

# Confluence of an extension of Combinatory Logic by Boolean constants

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# Combinatory Logic with Boolean constants

The system CL-pc

$$K_{xy} \rightarrow x$$

$$S_{xyz} \rightarrow xz(yz)$$

$$CT_{xy} \rightarrow x$$

$$CF_{xy} \rightarrow y$$

$$Cz_{xx} \rightarrow x$$

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Not confluent (Klop, 1980).

# Combinatory Logic with Boolean constants

Conditional linearization: the system  $CL\text{-}pc^L$

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Conditional linearization: the system  $CL\text{-}pc^L$

$$\begin{aligned}Kxy &\rightarrow x \\Sxyz &\rightarrow xz(yz) \\CTxy &\rightarrow x \\CFxy &\rightarrow y \\Czxy &\rightarrow x \quad \Leftarrow \quad x = y \\Czxy &\rightarrow y \quad \Leftarrow \quad x = y\end{aligned}$$

Confluent (de Vrijer, 1990).

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$$\begin{aligned}Kxy &\rightarrow x \\Sxyz &\rightarrow xz(yz) \\CTxy &\rightarrow x \\CFxy &\rightarrow y \\Czxy &\rightarrow x \quad \Leftarrow \quad x = y \\Czxy &\rightarrow y \quad \Leftarrow \quad x = y\end{aligned}$$

Confluent (de Vrijer, 1990).

The confluence proof of de Vrijer is essentially based on a “semantic” argument.

# Combinatory Logic with Boolean constants

Conditional linearization: the system  $CL\text{-pc}^1$

$$\begin{aligned}K_{xy} &\rightarrow x \\S_{xyz} &\rightarrow xz(yz) \\CT_{xy} &\rightarrow x \\CF_{xy} &\rightarrow y \\Czxy &\rightarrow x \quad \Leftarrow \quad x = y\end{aligned}$$

# Combinatory Logic with Boolean constants

Conditional linearization: the system  $CL\text{-pc}^1$

$$\begin{aligned}K_{xy} &\rightarrow x \\S_{xyz} &\rightarrow xz(yz) \\CT_{xy} &\rightarrow x \\CF_{xy} &\rightarrow y \\Cz_{xy} &\rightarrow x \quad \Leftarrow \quad x = y\end{aligned}$$

Confluence of  $CL\text{-pc}^1$  appears as problem 15 on the RTA list of open problems.



# Combinatory Logic with Boolean constants

Conditional linearization: the system  $CL\text{-pc}^1$

$$\begin{aligned}Kxy &\rightarrow x \\Sxyz &\rightarrow xz(yz) \\CTxy &\rightarrow x \\CFxy &\rightarrow y \\Czxy &\rightarrow x \quad \Leftarrow \quad x = y\end{aligned}$$

Confluence of  $CL\text{-pc}^1$  appears as problem 15 on the RTA list of open problems.

We give a “syntactic” proof of confluence of  $CL\text{-pc}^1$ . This proof also works for  $CL\text{-pc}^L$ .

# Combinatory Logic with Boolean constants

## Conditional linearization

### Lemma

*The following are equivalent.*

- ▶  $t_1 =_{\text{CL-pc}} t_2$ ,
- ▶  $t_1 =_{\text{CL-pc}^L} t_2$ ,
- ▶  $t_1 =_{\text{CL-pc}^1} t_2$ .

# Combinatory Logic with Boolean constants

## Confluence

De Vrijer shows  $T \not\equiv_{\text{CL-pc}^L} F$  by a model construction within the Graph Model  $P\omega$ . Then confluence of  $\text{CL-pc}^L$  is proven using a simple technical argument with an auxiliary term rewriting system.

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We show: if  $t \equiv_{\text{CL-pc}^1} F$  then  $t \rightarrow_{\text{CL-pc}^1}^* F$ .

# Combinatory Logic with Boolean constants

## Confluence

De Vrijer shows  $T \not\equiv_{\text{CL-pc}^L} F$  by a model construction within the Graph Model  $P\omega$ . Then confluence of  $\text{CL-pc}^L$  is proven using a simple technical argument with an auxiliary term rewriting system.

