

Membership problems in braid and Artin groups

Robert D. Gray¹

(joint work with Carl-Fredrik Nyberg-Brodde)

Warsaw Algebra Seminar, October 2025



Engineering and
Physical Sciences
Research Council



¹Research supported by EPSRC Fellowship EP/V032003/1 'Algorithmic, topological and geometric aspects of infinite groups, monoids and inverse semigroups'.

How little we really know!

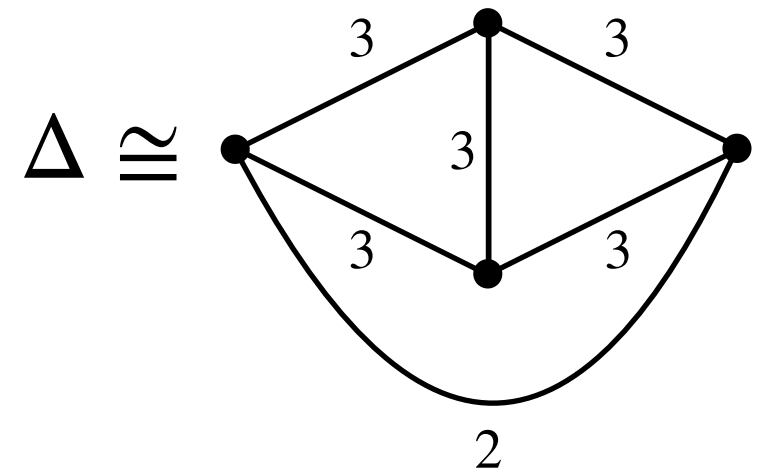
Let G be the group defined by the presentation:

$$G \cong \langle a, b, x, y \mid \begin{array}{l} ab = ba, \quad axa = xax, \quad aya = yay, \\ bxb = xbx, \quad byb = yby, \quad xyx = yxy \end{array} \rangle$$

Open problem

Does this group G have decidable word problem? That is, is there an algorithm that takes any two words α, β over $\{a, b, x, y\}^{\pm 1}$ and decides whether or not $\alpha = \beta$ in G ?

- This is a particular case of a more general problem which asks whether the word problem is decidable for all [Artin groups](#).



