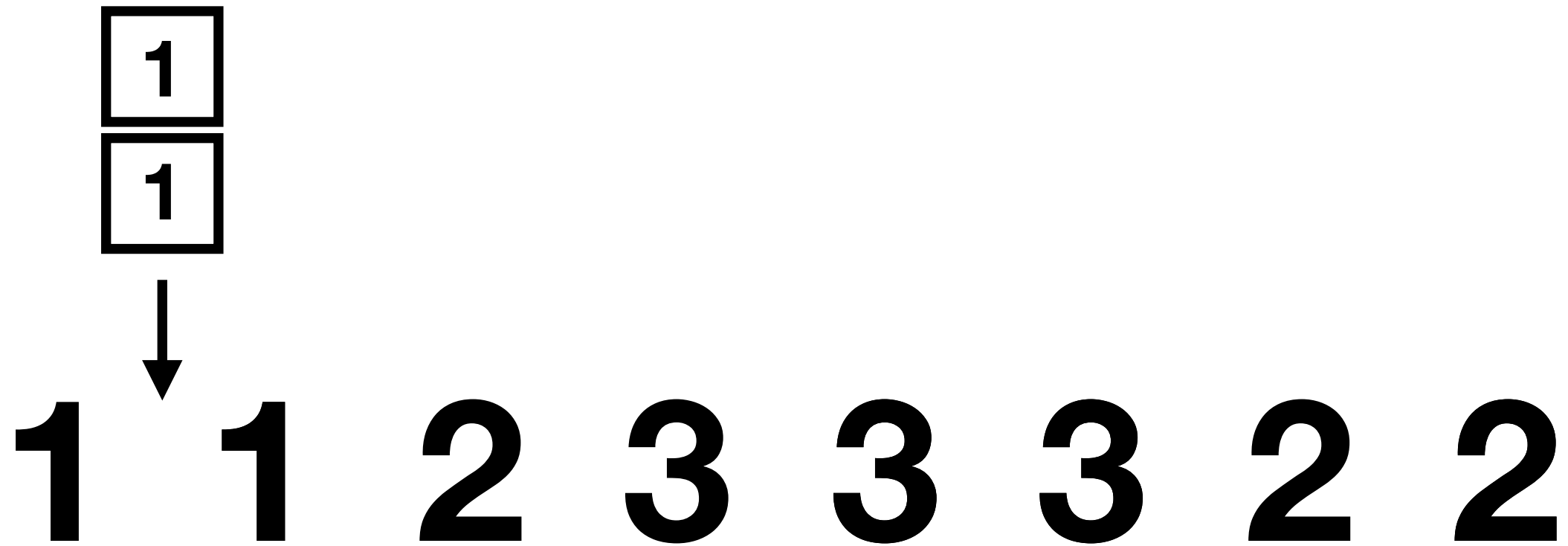
An example

Remove repetitions from the input



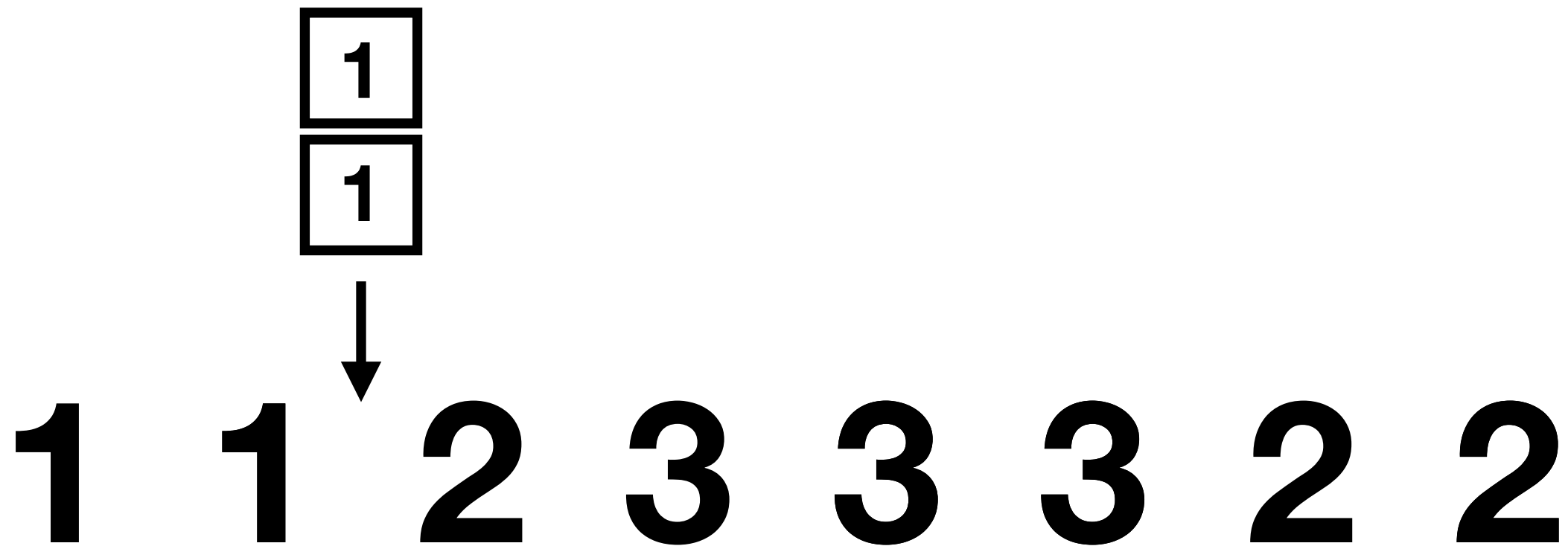
An example

Remove repetitions from the input



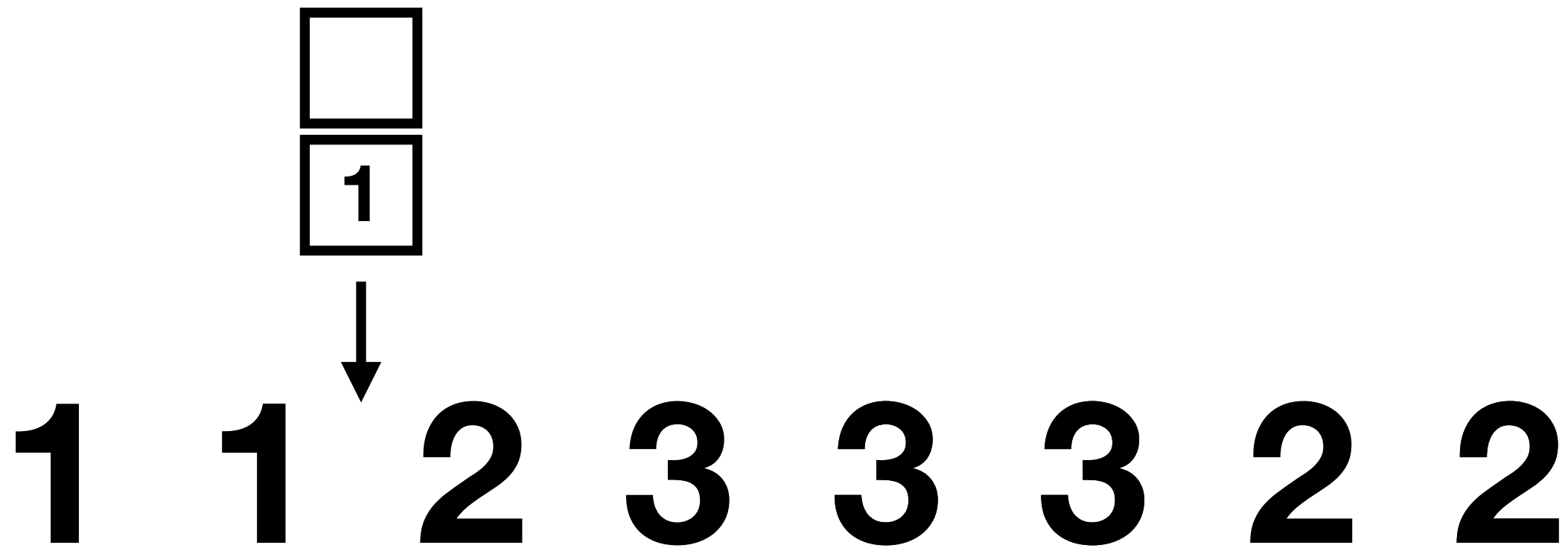
An example

Remove repetitions from the input



An example

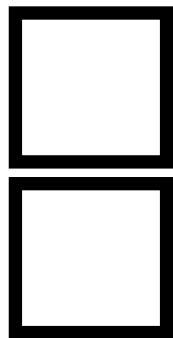
Remove repetitions from the input



An example

Remove repetitions from the input

1

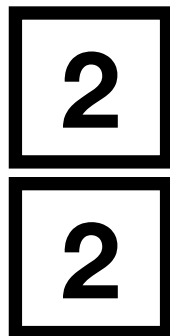


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1

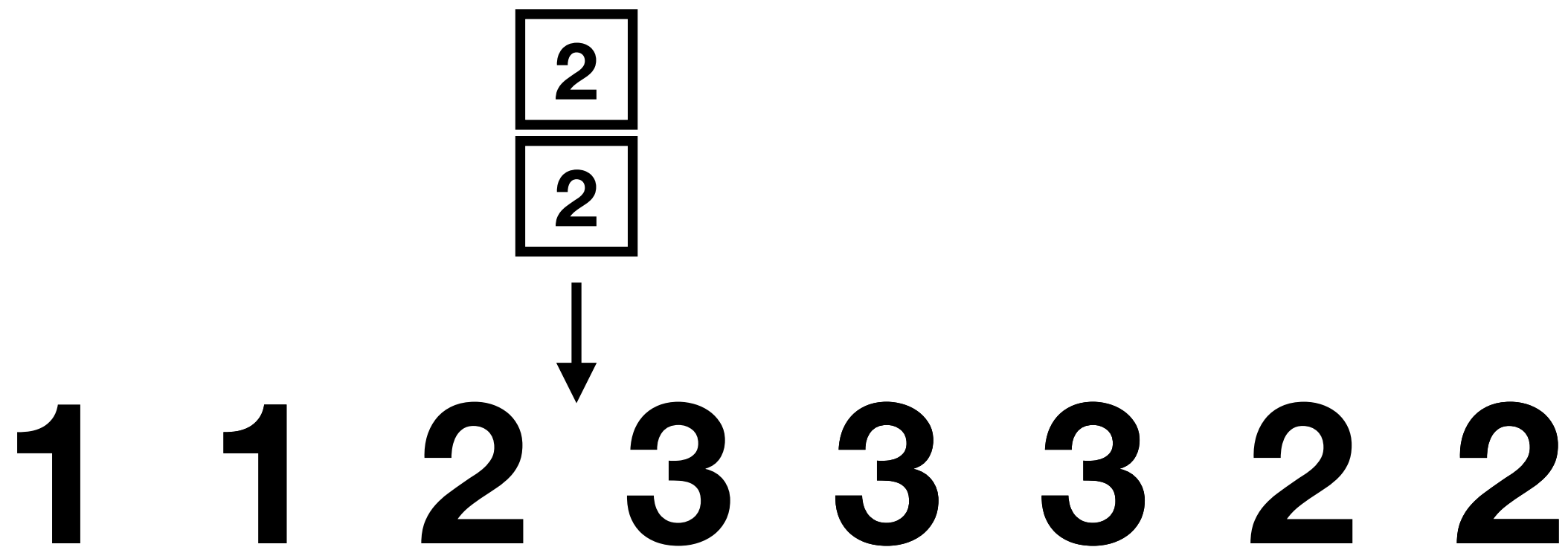


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