We show: if  $t \equiv_{\text{CL-pc}^1} F$  then  $t \rightarrow_{\text{CL-pc}^1}^* F$ .

Then confluence of  $\text{CL-pc}^1$  is proven using an argument with an auxiliary system similar to that of de Vrijer.

# Confluence of CL-pc<sup>1</sup>

Difficulty

$$\begin{aligned}Kxy &\rightarrow x \\Sxyz &\rightarrow xz(yz) \\CTxy &\rightarrow x \\CFxy &\rightarrow y \\Czxy &\rightarrow x \quad \Leftarrow \quad x = y\end{aligned}$$

If  $t_1 =_{\text{CL-pc}^1} t_2$  then  $CFt_1t_2 \rightarrow_{\text{CL-pc}^1} t_1$  and  $CFt_1t_2 \rightarrow_{\text{CL-pc}^1} t_2$ .

# Confluence of CL-pc<sup>1</sup>

If  $q =_{\text{CL-pc}^1} F$  then  $q \rightarrow_{\text{CL-pc}^1}^* F$ .

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# Confluence of $\text{CL-pc}^1$

If  $q =_{\text{CL-pc}} F$  then  $q \rightarrow_{\text{CL-pc}^1}^* F$ .

Given a conversion  $q =_{\text{CL-pc}} F$  we consider a certain set  $\mathcal{S}(q =_{\text{CL-pc}} F)$  of reductions in  $\text{CL-pc}^L$  from  $q$  to  $F$  such that it contains at least one reduction in  $\text{CL-pc}^1$ .

# Confluence of CL-pc<sup>1</sup>

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1. if  $\mathcal{S}(q =_{\text{CL-pc}} F)$  is defined and  $q \rightarrow_{\text{CL-pc}} q'$ , then we need to define  $\mathcal{S}(q' \text{ CL-pc} \leftarrow q =_{\text{CL-pc}} F)$ ,
2. if  $\mathcal{S}(q =_{\text{CL-pc}} F)$  is defined and  $q \text{ CL-pc} \leftarrow q'$ , then we need to define  $\mathcal{S}(q' \rightarrow_{\text{CL-pc}} q =_{\text{CL-pc}} F)$ .

## The system $\text{CL-pc}^5$

The set  $\mathcal{S}(q =_{\text{CL-pc}} F)$  is defined by labelling certain constants in  $q$  and defining permissible reductions on a labelled variant of  $q$  by a system  $\text{CL-pc}^5$ .

## The system CL-pc<sup>5</sup>

The set  $\mathcal{S}(q =_{\text{CL-pc}} F)$  is defined by labelling certain constants in  $q$  and defining permissible reductions on a labelled variant of  $q$  by a system CL-pc<sup>5</sup>.

$$\begin{array}{ll} C_1 T_1 xy \rightarrow x & C_2 zxy \rightarrow x \Leftarrow |x| =_{\text{CL-pc}^1} |y| \\ C_1 F_1 xy \rightarrow y & C_2 zxy \rightarrow y \Leftarrow |x| =_{\text{CL-pc}^1} |y| \\ K_1 xy \rightarrow x & \end{array}$$

and e.g.

$$\begin{array}{ll} S^{1,1} xy_1 \langle z_{0,1}, z_{1,1} \rangle \rightarrow xz_{0,1}(y_1 z_{1,1}) \\ S^{1,2,1} x \langle y_1, y_2 \rangle \langle z_{0,1}, z_{1,1}, z_{1,2}, z_{2,1} \rangle \rightarrow xz_{0,1} \langle y_1 \langle z_{1,1}, z_{1,2} \rangle, y_2 z_{2,1} \rangle \\ S^{2,2} xy_1 \langle z_{0,1}, z_{0,2}, z_{1,1}, z_{1,2} \rangle \rightarrow x \langle z_{0,1}, z_{0,2} \rangle (y_1 \langle z_{1,1}, z_{1,2} \rangle) \end{array}$$

where  $|z_{i,j}| =_{\text{CL-pc}^1} |z'_{i,j}|$  and  $|y_i| =_{\text{CL-pc}^1} |y_j|$ .