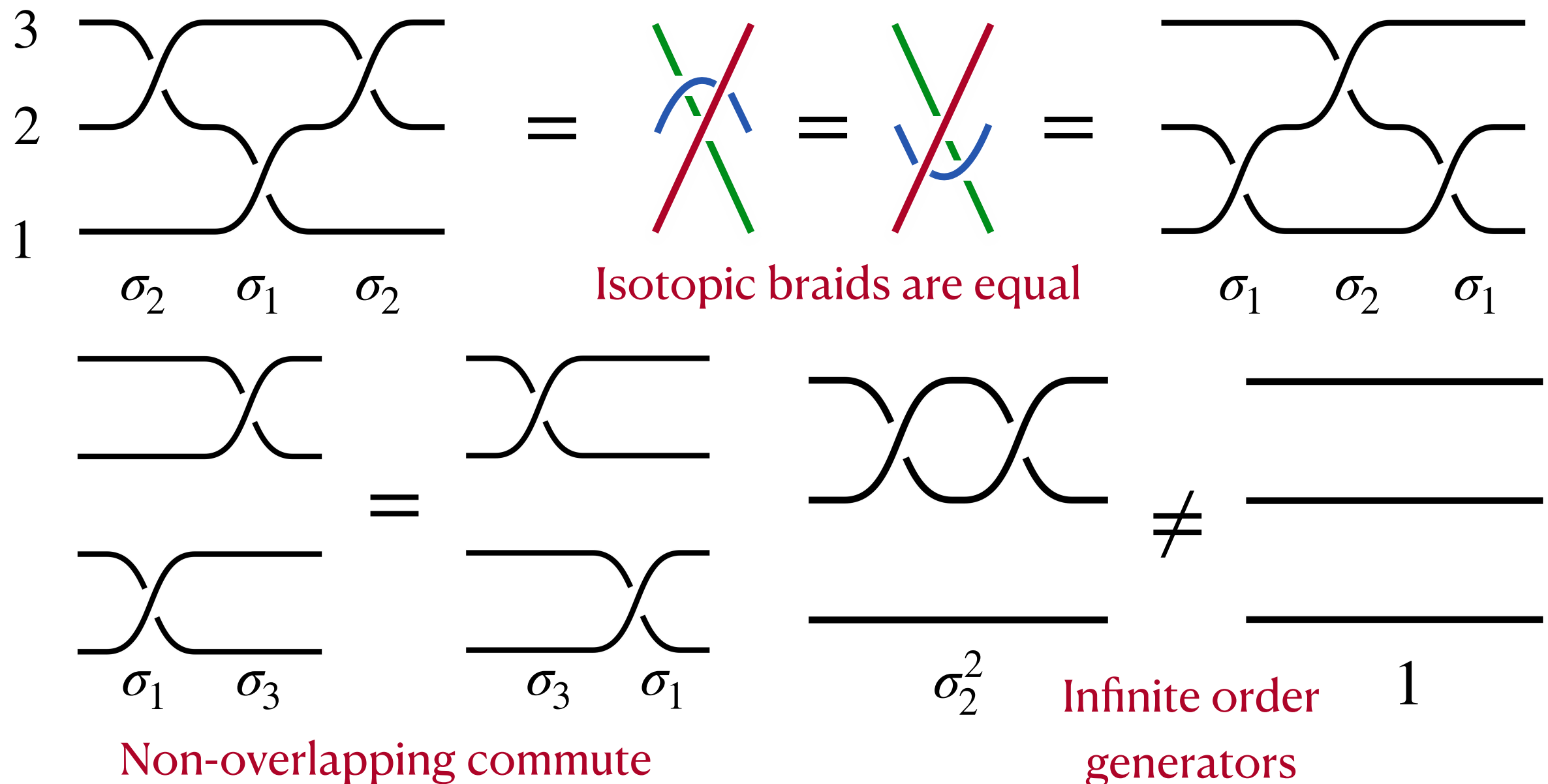
$$G = A(\Delta)$$

The braid group B_n

The braid group B_n has a finite Artin presentation

$$B_n = \langle \sigma_1, \dots, \sigma_{n-1} \mid \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}, \quad \sigma_i \sigma_j = \sigma_j \sigma_i \text{ if } |i - j| \geq 2 \rangle$$

Representing elements graphically



Braid group history and motivation

- Formally introduced and studied by [Artin \(1925\)](#) & implicit in work of [Hurwitz \(1891\)](#) as fundamental groups of configurations spaces.
- Connections with representation theory, algebraic geometry and topology, mathematical physics, Yang—Baxter equation, etc...

Algorithmic properties of Braid groups

- Word problem is decidable ([Artin \(1925\)](#)).
- Conjugacy problem is decidable ([Garside \(1969\)](#)).
- They are bi-automatic ([Charney \(1992\)](#)) implying that their word problems are solvable in quadratic time and their conjugacy problems are solvable in exponential time.
- The more general “equation solving” problem for braid groups remains open.

The membership problem in B_n

The **membership problem** in B_n asks if a given braid τ (the **target**) can be written as a product of some other given set of braids $\delta_1, \dots, \delta_k$ (the **generators**).

Examples

$$\delta_1 = \text{braid diagram: two strands crossing, top-left to bottom-right, bottom-left to top-right}$$

$$\delta_2 = \text{braid diagram: two strands crossing, top-right to bottom-left, bottom-right to top-left}$$

$$\tau = \text{braid diagram: a sequence of four crossings between two strands}$$

$$\delta_2 = \sigma_2 \sigma_2$$

$$\delta_1 = \sigma_1 \sigma_2$$

$$\delta_2 \delta_1 = \sigma_2 \sigma_2 \sigma_1 \sigma_2$$

$$= \tau$$

$$\delta_2 = \sigma_2 \sigma_2$$

$$\delta_1 = \sigma_1 \sigma_2$$

$$\delta_2 \delta_1 = \sigma_2 \sigma_2 \sigma_1 \sigma_2$$

$$= \sigma_2 \sigma_1 \sigma_2 \sigma_1 = \tau$$

On the other hand

$$\tau' := \text{braid diagram: two strands crossing, top-right to bottom-left, bottom-right to top-left} = \delta_2^{-1}$$

cannot be written as a product of δ_1, δ_2 (e.g. since defining relations are length preserving)

The membership problem in B_n

Definition the braid group B_n has **decidable** (subsemigroup) **membership problem** if there is an algorithm solving:

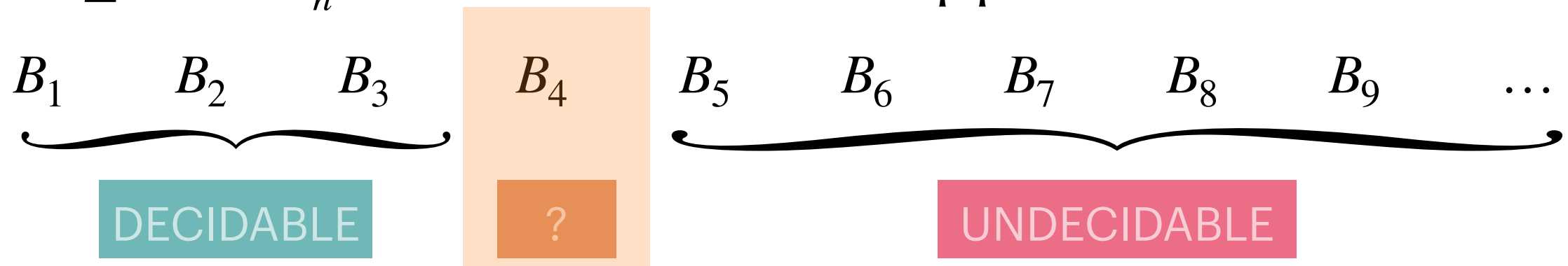
INPUT: A element $\tau \in B_n$ and elements $\delta_1, \dots, \delta_k \in B_n$.

QUESTION: $\tau \in \langle \delta_1, \dots, \delta_k \rangle \leq B_n$?

where $\langle \delta_1, \dots, \delta_k \rangle \leq B_n$ is the subsemigroup generated by $\{\delta_1, \dots, \delta_k\}$.

Theorem (Potapov 2013)

1. If $n \leq 3$ then B_n has decidable membership problem.
2. If $n \geq 5$ then B_n has undecidable membership problem.



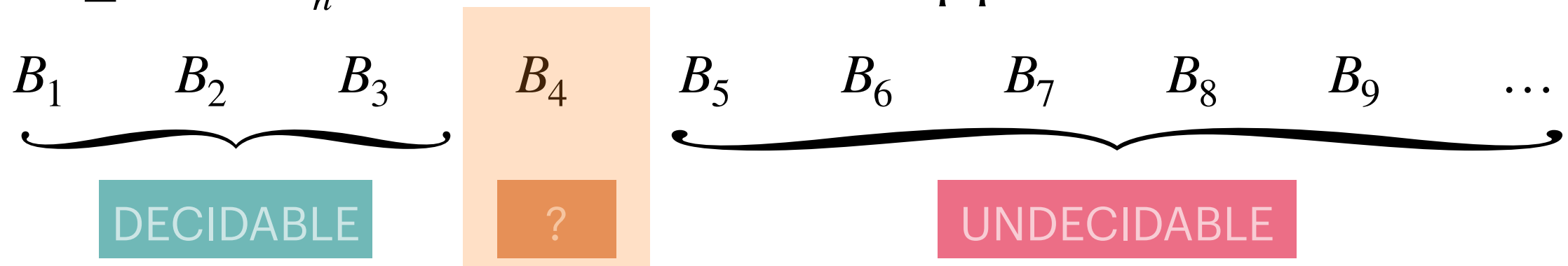
Question (Potapov 2013 + Ko & Potapov 2017)

Does B_4 has decidable membership problem?

The membership problem for small n

Theorem (Potapov 2013)

1. If $n \leq 3$ then B_n has decidable membership problem.
2. If $n \geq 5$ then B_n has undecidable membership problem.



Small n $|B_1| = 1$ and $B_2 \cong \mathbb{Z}$ ✓

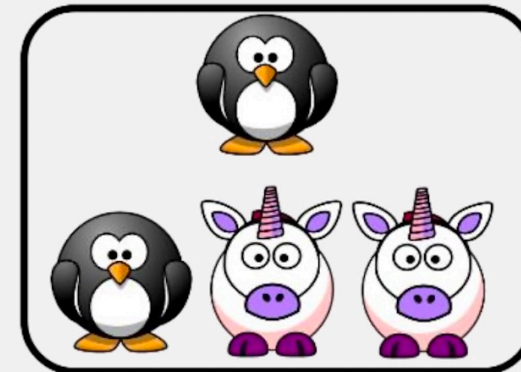
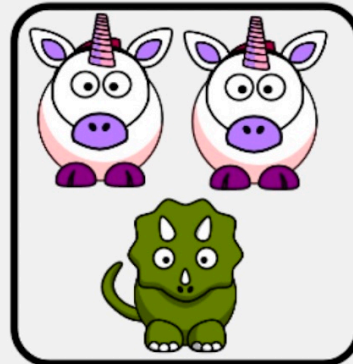
$B_3 \cong \langle x, y \mid x^2 = y^3 \rangle$ is a torus knot group

