

Equilibrium selection in evolutionary games with random matching of players

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Abstract

We discuss stochastic dynamics of populations of individuals playing games. Our models possess two evolutionarily stable strategies: an efficient one, where a population is in a state with the maximal payoff (fitness) and a risk-dominant one, where players are averse to risks. We assume that individuals play with randomly chosen opponents (they do not play against average strategies as in the standard replicator dynamics). We show that the long-run behavior of a population depends on its size and the mutation level. © 2004 Elsevier Ltd. All rights reserved.

Keywords: Population dynamics; Evolutionarily stable strategy; Equilibrium selection; Stochastic stability

1. Introduction

The long-run behavior of interacting individuals can often be described within game-theoretic models. The basic notion here is that of a Nash equilibrium. This is a state of population—an assignment of strategies to players—such that no player, for fixed strategies of his opponents, has an incentive to deviate from his current strategy; the change can only diminish his payoff. Nash equilibrium is supposed to be a result of decisions of rational players. Maynard Smith (1974, 1982) has refined this concept of equilibrium to include the stability of Nash equilibria against mutants. He introduced the fundamental notion of an evolutionarily stable strategy. If everybody plays such a strategy, then the small number of mutants playing a different strategy is eliminated from the population. The dynamical interpretation of the evolutionarily stable strategy was later provided by several authors (Taylor and Jonker, 1978; Hofbauer et al., 1979; Zeeman, 1981). They proposed a system of differential

or difference equations, the so-called replicator equations, which describe the time evolution of frequencies of strategies. It is known that any evolutionarily stable strategy is an asymptotically stable stationary point of such dynamics (Hofbauer and Sigmund, 1988; Weibull, 1997).

Here we will discuss a stochastic adaptation dynamics of a population of players interacting in discrete moments of time. We will analyse two-player games with two strategies and two evolutionarily stable strategies. The efficient strategy (also called payoff dominant) when played by the whole population results in its highest possible payoff (fitness). The risk-dominant one is played by individuals averse to risks. The strategy is risk dominant if it has a higher expected payoff against a player playing both strategies with equal probabilities. We will address the problem of equilibrium selection—a strategy which will be played in the long run with a high frequency.

We will review two models of adaptive dynamics of a population of a fixed number of individuals. In both of them, the selection part of the dynamics ensures that if the mean payoff of a given strategy at the time t is bigger than the mean payoff of the other one, then the number

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of individuals playing the given strategy should increase in $t + 1$. In the first model, introduced by [Kandori et al. \(1993\)](#), one assumes (as in the standard replicator dynamics) that individuals receive average payoffs with respect to all possible opponents—they play against the average strategy. In the second model, introduced by [Robson and Vega-Redondo \(1996\)](#), at any moment of time, individuals play only one game with randomly chosen opponents. In both models, players may mutate with a small probability; hence the population may move against a selection pressure. To describe the long-run behavior of such stochastic dynamics, [Foster and Young \(1990\)](#) introduced a concept of stochastic stability. A configuration of a system is stochastically stable if it has a positive probability in the stationary state of the above dynamics in the limit of no mutations. It means that in the long run we observe it with a positive frequency. [Kandori et al. \(1993\)](#) showed that in their model, the risk-dominant strategy is stochastically stable—if the mutation level is small enough we observe it in the long run with the frequency close to one. In the model of [Robson and Vega-Redondo \(1996\)](#), the efficient strategy is stochastically stable. It is one of the very few models in which an efficient strategy is stochastically stable in the presence of a risk-dominant one. The population evolves in the long run to a state with the maximal fitness.

The main goal of our paper is to investigate the effect of the number of players on the long-run behavior of the Robson–Vega-Redondo model. We will discuss parallel and sequential dynamics, and the one, where each individual enjoys each period a revision opportunity with some independent probability. We will show that in the last two dynamics, for any arbitrarily low but a fixed level of mutations, if the number of players is sufficiently big, a risk-dominant strategy is played in the long run with a frequency close to one—a stochastically stable efficient strategy is observed with a very low frequency. It means that when the number of players increases, the population undergoes a transition between an efficient payoff-dominant equilibrium and a risk-dominant one. We will also show that for some range of payoff parameters, stochastic stability itself depends on the number of players. If the number of players is below a certain value (which may be arbitrarily large), then a risk-dominant strategy is stochastically stable. An efficient strategy becomes stochastically stable only if n is large enough, as proved by [Robson and Vega-Redondo \(1996\)](#).

