Coq - introduction

Daria Walukiewicz-Chrząszcz

19 march 2019

Coq — interactive proof assistant

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(***********)
(*  v      *)
(* <O___,, *)
(*   \VV/  *)
(*    //   *)
(*         *)
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http://coq.inria.fr/

- rich (pure) functional programming language
- rich logical language
- user writes proofs
- Coq makes sure every step is correct
- and solves subgoals for which automated proving algorithms have been implemented
- (proved to be correct) program can be extracted to Ocaml, Haskell, Scheme...

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Proving in Coq

Curry-Howard isomorphism

proof	of	a	given	formula
$\updownarrow$				$\updownarrow$
term		of the corresponding		type

$$\lambda x^{A\to B\to C}\lambda y^{A\to B}\lambda z^{A}\, xz(yz) \;:\; (A\to B\to C)\to(A\to B)\to(A\to C)$$

Correctness of Coq relies on correctness of type-checking

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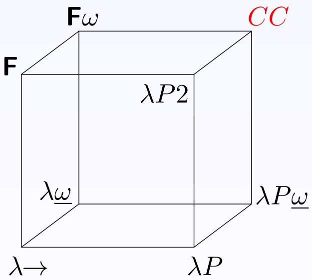
Coq — formalism

Coq — calculus of constructions (CC) + inductive definitions



$\uparrow$ polimorphism

$\nearrow$ type constructors

$\rightarrow$ dependent types

Big picture of Coq architecture

The De Bruijn principle (“small” core, externally checkable terms)

- core / kernel (≈20KLOC), responsible for:
  - CIC typing
  - reduction
  - environment (definitions, axioms etc).
  - modules
- the rest (≈230KLOC), responsible for:
  - user interface
  - file management
  - sections
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  - proof mode (plus tactics, tactic language)
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Coq — a bit of history

1984 CoC - calculus of constructions - G. Huet, T. Coquand

1989 first public release (version 4.10)

1991 Coq - calculus of inductive constructions - C. Paulin (version 5.6)

⋮

2000 version 7.0 with new (safer) architecture

2003 version 7.4 with modules

2004 version 8.0 with new syntax

2009 version 8.2 with "type classes"

2012 version 8.4 with eta-reduction, structural proof syntax...

2018 version 8.7.2 — fixes a critical bug in the universes (present since 8.5)

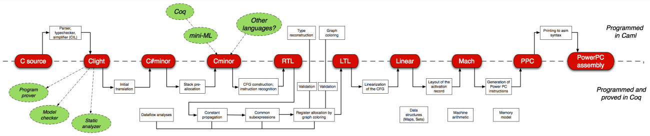
Coq — famous formalizations

- Fundamental theorem of algebra, Nijmegen 2000
- JavaCard Platform formalization, Trusted Logic 2003

  > *September 2007: a big step in program certification in the real world: The Technology and Innovation group at Gemalto has successfully completed a Common Criteria (CC) evaluation on a JavaCard based commercial product. This evaluation is the world's first CC certificate of a Java product involving EAL7 components.* (the official press release)

- Four color theorem, Cambridge 2004
- CompCert certified Clight compiler, 2008-now

  > *The main result of the project is the CompCert C verified compiler, a high-assurance compiler for almost all of the ISO C90 / ANSI C language, generating efficient code for the PowerPC, ARM and x86 processors.*



Coq — programming language

- predicative sorts `Set` and `Type`
- abstraction and application
- inductive types,
- (structural) recursion
- polimorphism
- dependant types and dependent pattern-matching
- modules i functors
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- intuitionistic higher-order logic
- impredicative sort `Prop`
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- interactive proof mode (goal management)
- built-in tactics (constructing a bit of proof-term): intro, apply, etc.
- automatic ad-hoc tactics: auto, intuition, etc.
- decision procedures: omega, ring, field, tauto, etc.
- tactic language (`Ltac mytactic:=...`)

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- program is extracted from the proof
- extracted program satisfies its specification *by definition*
- extraction — "elimination" of logical parts from the proof-term
- extraction possible because proofs are done in constructive logic (excluded-middle and double negation laws do not hold)
- target languages: O'Caml, Haskell, Scheme

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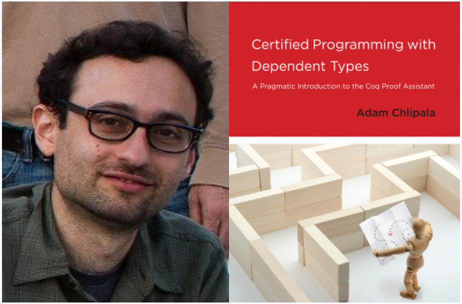