$\Rightarrow B_3$ is virtually $F_n \times \mathbb{Z}$ with F_n a free group (by e.g. [Niblo and Wise \(2001\)](#))

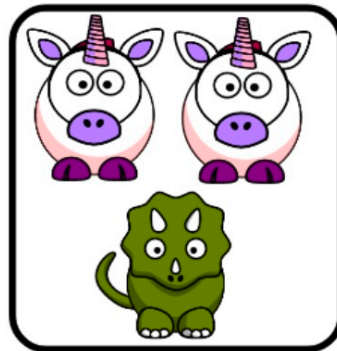
$\Rightarrow B_3$ has decidable membership problem (by [Benois \(1969\)](#) & [Kambites et al. \(2007\)](#))

Animal dominoes

Dominoes

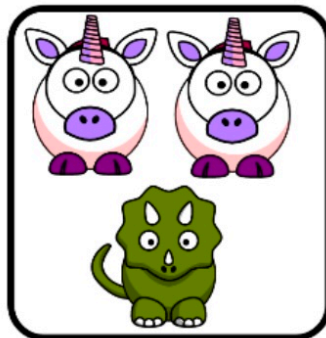
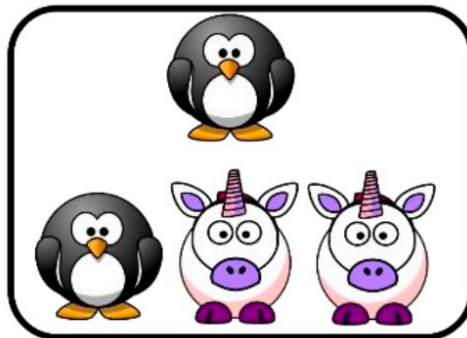


The game: Line up dominoes so the top sequence matches the bottom sequence.



Incorrect!

The top does not match the bottom



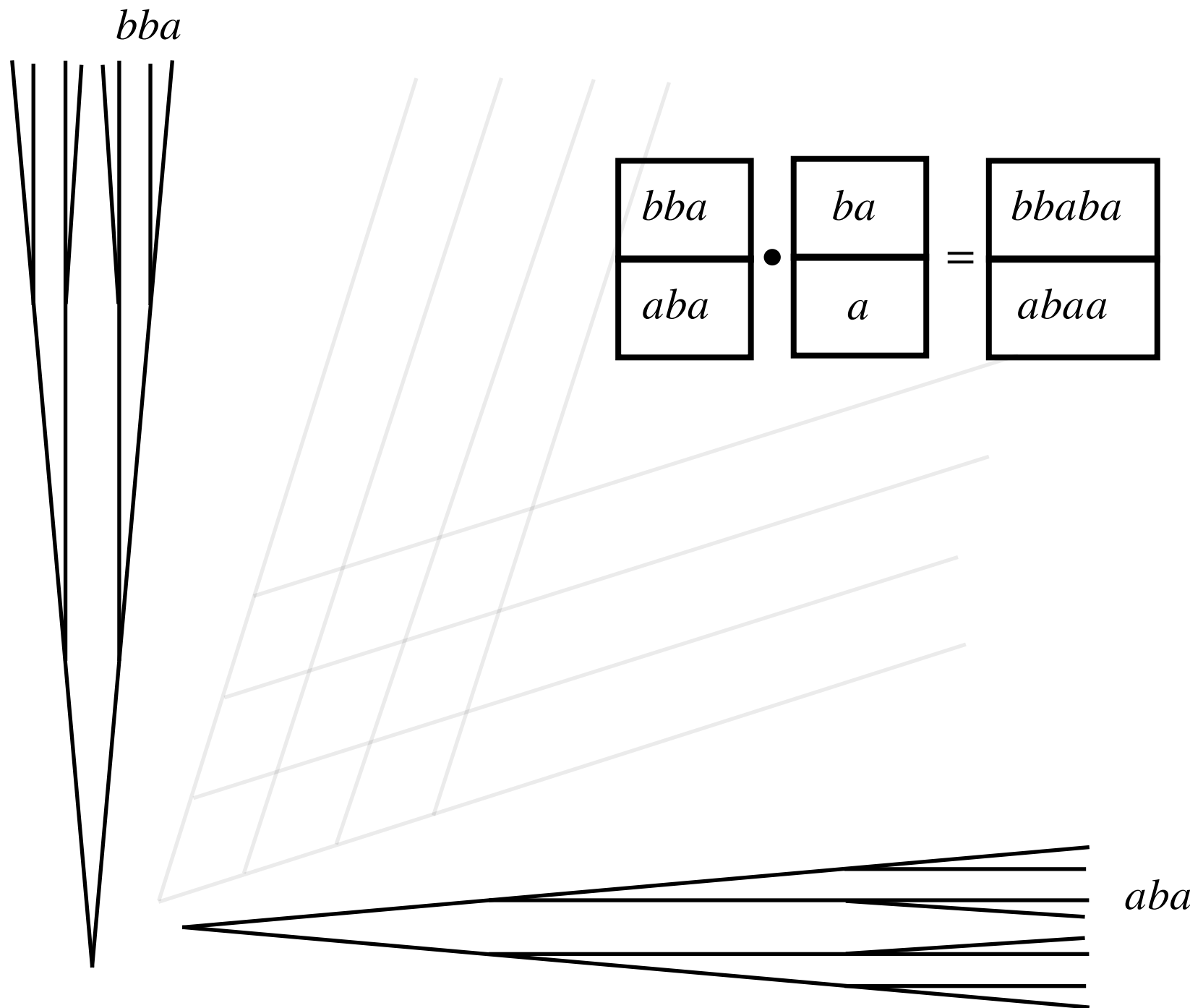
Correct!