In Section 2, we introduce Kandori–Mailath–Rob and Robson–Vega-Redondo models and review their main properties. In Section 3, we analyse the Robson–Vega-Redondo model in the limit of the infinite number of players and show our main results. Discussion follows in Section 4.

2. Models of adaptive dynamics with mutations

We will consider a finite population of n individuals who have at their disposal one of the two strategies: A and B . At every discrete moment of time, $t = 1, 2, \dots$, they are randomly paired (we assume that n is even) to play a two-player symmetric game with payoffs given by the following matrix:

$$U = \begin{matrix} & A & B \\ A & a & b \\ B & c & d \end{matrix}$$

where the ij entry, $i, j = A, B$, is the payoff of the first (row) player when he plays the strategy i and the second (column) player plays the strategy j . We assume that both players are the same and hence payoffs of the column player are given by the matrix transposed to U ; such games are called symmetric.

We assume that $a > c$ and $d > b$, therefore, both A and B are evolutionarily stable strategies, and $a + b < c + d$, so the strategy B has a higher expected payoff against a player playing both strategies with the probability $\frac{1}{2}$. We say that B risk dominates the strategy A (the notion of the risk-dominance was introduced and thoroughly studied by [Harsányi and Selten \(1988\)](#)). We also assume that $a > d$; hence we have a selection problem of choosing between the risk-dominant B and the so-called payoff-dominant or efficient strategy A .

At every discrete moment of time t , the state of our population is described by the number of individuals, z_t , playing A . Formally, by the state space we mean the set $\Omega = \{z, 0 \leq z \leq n\}$.

Now we will describe the dynamics of our system. It consists of two components: selection and mutation. The selection mechanism ensures that if the mean payoff of a given strategy, $\pi_i(z_t)$, $i = A, B$, at the time t is bigger than the mean payoff of the other one, then the number of individuals playing the given strategy should increase in $t + 1$. In their paper, [Kandori et al. \(1993\)](#) write

$$\begin{aligned} \pi_A(z_t) &= \frac{a(z_t - 1) + b(n - z_t)}{n - 1}, \\ \pi_B(z_t) &= \frac{cz_t + d(n - z_t - 1)}{n - 1}, \end{aligned} \tag{2.1}$$

provided $0 < z_t < n$.

It means that in every time step, players are paired infinitely many times to play the game or equivalently, each player plays with every other player and his payoff is the sum of corresponding payoffs. This model may be therefore considered as an analog of replicator dynamics for populations with fixed numbers of players.

The selection dynamics is formalized in the following way:

$$\begin{aligned} z_{t+1} &> z_t && \text{if } \pi_A(z_t) > \pi_B(z_t), \\ z_{t+1} &< z_t && \text{if } \pi_A(z_t) < \pi_B(z_t), \\ z_{t+1} &= z_t && \text{if } \pi_A(z_t) = \pi_B(z_t), \\ z_{t+1} &= z_t && \text{if } z_t = 0 \text{ or } z_t = n. \end{aligned} \tag{2.2}$$

Now mutations are added. Players may switch to new strategies with the probability ε . It is easy to see that for any two states of the population there is a positive probability of the transition between them in some finite number of time steps. We have therefore obtained an irreducible Markov chain with $n + 1$ states. It has a unique stationary probability distribution (a stationary state) which we denote by μ_n^ε . It was shown (Kandori et al., 1993) that $\lim_{\varepsilon \rightarrow 0} \mu_n^\varepsilon(0) = 1$ which means that in the long run, in the limit of no mutations, all players play the risk-dominant strategy B . We say that the risk-dominant strategy is *stochastically stable*.