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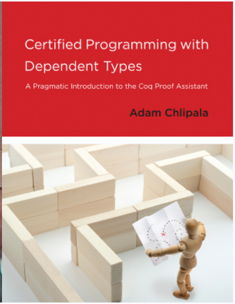

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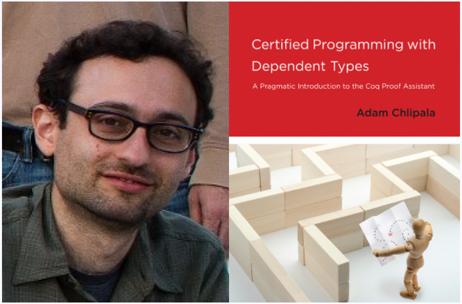

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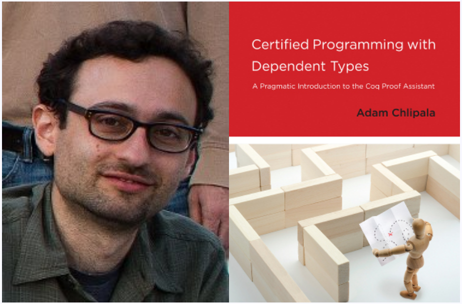

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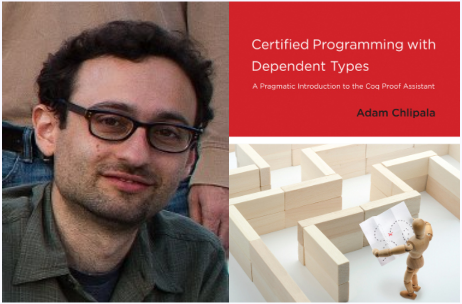

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Typing

$$\texttt{environment} \vdash \texttt{term} : \texttt{type}$$

environment: global and local declarations and definitions

types are terms and have types, ex. `nat:Set`, `Set:Type`$_1$

but there are terms that are not types, ex. `fun n:nat => n`

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Coq — formalism: fun for all

simple types abstraction rule:

$$\frac{\Gamma, x{:}A \vdash M : B}{\Gamma \vdash \lambda x{:}A.M \;:\; A \to B}$$

dependent types abstraction rule:

$$\frac{\Gamma, x{:}A \vdash M : B(x)}{\Gamma \vdash \lambda x{:}A.M \;:\; \forall x{:}A.B(x)}$$

Shorthand: $A \to B$ is $\forall x{:}A.B$, where $x \notin FV(B)$

concrete Coq syntax:

`fun n:nat =>` M `: forall n:nat, vector n`

application rule:

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Coq — formalism: fun for all

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Coq — typing rules: sorts

- Sorts in Coq:

$$\begin{matrix} \mathit{Prop} \\ \mathit{Set} \end{matrix} \; : \mathit{Type}_1 : \mathit{Type}_2 : \ldots$$

- Cummulativity (or sub-sorting):

$$\mathit{Prop} \leq \mathit{Set} \leq \mathit{Type}_1 \leq \mathit{Type}_2 \leq \ldots$$

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Coq — products

product rule

$$\frac{\Gamma \vdash A : s_1 \qquad \Gamma, x{:}A \vdash B : s_2}{\Gamma \vdash \forall x{:}A.B : s_2}$$ if s_1 and s_2 satisfy ...

- $s_1 \leq s_2$, or
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Coq — reductions

- beta

  $(\lambda x{:}A.M)N \longrightarrow_{\beta} M[N/x]$

- eta expansion (if M is of a functional type)

  $M \longrightarrow_{\eta} \lambda x{:}A.Mx$

- delta

  (definition unfolding)

- zeta

  `(let x:=N in M)` $\longrightarrow_{\zeta}$ `M[N/x]`

- iota

  (inductive types reductions — soon :)

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`vector nat 4` $=_{iota}$ `vector nat (2+2)`

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Coq implements intuitionistic logic

- forall and implication are built-in
- False, conjunction, disjunction cannot be defined from $\rightarrow$
- they are defined as inductive types
- negation is defined $\neg\phi \equiv \phi \rightarrow False$
- existential quantifier cannot be defined from universal one
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- one can use classical logic - axioms needed (ex: excluded middle)
- proof-checking is decidable (not provability)

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True and False

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Inductive False : Prop :=.
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Inductive True : Prop :=
   I : True.
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True and False

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Inductive False : Prop :=.
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Conjunction and disjunction

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Inductive and (A B : Prop) : Prop :=
   conj : A -> B -> and A B
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$\wedge$ is an infix notation for and

```
Inductive or (A B : Prop) : Prop :=
   or_introl : A -> or A B
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Existential quantifier

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Inductive ex (A : Type) (P : A -> Prop) : Prop :=
   ex_intro : forall x : A, P x -> ex A P.
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`exists y, P y` is a notation for `ex`

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