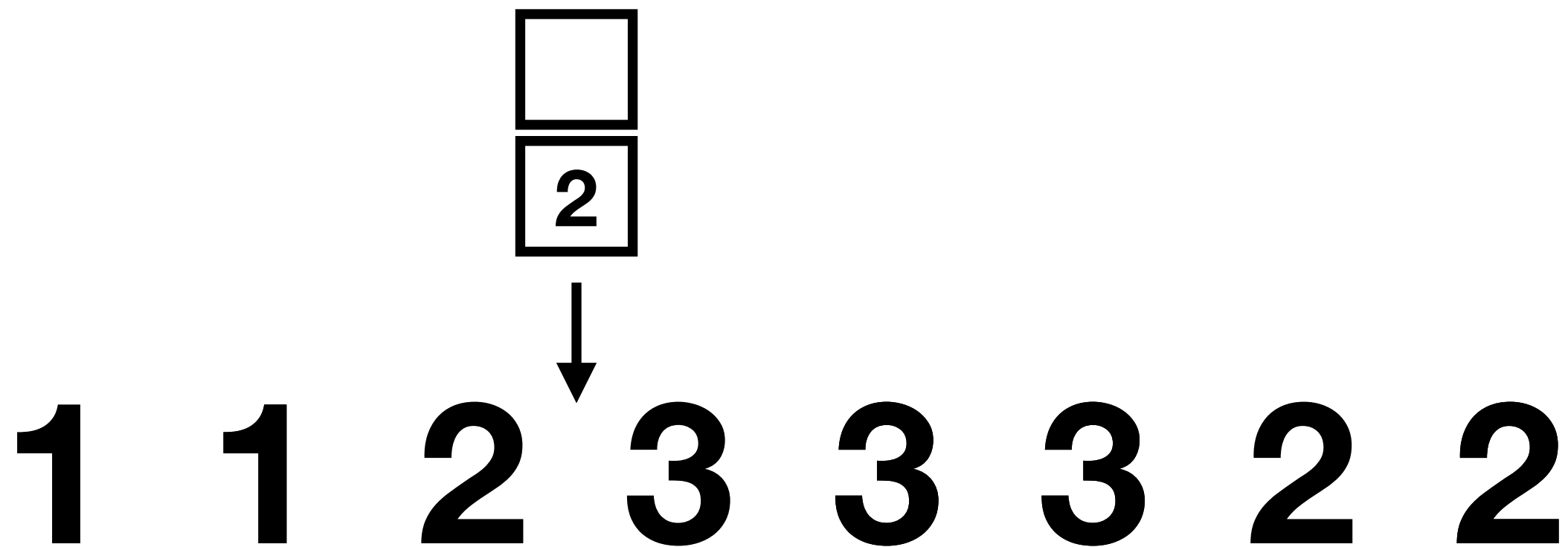
1



An example

Remove repetitions from the input

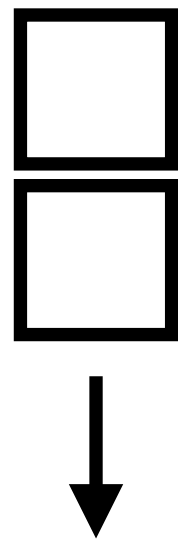
1



An example

Remove repetitions from the input

1 2

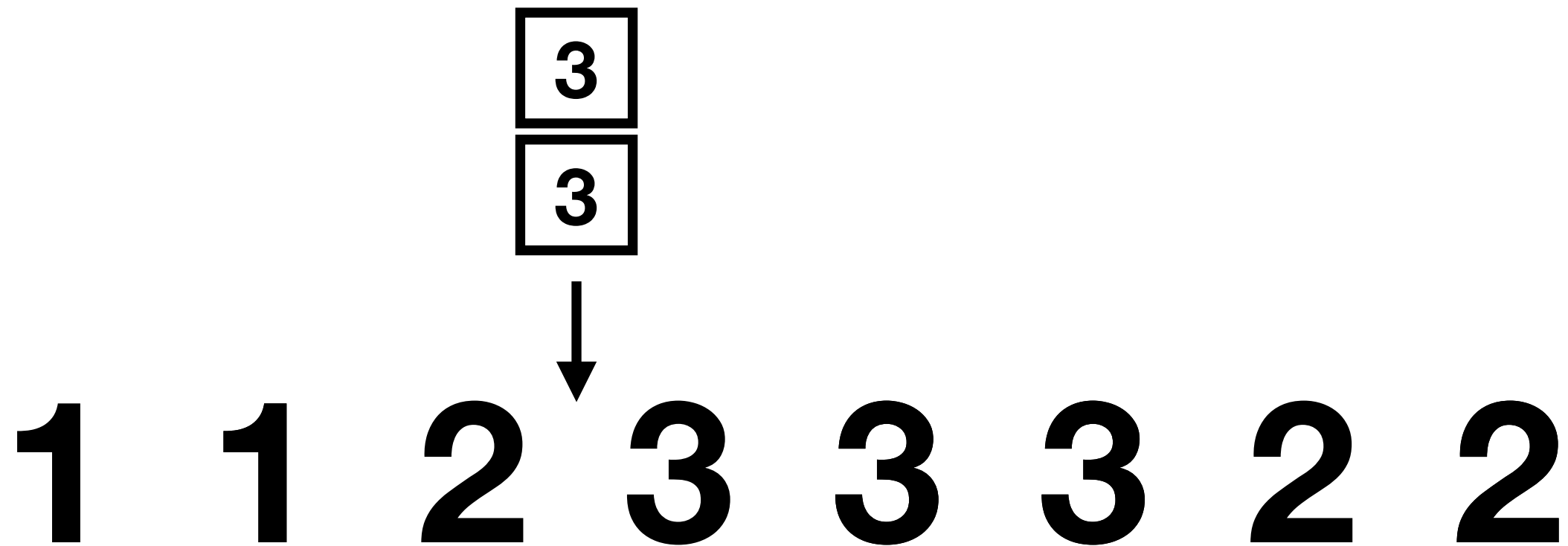


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

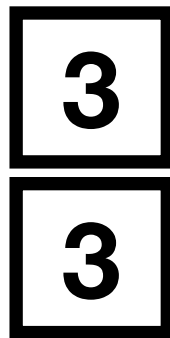
1 2



An example

Remove repetitions from the input

1 2

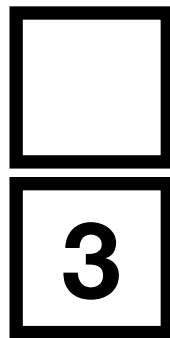


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2

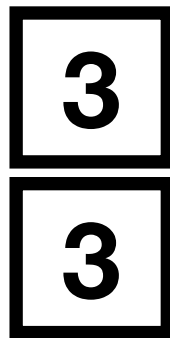


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2