The top matches the bottom

penguin.unicorn.unicorn.dinosaur.penguin

This is the **Post Correspondence Problem** and is known to be undecidable e.g. there is no algorithm that will take a finite of dominoes and tell you whether or not there is a correct line using these dominoes (each domino can be used as many times as you want).

Large n : Embedding pairs of words $F_2 \times F_2 \hookrightarrow B_n$



- $F_2 = F(a, b)$ = free group
- [Makanina \(1981\)](#) proved the **direct product** $F_2 \times F_2 \hookrightarrow B_n$ when $n \geq 5$.
- Can be used to simulate animal dominoes game i.e. post correspondence problem.
- [Mikhailova \(1958\)](#) showed $F_2 \times F_2$ contains a fixed subgroup in which membership is undecidable.

[Corollary \(Makanina \(1981\) & Mikhailova \(1958\)\)](#)

If $n \geq 5$ then B_n has undecidable membership problem i.e. there is **no algorithm** that takes input $\delta_1, \dots, \delta_k \in B_n$ and $\tau \in B_n$ and decides whether $\tau \in \langle \delta_1, \dots, \delta_k \rangle$.

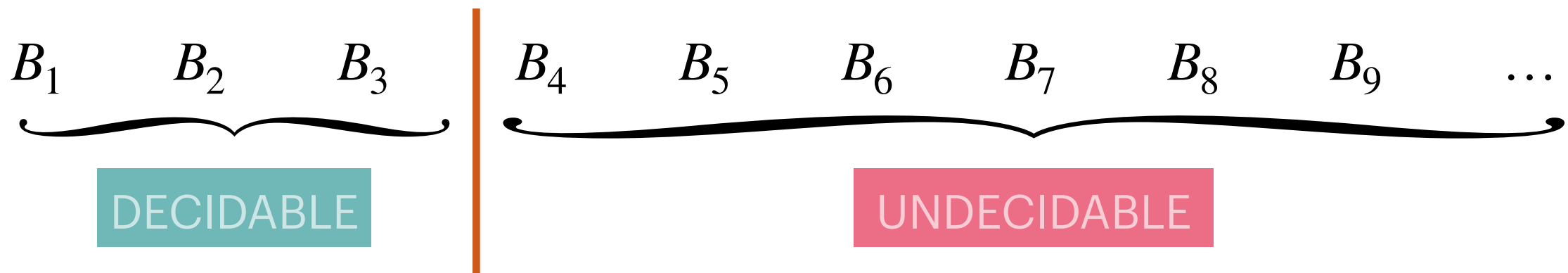
The membership problem in B_4

Quoting from Ko & Potapov 2017

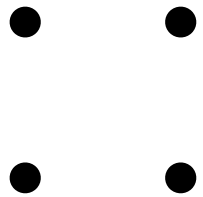
“...there is **no embedding** from a set of pairs of words into B_4 ; by [Akimenkov \(1991\)](#). Hence, these problems might be decidable for B_4 since our undecidability proofs for B_5 essentially rely on the embedding from a set of pairs of words into B_5 .”

