

# Parameterized algorithms — tutorial 1

## Branching algorithms

**Problem 1.** Prove that in any graph  $G$  there are at most  $2^k$  inclusion-wise minimal vertex covers of size at most  $k$ , and they can be enumerated in time  $2^k \cdot n^{\mathcal{O}(1)}$ . Is the  $2^k$  bound optimal?

**Problem 2.** Prove that CLIQUE parameterized by  $n - k$  is FPT.

**Problem 3.** Suppose  $\mathcal{F}$  is a finite family of graphs. A graph  $G$  is  $\mathcal{F}$ -free if  $G$  does not contain any graph from  $\mathcal{F}$  as an induced subgraph. In the  $\mathcal{F}$ -FREE VERTEX DELETION problem we are given a graph  $G$  and integer  $k$ , and we ask whether one can remove at most  $k$  vertices from  $G$  to obtain an  $\mathcal{F}$ -free graph. Prove that there is a constant  $c$ , depending only on  $\mathcal{F}$ , such that  $\mathcal{F}$ -FREE VERTEX DELETION can be solved in time  $c^k \cdot n^{\mathcal{O}(1)}$ . Prove the same for the  $\mathcal{F}$ -FREE EDGE DELETION,  $\mathcal{F}$ -FREE EDGE COMPLETION, and  $\mathcal{F}$ -FREE EDGE EDITING problems, where instead of removing vertices we may remove edges, add edges, or add and remove edges.

**Problem 4.** Prove that for every constant  $d \in \mathbb{N}$ , the  $d$ -REGULAR VERTEX DELETION problem, where given  $G$  and  $k$  we want to remove at most  $k$  vertices from  $G$  to obtain a  $d$ -regular graph, is FPT when parameterized by  $k$ .

**Problem 5.** A directed graph  $D$  is a *tournament* if between every pair of vertices there is exactly one arc. In DIRECTED FEEDBACK VERTEX SET we are given a directed graph  $D$  and integer  $k$ , and we want to remove at most  $k$  vertices from  $D$  to obtain a DAG. Prove that DFVS on tournaments is FPT when parameterized by  $k$ .

**Problem 6.** Prove that for every constant  $p \in \mathbb{N}$ , the following problem can be solved in time  $p^k \cdot \|\varphi\|^{\mathcal{O}(1)}$ : given a boolean formula  $\varphi$  in  $p$ -CNF, decide whether there exists an assignment that satisfies  $\varphi$  and sets at most  $k$  variables to true.

**Problem 7.** Prove that the following problem can be solved in time  $2^k \cdot \|\varphi\|^{\mathcal{O}(1)}$ : given a boolean formula  $\varphi$  in CNF, decide whether there exists an assignment that satisfies at most  $k$  clauses of  $\varphi$ .

**Problem 8.** We consider the INDEPENDENT SET problem: given  $G$  and  $k$ , decide whether there is a subset of  $k$  pairwise non-adjacent vertices in  $G$ . Prove that this problem can be solved on graphs of maximum degree 4 in time  $2.31^k \cdot n^{\mathcal{O}(1)}$ .

**Problem 9.** Prove that TRIANGLE-FREE VERTEX DELETION can be solved in time  $2.562^k \cdot n^{\mathcal{O}(1)}$ .

**Problem 10.** In the CLOSEST STRING problem we are given strings  $s_1, \dots, s_n$  over some alphabet  $\Sigma$ , each of length  $L$ , and a number  $d$ . The question is whether there exists a string  $t \in \Sigma^L$  that is at Hamming distance (i.e. number of differing symbols) at most  $d$  from each of the strings  $s_1, \dots, s_n$ . Prove that this problem can be solved in time  $(2d)^d \cdot (nL)^{\mathcal{O}(1)}$ .