1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2

3

3

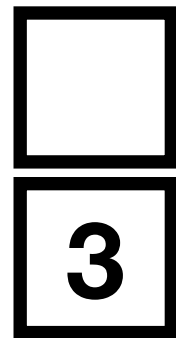


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2



1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2

3

3

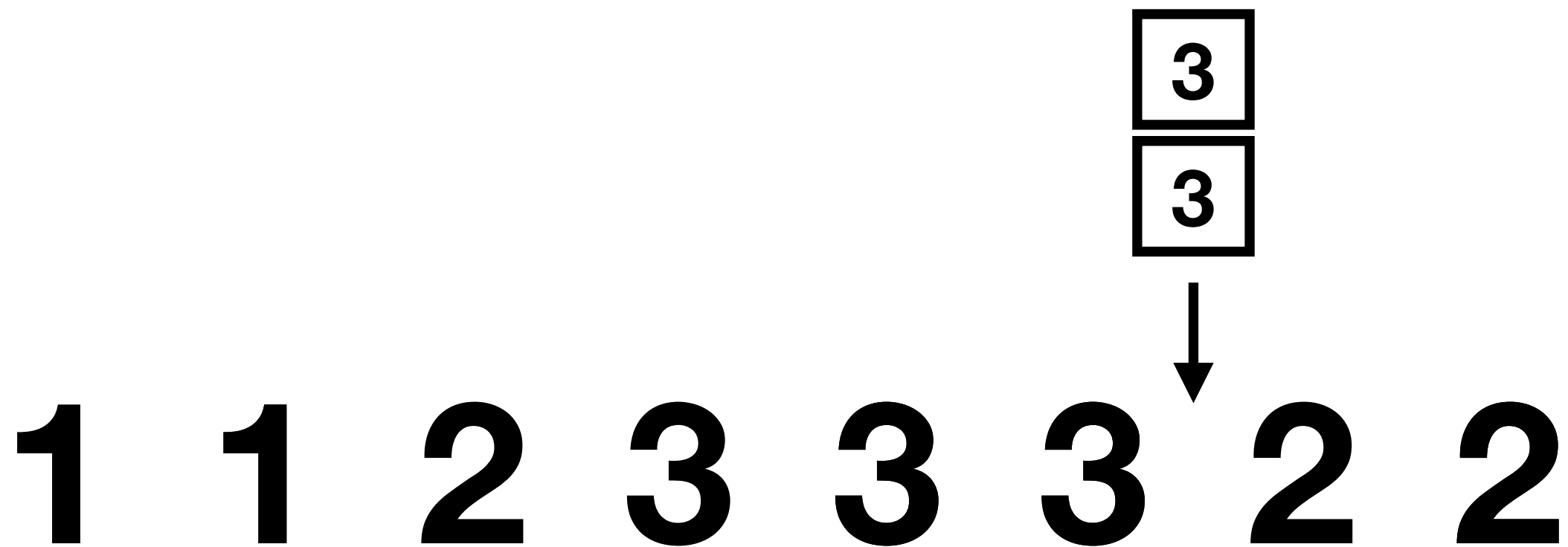


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

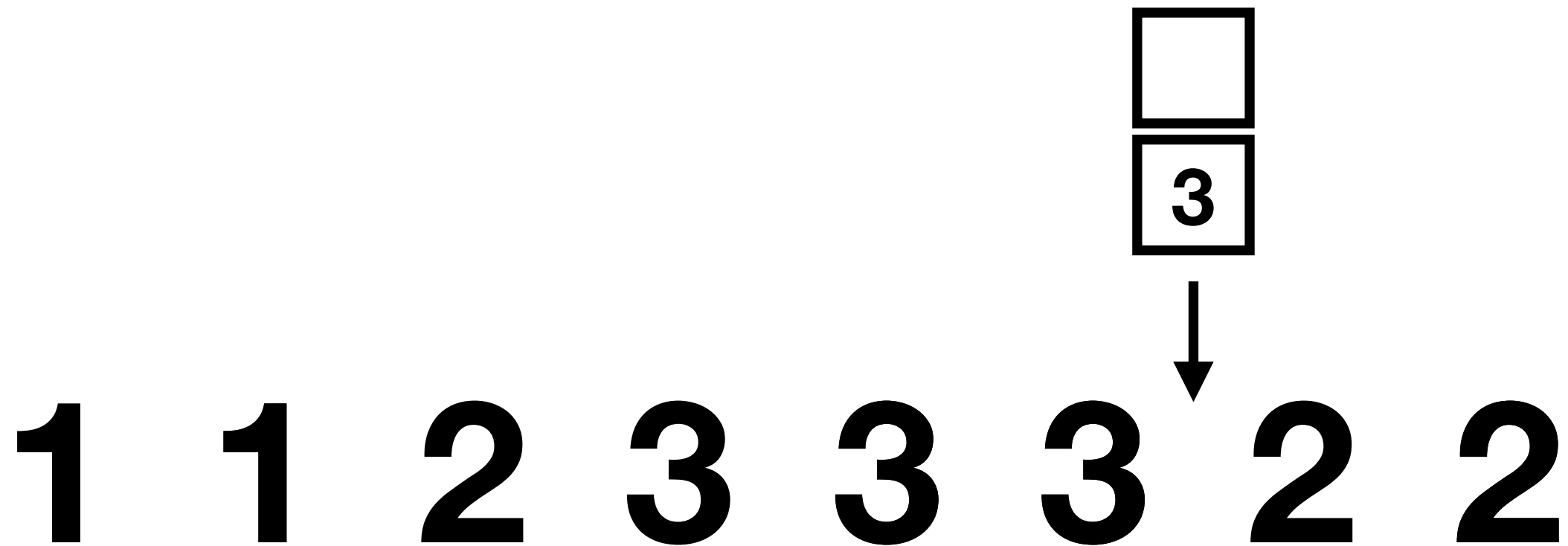
1 2



An example

Remove repetitions from the input

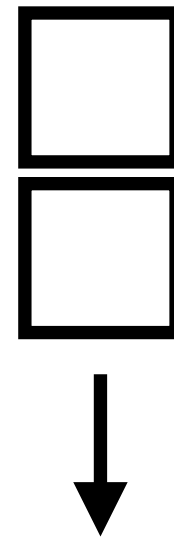
1 2



An example

Remove repetitions from the input

1 2 3

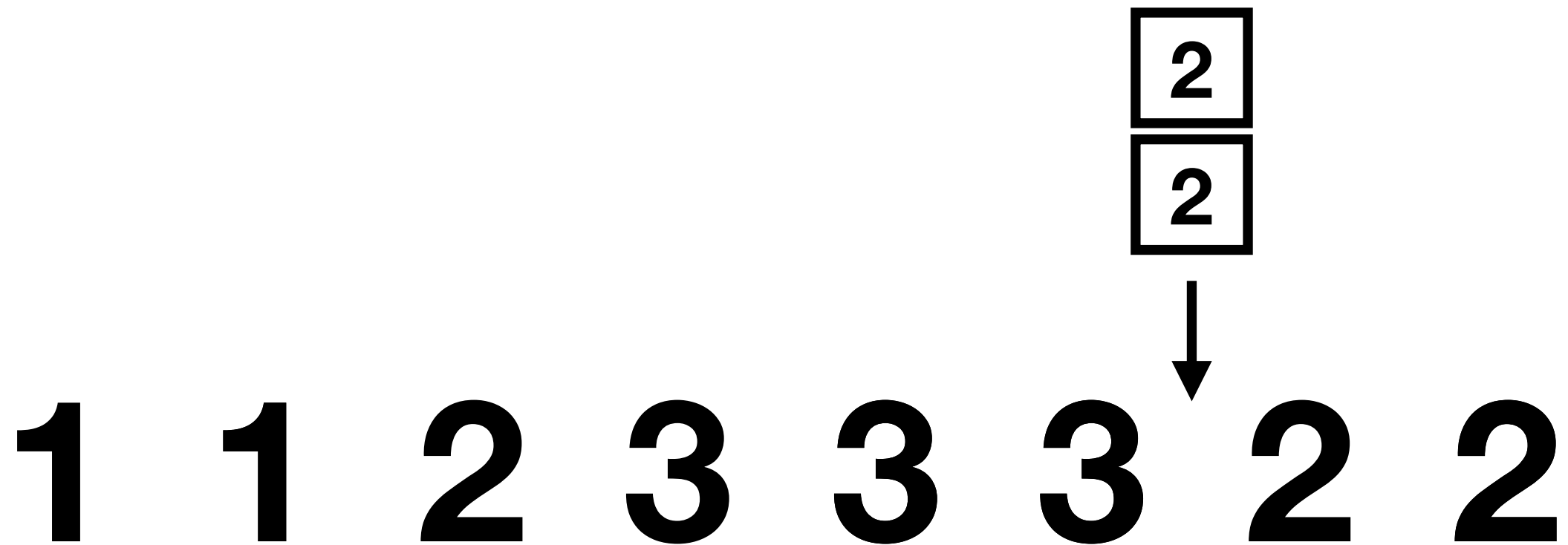


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