## Significant and insignificant subterms

- ▶ Significant (sub)terms (*s*-terms): “labelled” subterms, e.g.,  $C_1t_1t_2t_3$ ,  $C_1t$ ,  $K_1t$ , etc., CL-*pc*-reductions/expansions on their erasures are translated to *s*-reductions/expansions (i.e. reductions/expansions in CL-*pc*<sup>s</sup>).

## Significant and insignificant subterms

- ▶ Significant (sub)terms (*s*-terms): “labelled” subterms, e.g.,  $C_1 t_1 t_2 t_3$ ,  $C_1 t$ ,  $K_1 t$ , etc., CL-*pc*-reductions/expansions on their erasures are translated to *s*-reductions/expansions (i.e. reductions/expansions in CL-*pc*<sup>5</sup>).
- ▶ Insignificant (sub)terms (*i*-terms): “unlabelled” subterms, CL-*pc*-reductions/expansions inside them don't matter – they are translated to *i*-reductions/expansions (i.e. reductions/expansions in CL-*pc*<sup>1</sup>).

# Translating the CL- $\rho$ -conversion

- ▶ CL- $\rho$ -reductions/expansions of  $i$ -terms are translated to  $i$ -reductions/expansions.

## Translating the CL- $\text{pc}$ -conversion

- ▶ CL- $\text{pc}$ -reductions/expansions of  $i$ -terms are translated to  $i$ -reductions/expansions.
- ▶ CL- $\text{pc}$ -reductions of (erasures of)  $s$ -terms are translated to  $s$ -reductions.



# Translating the CL- $\text{pc}$ -conversion

- ▶ CL- $\text{pc}$ -reductions/expansions of  $i$ -terms are translated to  $i$ -reductions/expansions.
- ▶ CL- $\text{pc}$ -reductions of (erasures of)  $s$ -terms are translated to  $s$ -reductions.
- ▶ CL- $\text{pc}$ -expansions of (erasures of)  $s$ -terms are translated to  $a$ -expansions.

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- ▶ CL- $\text{pc}$ -reductions/expansions of  $i$ -terms are translated to  $i$ -reductions/expansions.
- ▶ CL- $\text{pc}$ -reductions of (erasures of)  $s$ -terms are translated to  $s$ -reductions.
- ▶ CL- $\text{pc}$ -expansions of (erasures of)  $s$ -terms are translated to  $a$ -expansions.

$a$ -expansion is like  $s$ -expansion but it also specifies the labelling of the expansion (and some other technicalities).

## Example

$$\begin{aligned} F \leftarrow C(KF\Omega)FF \leftarrow C(KF\Omega)F(CTF(KF\Omega)) \rightarrow CFF(CTF(KF\Omega)) \leftarrow \\ CFF(C(KTF)F(KF\Omega)) \rightarrow C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

## Example

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# Example

$$\begin{aligned} F_1 & \dashv\vdash C_2(KF\Omega)F_1F_1 \\ & \succ C(KF\Omega)FF \leftarrow C(KF\Omega)F(CTF(KF\Omega)) \rightarrow CFF(CTF(KF\Omega)) \leftarrow \\ & CFF(C(KTF)F(KF\Omega)) \rightarrow C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ & KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1F_1$$

$$\begin{aligned} &> C(KF\Omega)F \leftarrow C(KF\Omega)F(CTF(KF\Omega)) \rightarrow CFF(CTF(KF\Omega)) \leftarrow \\ &CFF(C(KTF)F(KF\Omega)) \rightarrow C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ &KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \text{ } \leftarrow C_2(KF\Omega)F_1 F_1 \text{ } \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \\ \succ C(KF\Omega)F(CTF(KF\Omega)) \rightarrow CFF(CTF(KF\Omega)) \leftarrow \\ CFF(C(KTF)F(KF\Omega)) \rightarrow C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

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$$\begin{aligned} F_1 \dashv\vdash C_2(KF\Omega)F_1 F_1 \dashv\vdash C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \\ \succ C(KF\Omega)F(CTF(KF\Omega)) \rightarrow CFF(CTF(KF\Omega)) \dashv\vdash \\ CFF(C(KTF)F(KF\Omega)) \rightarrow C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \dashv\vdash \\ KF(CF) \dashv\vdash SKCF \dashv\vdash SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$



