

# Infinite Automata 2025/26

## Exercise Sheet 9

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**Exercise 9.1.** Show that semilinear sets are closed under complement. In other words, if  $S \subseteq \mathbb{N}^d$  is a semilinear set, then  $\mathbb{N}^d \setminus S$  is a semilinear set.

*Hint 1. It suffices to prove that the complement of a linear set with linearly independent periods is a semilinear set.*

*Hint 2. Let  $C$  be the closure of a given linear set  $L$  with linearly independent periods (i.e.  $C$  is the smallest integer cone that contains  $L$  and does not have any gap inside). The complement of  $\mathbb{N}^d \setminus L$  can be described as  $(\mathbb{N}^d \setminus C) \cup (C \setminus L)$ .*

*Hint 3. Notice that the set  $C \setminus L$  is a linear set with the same periods as  $L$  but different bases.*

**Exercise 9.2.** Show that semilinear sets are closed under difference. In other words, if  $S_1, S_2 \subseteq \mathbb{N}^d$  are semilinear sets then  $S_1 \setminus S_2$  is a semilinear set.

**Exercise 9.3.** Show that the coverability problem for binary-encoded 1-VASS is in PSPACE.

*Hint. Show that if there is a run witnessing coverability, then there must be a run witnessing coverability that has at most exponential length.*

**Exercise 9.4.** Show that the coverability problem for binary-encoded 1-VASS is in NP.

*Hint. Prove that if there is a run witnessing coverability, then there must be a run witnessing coverability that take the form: a short path, followed by a short positive cycle taken at most exponentially many times, followed by a short path.*

**Exercise 9.5.** Show that the coverability problem for binary-encoded 1-VASS is in P.

*Hint. Prove that, actually the witnesses consisting of a short path, followed by a short positive cycle, followed by a short path can be checked by a polynomial time dynamic programming algorithm.*