The general setup in the Robson–Vega-Redondo model (1996) is the same. However, individuals are paired only once at every time step and play only one game before the selection process takes place. Let p_t denote the random variable which describes the number of cross-pairings, i.e. the number of pairs of matched individuals playing different strategies at the time t . Let us notice that p_t depends on z_t . For a given realization of p_t and z_t , mean payoffs obtained by each strategy are as follows:

$$\begin{aligned} \tilde{\pi}_A(z_t, p_t) &= \frac{a(z_t - p_t) + bp_t}{z_t}, \\ \tilde{\pi}_B(z_t, p_t) &= \frac{cp_t + d(n - z_t - p_t)}{n - z_t}, \end{aligned} \tag{2.3}$$

provided $0 < z_t < n$. Then the authors show that the payoff-dominant strategy is stochastically stable. We will outline their proof.

First of all, one can show that there exists k such that if n is large enough and $z_t \geq k$, then there is a positive probability (a certain realization of p_t) that after a finite number of steps of the mutation-free selection dynamics, all players will play A . Likewise, if $z_t < k$ (for any $k \geq 1$), then if the number of players is large enough, then after a finite number of steps of the mutation-free selection dynamics all players will play B . In other words, $z = 0$ and $z = n$ are the only absorbing states of the mutation-free dynamics. Moreover, if n is large enough, then if $z_t \geq n - k$, then the mean payoff obtained by A is always (for any realization of p_t) bigger than the mean payoff obtained by B (in the worst case all B -players play with A -players). Therefore, the size of the basin of attraction of the state $z = 0$ is at most $n - k - 1$ and that of $z = n$

is at least $n - k$. Observe that mutation-free dynamics is not deterministic (p_t describes the random matching) and therefore basins of attraction may overlap. It follows that the system needs at least $k + 1$ mutations to evolve from $z = n$ to 0 and at most k mutations to evolve from $z = 0$ to n . Now using a tree representation of stationary states of irreducible Markov chains (Freidlin and Wentzell, 1970, 1984; see also Appendix B), Robson and Vega-Redondo finish the proof and show that the efficient strategy is stochastically stable.

However, as outlined above, their proof requires the number of players to be sufficiently large. We will now show that for some payoff parameters, a risk-dominant strategy is stochastically stable if the number of players is below a certain value which can be arbitrarily big.

Let $c > b$. Now if the population consists of only one B -player and $n - 1$ A -players and if $c > [a(n - 2) + b]/(n - 1)$, that is, $n < (2a - c - b)/(a - c)$, then $\tilde{\pi}_B > \tilde{\pi}_A$. It means that one needs only one mutation to evolve from $z = n$ to 0 . It is easy to see that two mutations are necessary to evolve from $z = 0$ to n . Using again the tree representation of stationary states one can prove the following theorem.

Theorem 1. *If $n < \frac{2a-c-b}{a-c}$, then the risk-dominant strategy B is stochastically stable in the case of random matching of players.*

To see stochastically stable states, we need to take the limit of no mutations. We will now examine the long-run behavior of the Robson–Vega-Redondo model for a fixed level of mutations in the limit of the infinite number of players.

3. Long-run behavior in the limit of infinitely many players

We will consider three specific cases of the selection rule (2.2).

In the parallel dynamics, everyone in the selection process chooses at the same time (all players are synchronized) a strategy with the bigger average payoff. It means that after mutations have taken place, the selection moves the population to one of the two extreme states, $z = 0$ or n . Our system becomes then a two-state Markov chain with a unique stationary state μ_n^ε (a similar model was discussed in Vega-Redondo (1996)). We will show that for any number of players, if the mutation level is sufficiently small, then in the long run almost all individuals play the payoff-dominant strategy. The same result holds for any small mutation level if the number of players is large enough.