## Example

$$\begin{aligned} F_1 \text{ }_a \leftarrow C_2(KF\Omega)F_1 F_1 \text{ }_a \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow; \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \\ \succ CFF(CTF(KF\Omega)) \leftarrow CFF(C(KTF)F(KF\Omega)) \rightarrow \\ C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

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# Example

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$$\begin{aligned} &F_1 \leftarrow C_2(KF\Omega)F_1 F_1 \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ &C_2FF_1(C_1T_1F_1(KF\Omega)) \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ &C_1(K_1T_1F)F_1(KF\Omega) \\ &\succ C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ &KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

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$$\begin{aligned} F_1 \leftarrow C_2(KF\Omega)F_1 F_1 \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \\ \succ C(KTF)F(KF\Omega) \rightarrow C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1 F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ a} \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \\ \succ C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1 F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ a} \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \\ \succ C(KTF)FF \rightarrow F \leftarrow \\ KF(CF) \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$



# Example

$$\begin{aligned} F_1 \xrightarrow{a} C_2(KF\Omega)F_1 & F_1 \xrightarrow{a} C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) & \xrightarrow{a} C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) & \rightarrow_i C_1(K_1T_1F)F_1F \xrightarrow{s^*} F_1 \\ \succ F \leftarrow KF(CF) & \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \xrightarrow{a} C_2(KF\Omega)F_1 & F_1 \xrightarrow{a} C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) & \xrightarrow{a} C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) & \rightarrow_i C_1(K_1T_1F)F_1F \xrightarrow{s}^* F_1 \\ \succ F \leftarrow KF(CF) & \leftarrow SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

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# Example

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# Example

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# Example

$$\begin{aligned} F_1 \text{ }_a\leftarrow C_2(KF\Omega)F_1 F_1 \text{ }_a\leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ }_a\leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \rightarrow_s^* F_1 \text{ }_a\leftarrow \\ K_1F_1(CF) \text{ }_a\leftarrow S^{1,1}K_1C(F_1, F) \\ \succ SKCF \leftarrow SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1 F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ a} \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \rightarrow_s^* F_1 \text{ a} \leftarrow \\ K_1F_1(CF) \text{ a} \leftarrow S^{1,1}K_1C\langle F_1, F \rangle \text{ a}, i \leftarrow^* \\ S^{1,1}K_1C\langle K_1F_1\Omega, KF\Omega \rangle \\ \succ SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1 F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ a} \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \rightarrow_s^* F_1 \text{ a} \leftarrow \\ K_1F_1(CF) \text{ a} \leftarrow S^{1,1}K_1C\langle F_1, F \rangle \text{ a}, i \leftarrow^* \\ S^{1,1}K_1C\langle K_1F_1\Omega, KF\Omega \rangle \\ \succ SKC(KF\Omega) \rightarrow SKCF \end{aligned}$$



# Example

$$\begin{aligned} F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1 F_1 \text{ a} \leftarrow C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \rightarrow_i \\ C_2FF_1(C_1T_1F_1(KF\Omega)) \text{ a} \leftarrow C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \rightarrow_s \\ C_1(K_1T_1F)F_1(KF\Omega) \rightarrow_i C_1(K_1T_1F)F_1F \rightarrow_s^* F_1 \text{ a} \leftarrow \\ K_1F_1(CF) \text{ a} \leftarrow S^{1,1}K_1C\langle F_1, F \rangle \text{ a}, i \leftarrow^* \\ S^{1,1}K_1C\langle K_1F_1\Omega, KF\Omega \rangle \rightarrow_{s,i}^* S^{1,1}K_1C\langle F_1, F \rangle \\ \succ SKCF \end{aligned}$$

