# The pumping lemma is incorrect? 

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August 3, 2011

We will show that Lemma 9.5 in [1] is false. This lemma says that in each long enough run $r$ of any automaton there exists a pumping pair of configurations $u, v$. From the definition of a pumping pair we use only the following:

- $u$ is (strictly) before $v$ in the run,
- $\rho r(u)=\rho r(v)$ (the state in $u$ and in $v$ is the same),
- $\pi r(u) \triangleleft_{1} \pi r(v)$.

Consider an automaton $\mathcal{A}$ of level 3, which realizes the following program:
repeat forever
push $_{2}$
push $_{3}$
pop $_{1}$
push $_{3}$
pop $_{2}$
push $_{3}$

Thus it has 6 states and one loop of transitions between them. The stack alphabet contains only the $a$ symbol. The automaton does not read any input (it has only $\epsilon$-transitions). Take the initial configuration $[[[a a]]]$ (one order 1 stack with two symbols). Started from it, the automaton has exactly one infinite run. Hence from Lemma 9.5 there is a pumping pair $u, v$ in it. We will show that this is not true.

How our automaton works? First observe that it never makes any pop $3_{3}$ operation. Hence only the topmost order 2 stack is accessed. By making a push ${ }_{3}$ operation we keep a history of the current contents of the topmost order 2 stack.

Now observe how the topmost order 2 stack changes. It has tree possible contents, between which we loop:

$$
\begin{aligned}
x & =[[a a]], \\
y & =[[a a][a a]], \\
z & =[[a a][a]] .
\end{aligned}
$$

We have $x \triangleleft_{1} y$ and $z \triangleleft_{1} y$, but $x$ and $z$ are $\triangleleft_{1}$-incomparable.
Assume we have a pumping pair $u, v(u$ is before $v)$. Let $\pi r(u)=\xi_{1} \ldots \xi_{k}$ and $\pi r(v)=\zeta_{1} \ldots \zeta_{l}$. The configurations $u, v$ have the same state, which means that $\xi_{k}=\zeta_{l}$. As $v$ is strictly after $u$, there is $l \geq k+3$. Because $\pi r(u) \triangleleft_{1} \pi r(v)$, it has to be

$$
\xi_{k} \triangleleft_{1} \zeta_{k} \quad \text { and } \quad \xi_{k} \triangleleft_{1} \zeta_{k+1} \quad \text { and } \quad \xi_{k} \triangleleft_{1} \zeta_{k+2}
$$

We know that $\zeta_{1}, \zeta_{2}, \ldots, \zeta_{l-1}=x, y, z, x, y, z, \ldots$ (and $\zeta_{l}$ is either equal to $\zeta_{l-1}$, or is the next symbol). This means that

$$
\xi_{k} \triangleleft_{1} x \quad \text { and } \quad \xi_{k} \triangleleft_{1} y \quad \text { and } \quad \xi_{k} \triangleleft_{1} z
$$

But none of $x, y, z$ satisfies this. Hence there is no pumping pair.

## References

[1] A. Blumensath. On the structure of graphs in the caucal hierarchy. Theor. Comput. Sci., 400(1-3):19-45, 2008.