Theorem 2. *In parallel dynamics,*

$$\lim_{\varepsilon \rightarrow 0} \mu_n^\varepsilon(n) = 1 \quad \text{for every } n,$$

$$\lim_{n \rightarrow \infty} \mu_n^\varepsilon(n) = 1 \quad \text{for every } \varepsilon < 1.$$

Proof. We are looking for a unique stationary state μ_n^ε of a two-state Markov chain. Let us denote by p_{0n} a transition probability from the state $z = 0$ to n and by p_{n0} from $z = n$ to 0 . We have

$$\mu_n^\varepsilon(n) = \frac{p_{0n}}{p_{0n} + p_{n0}}. \tag{3.1}$$

For the transition from $z = 0$ to n it is enough that two players mutate from B to A and then they are paired to play a game. It follows that

$$p_{0n} > \varepsilon^2 \frac{1}{n-1}. \tag{3.2}$$

Transition from $z = n$ to 0 requires at least γn mutations (for some γ) which means that

$$p_{n0} < \varepsilon^{\gamma n}. \tag{3.3}$$

It follows from (3.1–3.3) that

$$\mu_n^\varepsilon(n) > \frac{1}{1 + (n-1)\varepsilon^{\gamma n-2}}. \tag{3.4}$$

Hence, $\mu_n^\varepsilon(n)$ is arbitrarily close to one if ε is sufficiently small or n is sufficiently big. \square

Now, we will analyse the other extreme case of a selection rule (2.2)—a sequential dynamics, where in one time unit only one player can change his strategy. Although our dynamics is discrete in time, it captures the essential features of continuous-time models, where every player has an exponentially distributed waiting time to a moment of a revision opportunity. The probability that two or more players revise their strategies at the same time is therefore equal to zero—this is an example of a birth and death process.

The number of A -players in the population may increase by one in $t + 1$, if a B -player is chosen in t which happens with the probability $(n - z_t)/n$. Analogously, the number of B -players in the population may increase by one in $t + 1$, if an A -player is chosen in t which happens with the probability $(z_t)/n$.

The player who has a revision opportunity chooses in $t + 1$ with the probability $1 - \varepsilon$ the strategy with a higher average payoff in t and the other one with the probability ε .

Let $r(k) = P(\tilde{\pi}_A(z_t, p_t) > \tilde{\pi}_B(z_t, p_t))$ and $l(k) = P(\tilde{\pi}_A(z_t, p_t) < \tilde{\pi}_B(z_t, p_t))$. The sequential dynamics is described by the following transition probabilities:

if $z_t = 0$, then $z_{t+1} = 1$ with the probability ε and $z_{t+1} = 0$ with the probability $1 - \varepsilon$,

if $z_t = n$, then $z_{t+1} = n - 1$ with the probability ε and $z_{t+1} = n$ with the probability $1 - \varepsilon$,

if $z_t \neq 0, n$, then $z_{t+1} = z_t + 1$ with the probability

$$r(k) \frac{n - z_t}{n} (1 - \varepsilon) + (1 - r(k)) \frac{n - z_t}{n} \varepsilon$$

and $z_{t+1} = z_t - 1$ with the probability

$$l(k) \frac{z_t}{n} (1 - \varepsilon) + (1 - l(k)) \frac{z_t}{n} \varepsilon.$$

In the dynamics intermediate between the parallel and sequential one, at time period, each individual has a revision opportunity with some probability $\tau < 1$. Each chosen player follows independently the same rule as in the sequential dynamics. The probability that in one period, a given player will have a revision opportunity should be proportional to the length of the period (which we normalized to 1 in our models). For a fixed ε and an arbitrarily large but fixed n , we consider the limit of continuous time, $\tau \rightarrow 0$, and show that the limiting behavior is already obtained for a sufficiently small τ , namely $\tau < \varepsilon/n^3$.

For an interesting discussion on the importance of the order of taking different limits ($\tau \rightarrow 0$, $n \rightarrow \infty$, and $\varepsilon \rightarrow 0$) in evolutionary models (especially in the aspiration and imitation model), see Samuelson (1997).

In the intermediate dynamics, instead of $(n - z_t)/n$ and z_t/n probabilities we have more involved combinatorial factors. In order to get rid of these inconvenient factors, we will enlarge the state space of the population. The state space Ω' is the set of all configurations of players, that is, all possible assignments of strategies to individual players. Therefore, a state $z_t = k$ in Ω consists of $\binom{n}{k}$ states in Ω' . The sequential dynamics is no longer a birth and death process on Ω' . However, we will be able to treat both dynamics in the same framework.