# Example

$$\begin{aligned} F_1 \xrightarrow{a\leftarrow} C_2(KF\Omega) F_1 & F_1 \xrightarrow{a\leftarrow} C_2(KF\Omega) F_1 (C_1 T_1 F_1 (KF\Omega)) \xrightarrow{i} \\ C_2 F F_1 (C_1 T_1 F_1 (KF\Omega)) & \xrightarrow{a\leftarrow} C_2 F F_1 (C_1 (K_1 T_1 F) F_1 (KF\Omega)) \xrightarrow{s} \\ C_1 (K_1 T_1 F) F_1 (KF\Omega) & \xrightarrow{i} C_1 (K_1 T_1 F) F_1 F \xrightarrow{s^*} F_1 \xrightarrow{a\leftarrow} \\ K_1 F_1 (CF) \xrightarrow{a\leftarrow} S^{1,1} K_1 C \langle F_1, F \rangle & \xrightarrow{a,i^*} \\ S^{1,1} K_1 C \langle K_1 F_1 \Omega, KF\Omega \rangle & \xrightarrow{s,i^*} S^{1,1} K_1 C \langle F_1, F \rangle \end{aligned}$$

# Confluence of CL-pc<sup>1</sup>

We show: for  $t$  “standard” (obtained by the process just indicated): if  $t \rightarrow_s^! t'$  then  $t' \equiv F_1$ .

# Confluence of CL- $\text{pc}^1$

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From this it follows that there is an  $s$ -reduction which may be “erased” to a CL- $\text{pc}^1$ -reduction from  $t$  to  $F$ .

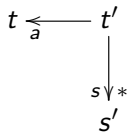
## Example

$$\begin{aligned} F_1 \xleftarrow{a} C_2(KF\Omega) F_1 & \xleftarrow{a} C_2(KF\Omega) F_1 (C_1 T_1 F_1(KF\Omega)) \xrightarrow{i} \\ C_2 F F_1 (C_1 T_1 F_1(KF\Omega)) & \xleftarrow{a} C_2 F F_1 (C_1 (K_1 T_1 F) F_1(KF\Omega)) \xrightarrow{s} \\ C_1 (K_1 T_1 F) F_1(KF\Omega) & \xrightarrow{i} C_1 (K_1 T_1 F) F_1 F \xrightarrow{s}^* F_1 \xleftarrow{a} \\ K_1 F_1(CF) & \xleftarrow{a} S^{1,1} K_1 C \langle F_1, F \rangle \xleftarrow{a,i}^* \\ S^{1,1} K_1 C \langle K_1 F_1 \Omega, KF\Omega \rangle & \xrightarrow{s,i}^* S^{1,1} K_1 C \langle F_1, F \rangle \end{aligned}$$

## Example

$$\begin{aligned} F_1 \xrightarrow{s} C_2(KF\Omega) F_1 & \xrightarrow{a} C_2(KF\Omega) F_1(C_1 T_1 F_1(KF\Omega)) \xrightarrow{i} \\ C_2 F F_1(C_1 T_1 F_1(KF\Omega)) & \xrightarrow{a} C_2 F F_1(C_1(K_1 T_1 F) F_1(KF\Omega)) \xrightarrow{s} \\ C_1(K_1 T_1 F) F_1(KF\Omega) & \xrightarrow{i} C_1(K_1 T_1 F) F_1 F \xrightarrow{s}^* F_1 \xrightarrow{a} \\ K_1 F_1(CF) & \xrightarrow{a} S^{1,1} K_1 C\langle F_1, F \rangle \xrightarrow{a,i}^* \\ S^{1,1} K_1 C\langle K_1 F_1 \Omega, KF\Omega \rangle & \xrightarrow{s,i}^* S^{1,1} K_1 C\langle F_1, F \rangle \end{aligned}$$

# Commutation of $a$ -reduction and $s$ -reduction







# Commutation of $a$ -reduction and $s$ -reduction

$$\begin{array}{ccc} t & \xleftarrow{a} & t' \\ \downarrow & & \downarrow \\ s & \xleftarrow{a} & s' \end{array}$$

The diagram shows a commutative square. The top-left node is  $t$ , the top-right node is  $t'$ , the bottom-left node is  $s$ , and the bottom-right node is  $s'$ . A solid arrow labeled  $a$  points from  $t'$  to  $t$ . A solid arrow labeled  $a$  points from  $s'$  to  $s$ . A solid arrow labeled  $s$  points from  $t$  down to  $s$ . A solid arrow labeled  $s$  points from  $t'$  down to  $s'$ . A dashed arrow labeled  $s$  points from  $s$  to  $s'$ . A dashed arrow labeled  $a$  points from  $s$  to  $t$ .