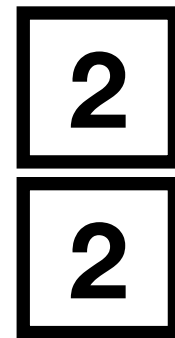
1 2 3



An example

Remove repetitions from the input

1 2 3

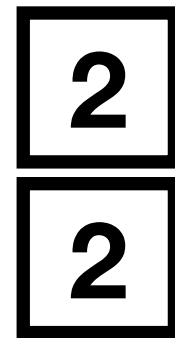


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3

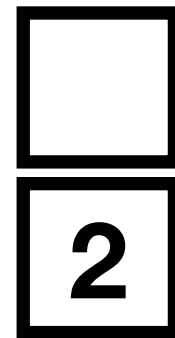


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3

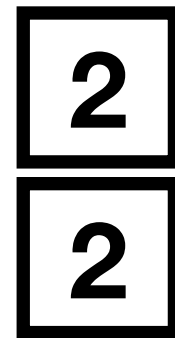


1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3



1 1 2 3 3 3 2 2

An example

Remove repetitions from the input

1 2 3

1 1 2 3 3 3 2 2



An example

Remove repetitions from the input

1 2 3 2

1 1 2 3 3 3 2 2



An example

Remove repetitions from the input