We will show that for any arbitrarily low but fixed level of mutation, if the number of players is large enough, then in the long run only a small fraction of the population play the payoff-dominant strategy. The smaller the mutation level is, the fewer players use the payoff-dominant strategy.

The following two theorems are proved in Appendix C.

Theorem 3. *In the sequential dynamics, for any $\delta > 0$ and $\beta > 0$ there exist $\varepsilon(\delta, \beta)$ and $n(\varepsilon)$ such that for any $n > n(\varepsilon)$*

$$\mu_n^\varepsilon(z \leq \beta n) > 1 - \delta.$$

Theorem 4. *In the intermediate dynamics, for any $\delta > 0$ and $\beta > 0$ there exist $\varepsilon(\delta, \beta)$ and $n(\varepsilon)$ such that for any $n > n(\varepsilon)$ and $\tau < \frac{\varepsilon}{n^3}$*

$$\mu_n^\varepsilon(z \leq \beta n) > 1 - \delta.$$

Let us note that the above theorems concern an ensemble of configurations, not an individual one. In the limit of the infinite number of players, that is the infinite number of configurations, every single configuration has zero probability in the stationary state. It is an ensemble of configurations that might be stable (Miękisz, 2004a; Miękisz, 2004b).

Let us now assume that at every time period, players are matched many times. It follows from the results in (Kandori et al., 1993; Robson and Vega-Redondo, 1996) analysed in (Vega-Redondo, 1996) that the limits of zero mutations and the infinite number of matching rounds per period do not commute. In the limit of the infinite number of matching rounds per period, individuals play against the average strategy and we obtain the Kandori–Mailath–Rob model and their conclusion follows. On the other hand, for any fixed number of matching rounds (the Robson–Vega-Redondo model), the limit of zero mutations gives us the stochastic stability of an efficient strategy. Here, we investigated the effect of the number of players on the long-run behavior in the random matching model. We showed that the limit of the infinite number of players has the same effect as the limit of the infinite number of matching rounds. In fact, the probability that the average payoff of strategy *A* is bigger than the average payoff of strategy *B* converges in both limits to 1 or 0, if the fraction of the population playing *A* is, respectively, right to or left to the unique mixed Nash equilibrium. Both limits are therefore alternative ways of representing the idea of a low matching-induced noise.

4. Conclusion

We studied the effect of the number of players on the long-run behavior in adaptive dynamics with mutations and random matching of players. We showed that in the sequential dynamics for any arbitrarily low but fixed level of mutation, if the number of players is large enough, then in the long run almost all of them play a risk-dominant strategy. The same result holds if at any period, each individual has a revision opportunity with some small probability. This is in contrast with the result of Robson and Vega-Redondo (1996) who for a fixed number of players take the limit of zero mutations and obtain stochastic stability of a payoff-dominant strategy. It means that when the number of players increases, the population undergoes a transition between an efficient payoff-dominant equilibrium and a risk-dominant one. Therefore, in any specific model, to describe its long-run behavior, one has to evaluate the number of players and the mutation level.

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Appendix A. Random variable of cross-pairings

We will first investigate the random variable p_t which describes the number of cross-pairings in a state z_t . Let $z_t = \alpha n$. Let P be the probability mass function of the random variable of cross-pairings p ; we skip the subscript t .

Proposition A.1.

$$P(|p - n\alpha(1 - \alpha)| > \beta n) \rightarrow_{n \rightarrow \infty} 0 \tag{A.1}$$

for any $\beta > 0$.