Hence if  $t \rightarrow_s^* s'$  with  $s'$  in  $s$ -normal form then  $s' = F_1$ .

## Example

$$\begin{aligned} F_1 \xrightarrow{s} C_2(KF\Omega) F_1 F_1 \xrightarrow{a} C_2(KF\Omega) F_1 (C_1 T_1 F_1 (KF\Omega)) \xrightarrow{i} \\ C_2 F F_1 (C_1 T_1 F_1 (KF\Omega)) \xrightarrow{a} C_2 F F_1 (C_1 (K_1 T_1 F) F_1 (KF\Omega)) \xrightarrow{s} \\ C_1 (K_1 T_1 F) F_1 (KF\Omega) \xrightarrow{i} C_1 (K_1 T_1 F) F_1 F \xrightarrow{s}^* F_1 \xrightarrow{a} \\ K_1 F_1 (CF) \xrightarrow{a} S^{1,1} K_1 C \langle F_1, F \rangle \xrightarrow{a,i}^* \\ S^{1,1} K_1 C \langle K_1 F_1 \Omega, KF\Omega \rangle \xrightarrow{s,i}^* S^{1,1} K_1 C \langle F_1, F \rangle \end{aligned}$$

## Example

$$\begin{aligned} F_1 \xleftarrow{s} C_2(KF\Omega)F_1 & F_1 \xleftarrow{s} C_2(KF\Omega)F_1(C_1T_1F_1(KF\Omega)) \xrightarrow{i} \\ C_2FF_1(C_1T_1F_1(KF\Omega)) & \xleftarrow{a} C_2FF_1(C_1(K_1T_1F)F_1(KF\Omega)) \xrightarrow{s} \\ C_1(K_1T_1F)F_1(KF\Omega) & \xrightarrow{i} C_1(K_1T_1F)F_1F \xrightarrow{s}^* F_1 \xleftarrow{a} \\ K_1F_1(CF) & \xleftarrow{a} S^{1,1}K_1C\langle F_1, F \rangle \xleftarrow{a,i}^* \\ S^{1,1}K_1C\langle K_1F_1\Omega, KF\Omega \rangle & \xrightarrow{s,i}^* S^{1,1}K_1C\langle F_1, F \rangle \end{aligned}$$

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## Example

$$\begin{aligned} F_1 \xleftarrow{s} C_2 F F_1 \xleftarrow{s} C_2 F F_1 (C_1 T_1 F_1 (K F \Omega)) \\ \xleftarrow{a} C_2 F F_1 (C_1 (K_1 T_1 F) F_1 (K F \Omega)) \xrightarrow{s} \\ C_1 (K_1 T_1 F) F_1 (K F \Omega) \xrightarrow{i} C_1 (K_1 T_1 F) F_1 F \xrightarrow{s}^* F_1 \xleftarrow{a} \\ K_1 F_1 (C F) \xleftarrow{a} S^{1,1} K_1 C \langle F_1, F \rangle \xleftarrow{a,i}^* \\ S^{1,1} K_1 C \langle K_1 F_1 \Omega, K F \Omega \rangle \xrightarrow{s,i}^* S^{1,1} K_1 C \langle F_1, F \rangle \end{aligned}$$

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# Example

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# Example

$$\begin{aligned} F_1 s \leftarrow C_1 T_1 F_1 F \quad F s \leftarrow C_1 (K_1 T_1 F) F_1 F \\ \rightarrow_s^* F_1 a \leftarrow \\ K_1 F_1 (CF) a \leftarrow S^{1,1} K_1 C \langle F_1, F \rangle a, i \leftarrow^* \\ S^{1,1} K_1 C \langle K_1 F_1 \Omega, KF \Omega \rangle \rightarrow_{s, i}^* S^{1,1} K_1 C \langle F_1, F \rangle \end{aligned}$$

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# Example

$$F_1 \xleftarrow{S} K_1 F_1(CF) \xleftarrow{S} S^{1,1} K_1 C \langle F_1, F \rangle$$

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- ▶ Formalization available at:

<http://www.mimuw.edu.pl/~lukaszcz/clc.tar.gz>