1 2 3 2

1 1 2 3 3 3 2 2

Krohn-Rhodes theorem

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Prime functions

(Equivariant) homomorphism

An equivariant function: $f : \Sigma \rightarrow \Gamma^*$

Extended to: $f^* : \Sigma^* \rightarrow \Gamma^*$

Prime functions

(Equivariant) homomorphism

Double the letters:

1 2 3 2

Prime functions

(Equivariant) homomorphism

1 1 2 2 3 3 2 2

1 2 3 2

Prime functions

(Equivariant) homomorphism

1 1 2 2 3 3 2 2

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Prime functions

Delay $\Sigma^* \rightarrow ((\Sigma + \{+, -\})^2)^*$

Remove repetitions:

1 1 2 3 3 3 2 2

Prime functions

Delay $\Sigma^* \rightarrow ((\Sigma + \{\vdash, \dashv\})^2)^*$

Remove repetitions:

⊢	1	1	2	3	3	3	2	2
1	1	2	3	3	3	2	2	⊣

Prime functions

Delay $\Sigma^* \rightarrow ((\Sigma + \{+, -\})^2)^*$

Remove repetitions:

1 2
2 3

3
2

2
1

Prime functions

Delay $\Sigma^* \rightarrow ((\Sigma + \{+, -\})^2)^*$

Remove repetitions:

1 2 3 2

Prime functions

Delay $\Sigma^* \rightarrow ((\Sigma + \{+, -\})^2)^*$

Remove repetitions:

1 2 3 2

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

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Letter propagation

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Remove letters from odd positions:

1 9 3 9 2 7

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Remove letters from odd positions:

1 1 1 1 1 1 (\mathbb{Z}_2)

1 9 3 9 2 7

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Remove letters from odd positions:

1 0 1 0 1 0

1 9 3 9 2 7

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Remove letters from odd positions:

0	0	0
9	9	7

Prime functions

Finite group on prefixes $(\Sigma \times G)^* \rightarrow (\Sigma \times G)^*$

Remove letters from odd positions:

9 9 7

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

$$F = \{1, a, b\}$$

$$1a = aa = ba = a$$

$$1b = bb = ab = b$$

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

1 2 5 7 7 7 8 8 9

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

┌ 1 2 5 7 7 7 8 8 9
1 2 5 7 7 7 8 8 9 └

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

1 1 1 1 a a 1 a 1 1
┌ 1 2 5 7 7 7 8 8 9
1 2 5 7 7 7 8 8 9 ─

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

1 1 1 1 a a a a a a

┌ 1 2 5 7 7 7 8 8 9

1 2 5 7 7 7 8 8 9 └

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

1 1 1
1 2 5
2 5 7

Prime functions

Flip-flop monoid on prefixes $(\Sigma \times F)^* \rightarrow (\Sigma \times F)^*$

Remove everything after the first repetition:

257

Original Krohn-Rhodes theorem

Every function f recognised by a one-way transducer iff
it is an element of $(prime)^*$

Original Krohn-Rhodes theorem

Every function f recognised by a one-way transducer can also be expressed as an element of $(prime)^*$

$$\text{remove_repetitions} = f^* \circ \text{delay}$$

$$f(\langle x, y \rangle) = \begin{cases} y & \text{if } x \neq y \\ \epsilon & \text{otherwise} \end{cases}$$

Prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Prime functions


Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Operations on 1 register:

 Read value

 Output value

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Operations on 1 register:

\uparrow Read value

\downarrow Output value

Subject to single-use restriction

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Change the last letter to the first letter:

1 2 5 7 7 9

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Change the last letter to the first letter:

┌ 1 2 5 7 7 9
1 2 5 7 7 9 └

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Change the last letter to the first letter:

┌
1
┌ 1 2 5 7 7 9
1 2 5 7 7 9 └

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

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1 2 5 7 7 1

Prime functions

Letter propagation $(\Sigma \times P\{\uparrow, \downarrow\})^* \rightarrow (\Sigma \times (\{.\} + \Sigma))^*$

Change the last letter to the first letter:

1 2 5 7 7 1

Krohn-Rhodes theorem with atoms

Every function f recognised by a single use, one-way iff it is an element of *(prime + letter propagation)**

Two more prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Iterated reverse

Iterated duplicate

Prime functions

Iterated reverse $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

1 2 5 # 7 9 # 9

Prime functions

Iterated reverse $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

5 2 1 # 9 7 # 9
1 2 5 # 7 9 # 9

Prime functions

Iterated reverse $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

5 2 1 # 9 7 # 9

Prime functions

Iterated duplicate $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

1 2 5 # 7 9 # 9

Prime functions

Iterated duplicate $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

1 2 5 1 2 5 # 7 9 7 9 # 9
1 2 5 # 7 9 # 9

Prime functions

Iterated duplicate $(\Sigma + \{\#\})^* \rightarrow (\Sigma + \{\#\})^*$

1 2 5 1 2 5 # 7 9 7 9 # 9

Two more prime functions

(Equivariant) homomorphism

Delay

Finite group on prefixes

Flip-flop on prefixes

Letter propagation

Iterated reverse

Iterated duplicate

Two-way Krohn-Rhodes theorem with atoms

single use, two way register transducer
=
*(two-way prime + letter propagation)**

Two corollaries

**One-way single-use register automata
are closed under compositions**

**Two-way single-use register automata
are closed under compositions**

One more corollary

All of the following recognise the same class of transductions:

(Two-way prime + letter propagation)*

original model by Krohn, Rhodes 1963

Two-way, single-use register transducers

original model by Shepherdson 1959

String streaming, single-use register transducers

original model by Alur, Cerny 2010

Regular list functions

original model by Bojańczyk, Daviaud, Krishna 2018

The general picture

1. (Two-way prime + letter propagation)*
2. Two-way, single-use register transducers
3. String streaming, single-use register transducers
4. Regular list functions

1. (Prime + reversed flip-flop + letter propagation + reversed letter propagation)*

1. (Prime + letter propagation)*
2. One-way, single-use register transducers

The general picture

1. (Two-way prime + letter propagation)*
2. Two-way, single-use register transducers
3. String streaming, single-use register transducers
4. Regular list functions
5. Rigid MSO~ transductions

1. (Prime + reversed flip-flop + letter propagation + reversed letter propagation)*

1. (Prime + letter propagation)*
2. One-way, single-use register transducers