Proof. Let the number of *A*-players be equal to $k = \alpha n$. We begin by dividing all players into two groups. We arrange them randomly in a row and pick the first $n/2$ ones to be members of the first group. This will be players who will choose randomly their opponents. Let X denote the random variable counting the number of *A*-players in this group, $X = X_1 + \dots + X_{n/2}$, where $X_i = 1$ if the i th player plays *A*; otherwise $X_i = 0$. The expected value and the variance of X_i are equal to $E(X_i) = \alpha$ and $Var(X_i) = \alpha(1 - \alpha)$, respectively. One can then have that

$$E(X) = E(X_i) + \dots + E(X_{n/2}) = \alpha n/2, \tag{A.2}$$

$$\begin{aligned} Var(X) &= Var(X_i) + \dots + Var(X_{n/2}) \\ &+ 2 \sum_{j < k} (E(X_j X_k) - E(X_j)E(X_k)) \\ &= \frac{n}{2} \alpha(1 - \alpha) \left(1 - \frac{n-2}{2(n-1)} \right). \end{aligned} \tag{A.3}$$

From the Czebyshev inequality we get that

$$P(|X - E(X)| > \beta_1 n) \leq \frac{Var(X)}{(\beta_1 n)^2} \rightarrow_{n \rightarrow \infty} 0 \tag{A.4}$$

for every $\beta_1 > 0$.

Now every player from the first group is randomly paired with a player from the second group. Let us first assume (for pedagogical reasons) that the number of *A*-players in the first group (and therefore in the second group) is exactly equal to $\alpha n/2$.

Let Y be the random variable describing the number of cross-pairings for a given realization of X . $Y = Y_i + \dots + Y_{n/2}$, where $Y_i = 1$ if the i th player has chosen the opponent with a different strategy; $Y_i = 0$ otherwise. The expected value of Y_i is equal to $E(Y_i) = 1 - \alpha$ if $X_i = 1$ and $E(Y_i) = \alpha$ if $X_i = 0$, and $Var(Y_i) = \alpha(1 - \alpha)$.

We get

$$E(Y) = E(Y_i) + \dots + E(Y_{n/2}) = \alpha(1 - \alpha)n \tag{A.5}$$

and using the formula for the variance of the sum of the random variables in (A.3) we obtain

$$\begin{aligned} \text{Var}(Y) &= \frac{n}{2}\alpha(1 - \alpha) + \frac{n\alpha(1 - \alpha)}{n - 2} \\ &+ \frac{n^2\alpha(1 - \alpha) - (3\alpha^2 - 3\alpha + 1)}{2n - 4}. \end{aligned} \tag{A.6}$$

Now let the number of *A*-players in the first group be equal to $(\alpha + \alpha_1)n/2$ for some α_1 . We get that

$$E(Y) = n\alpha(1 - \alpha) + n\alpha_1^2. \tag{A.7}$$

We again use the formula in (A.3) and get that

$$\text{Var}(Y) = C(\alpha, \alpha_1)O(n), \tag{A.8}$$

where $C(\alpha, \alpha_1)$ is some constant depending on α and α_1 and $\lim_{n \rightarrow \infty} O(n)/n = 1$. For any fixed number of *A*-players in the first group we use the Czebyshev inequality to get

$$P(|Y - E(Y)| > \beta_2 n) \rightarrow_{n \rightarrow \infty} 0. \tag{A.9}$$

for every $\beta_2 > 0$.

Now we set $\beta_1 = \alpha_1^2$ in (A.4). Then Proposition A.1 follows with $\beta = \beta_1 + \beta_2$. \square

Now for any state of the system, $z = k, k \neq 0, n$, we will calculate, in the limit of the infinite number of players, the probability, $r(k)$, that the average payoff of *A* is bigger than that of *B*. We have

$$r(k) = P\left(\frac{a(k - p) + bp}{k} > \frac{cp + d(n - k - p)}{n - k}\right). \tag{A.10}$$

Let $k = \alpha n$. It follows from (A.10) that

$$r(\alpha n) = P\left(p\left(\frac{d - c}{n(1 - \alpha)} + \frac{b - a}{\alpha n}\right) > d - a\right). \tag{A.11}$$

If $(d - c)/(1 - \alpha) + (b - a)/\alpha \geq 0$, then $r(\alpha n) = 1$ because $d < a$. This happens for $\alpha \geq (a - b)/(a - c + d - b) \equiv \gamma_1 > \frac{1}{2}$. Let us notice that if $c \leq d$, then $\gamma_1 \leq 1$, if $c > d$, then $\gamma_1 > 1$. For $\alpha < \gamma_1$, from (A.11) we get

$$r(\alpha n) = P\left(p < \frac{n(d - a)(1 - \alpha)\alpha}{(d - c)\alpha + (b - a)(1 - \alpha)}\right). \tag{A.12}$$

Now it follows from Proposition A.1 that if

$$\frac{(d - a)(1 - \alpha)\alpha}{(d - c)\alpha + (b - a)(1 - \alpha)} < \alpha(1 - \alpha), \tag{A.13}$$

which holds for $\alpha < \gamma_2 \equiv (d - b)/(d - b + a - c)$, then

$$\lim_{n \rightarrow \infty} r(\alpha n) = 0. \tag{A.14}$$

Note that $\gamma_2, (\frac{1}{2} < \gamma_2 < 1)$, is the unique mixed Nash equilibrium of the game. We have proved the following proposition.

Proposition A.2. *If $\alpha \geq \gamma_1$, then $r(\alpha n) = 1$, if $\gamma_2 < \alpha < \gamma_1$, then $\lim_{n \rightarrow \infty} r(\alpha n) = 1$, if $\alpha < \gamma_2$, then $\lim_{n \rightarrow \infty} r(\alpha n) = 0$.*

Appendix B. Stationary states of irreducible Markov chains

The following tree representation of stationary states of Markov chains was proposed by Freidlin and Wentzell (1970, 1984). Let (Ω, P) be an irreducible Markov chain with a state space Ω and transition probabilities given by $P: \Omega \times \Omega \rightarrow [0, 1]$. It has a unique stationary probability distribution called a stationary state. For $x \in \Omega$, an x -tree is a directed graph on Ω such that from every $y \neq x$ there is a unique path to x and there is no outgoing edge out of x . Denote the set of all x -trees by $T(x)$ and let

$$q(x) = \sum_{d \in T(x)} \prod_{(y, y') \in d} P(y, y'), \tag{B.1}$$

where the product is with respect to all edges of d . Now one can show that

$$\mu(x) = \frac{q(x)}{\sum_{y \in \Omega} q(y)} \tag{B.2}$$

for all $x \in \Omega$.

A state is an absorbing one if it attracts nearby states in the mutation-free dynamics. We assume that after a finite number of steps of the mutation-free dynamics we arrive at one of the absorbing states (there are no other recurrence classes) and stay there forever. Then it follows from the above tree representation that any state different from absorbing states has zero probability in the stationary distribution in the zero-mutation limit. Moreover, in order to study the zero-mutation limit of the stationary state, it is enough to consider paths between absorbing states. More precisely, we construct x -trees with absorbing states as vertices; the family of such x -trees is denoted by $\tilde{T}(x)$. Let

$$q_m(x) = \max_{d \in \tilde{T}(x)} \prod_{(y, y') \in d} \tilde{P}(y, y'), \tag{B.3}$$

where $\tilde{P}(y, y') = \max \prod_{(w, w')} P(w, w')$, where the product is taken along any path joining y with y' and the maximum is taken with respect to all such paths. Now we may observe that if $\lim_{\varepsilon \rightarrow 0} q_m(y)/q_m(x) = 0$, for any $y \neq x$, then x is stochastically stable. Therefore, we have to compare trees with the biggest products in (B.3); we call such trees maximal.

Appendix C. Proof of Theorem 3

Pick δ, β , and ε . It follows from the limiting properties of $r(k)$ that there is $n(\varepsilon, \delta)$ and $\frac{1}{2} < \gamma_3 < \gamma_2$ such that for all $n > n(\varepsilon, \delta)$ we have that $r(\alpha n) < \varepsilon$ if $\alpha \leq \gamma_3$.

For any state in Ω' with $z = k$, we will prove that

$$q(k) < 3\epsilon q(k-1), \quad 1 \leq k \leq \gamma_3 n, \tag{C.1}$$

$$q(k) < \frac{2q(k-1)}{\epsilon}, \quad \gamma_3 n < k \leq n. \tag{C.2}$$

It follows from (C.1) and (C.2) that

$$\begin{aligned} \mu_n^\epsilon(z \leq \beta n) &= \frac{\sum_{0 \leq k \leq \beta n} \binom{n}{k} q(k)}{\sum_{0 \leq k \leq \beta n} \binom{n}{k} q(k) + \sum_{k > \beta n}^k \binom{n}{k} q(k) + \sum_{k > \gamma_3 n} \binom{n}{k} q(k)} \\ &> \frac{1}{1 + \sum_{k > \beta n}^k \binom{n}{k} (3\epsilon)^k + (3\epsilon)^{\gamma_3 n - 1} \sum_{k > \gamma_3 n} \binom{n}{k} \left(\frac{\epsilon}{3}\right)^{k - \gamma_3 n}} \\ &> \frac{1}{1 + \sum_{k > \beta n}^k \binom{n}{k} (3\epsilon)^k e^{k \left(\frac{n}{k}\right)^k} + \epsilon^{(2\gamma_3 - 1)n - 1} 3^{n-1} \left(\frac{\epsilon}{3} + 1\right)^n} \\ &> 1 - \delta \end{aligned} \tag{C.3}$$

if ϵ is small enough. Smaller β is and closer γ_3 is to $\frac{1}{2}$, the smaller ϵ should be.

To prove (C.1 and C.2), with every k -tree ($1 \leq k \leq \gamma_3 n$) we will associate a $(k-1)$ -tree. Let ω be a k -tree. We reverse arrows on all edges on the unique path between $k-1$ and k (all other edges we leave unchanged). (C.1) follows from the bound

$$\frac{r(k)(1-\epsilon) + (1-r(k))\epsilon}{(1-r(k-1))(1-\epsilon) + r(k-1)\epsilon} < 3\epsilon$$

and (C.2) from the bound

$$\frac{r(k)(1-\epsilon) + (1-r(k))\epsilon}{(1-r(k-1))(1-\epsilon) + r(k-1)\epsilon} > \frac{\epsilon}{2}. \quad \square$$

Proof of Theorem 4. In intermediate dynamics, the probability of moving m units to the right if $r(k) < \epsilon$ or to the left if $1 - r(k) < \epsilon$ is not proportional to ϵ^m as in sequential dynamics. Therefore, to prove (C.1–C.2) we cannot simply reverse arrows on edges in constructing corresponding trees.

To prove (C.1), with every k -tree ($1 \leq k \leq \gamma_3 n$) we will again associate a $(k-1)$ -tree. Let ω be a k -tree. If on the unique path between $k-1$ and k there are only transitions which involve only one individual at any time period, then we reverse arrows on all edges on this path as in the proof of Theorem 3. Otherwise, let an edge $j \rightarrow l$ be the first edge which involves at least two players. If $j > k-1$, then we reverse all arrows between $k-1$ and j , cut the edge $j \rightarrow l$ and connect k to $k-1$. Because an edge was deleted, a correspondence between k and $(k-1)$ -trees is not one-to-one anymore. If the

edge $j \rightarrow l$ involves m players, then there are at most $\binom{n}{m}$ k -trees with the same corresponding $(k-1)$ -tree. By cutting the edge we decreased a probability by at least τ^m times. If $\tau n^2 < 1/2$, then the series $\sum_{m \geq 2} \tau^m \binom{n}{m}$ is bounded by τ . C.1 follows.

If $j \leq k-1$, then we cut the edges $k-1 \rightarrow$ and $j \rightarrow l$, connect j to a state with $z = j-1$ (only one player changes his strategy) and k to $k-1$. By the above procedure we decreased a probability by τ . There are at most n^3 k -trees with the same corresponding $(k-1)$ -tree. If $\tau n^3 < \epsilon$, then (C.1) follows.

(C.2) can be proved in an analogous way. Now Theorem 4 follows in the same way as Theorem 3. \square

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