Theorem (RDG and Nyberg-Brodde 2025)

The braid group B_4 has an undecidable membership problem. In fact, there is a fixed finitely generated subsemigroup of B_4 in which membership is undecidable. Hence **B_n has decidable membership problem if and only if $n \leq 3$.**



Membership problem in Right-angled Artin groups (RAAGs)



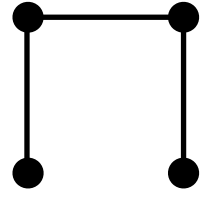
$$\langle a_i \mid \rangle$$

Benois (1969)

Free groups



DECIDABLE



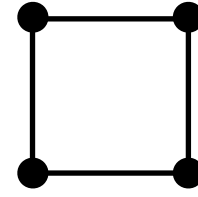
$$\langle a_i \mid a_i a_{i+1} = a_{i+1} a_i \rangle$$

Lohrey &
Steinberg (2008)

$$A(P_4)$$



UNDECIDABLE



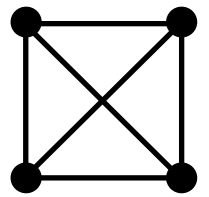
$$\langle a_i \mid a_i a_{i+1} = a_{i+1} a_i \pmod{4} \rangle$$

Mikhailova (1958)

$$A(C_4) = F_2 \times F_2$$



UNDECIDABLE



$$\langle a_i \mid a_i a_j = a_j a_i \rangle$$

Grunschlag (1999)

f.g. abelian groups



DECIDABLE

Definition $A(\Gamma) = \text{Gp}\langle V\Gamma \mid uv = vu \text{ if and only if } \{u, v\} \in E\Gamma \rangle$.

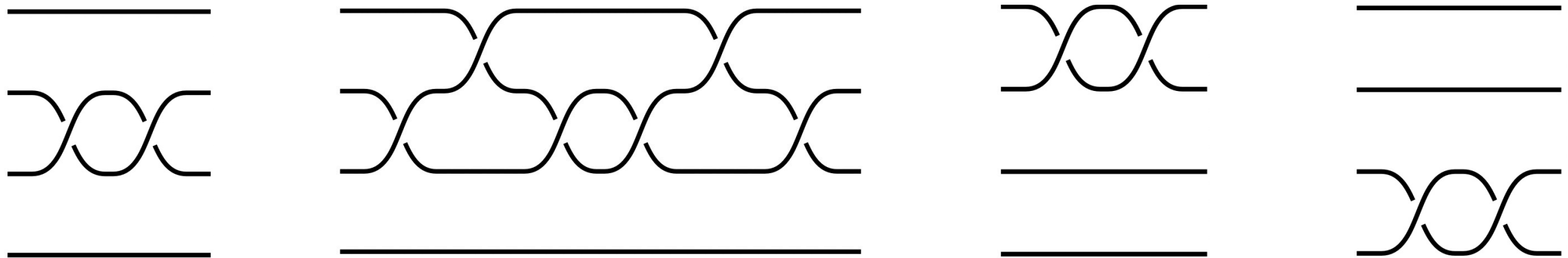
Theorem (Lohrey & Steinberg (2008)) A RAAG $A(\Gamma)$ has decidable membership problem $\Leftrightarrow \Gamma$ does not embed P_4 or C_4 as an induced subgraph.

- In fact they show that $A(P_4)$ has a fixed finitely generated subsemigroup in which membership is undecidable.

Tree groups and B_4

Theorem (Droms, Lewin & Servatius (1991))

The subgroup of B_4 generated by σ_2^2 , $(\sigma_2\sigma_3\sigma_2)^2$, σ_3^2 , σ_1^2 is isomorphic to $A(P_4)$.



Proof that B_4 has undecidable membership problem:

Lohrey & Steinberg (2008) $\Rightarrow \exists$ f.g. $S \leq A(P_4)$ in which membership is undecidable.

Droms, Lewin & Servatius (1991) $\Rightarrow A(P_4) \hookrightarrow B_4$.

Hence \exists f.g. $S \hookrightarrow A(P_4) \hookrightarrow B_4$ in which membership is undecidable.

Therefore B_4 has undecidable membership problem.



Artin groups

Let Γ be a finite graph with edges labelled by natural numbers greater than or equal to 2. The **Artin group** $A(\Gamma)$ is the group defined by the presentation with generating set the set $V\Gamma$ of vertices of Γ and a defining relation

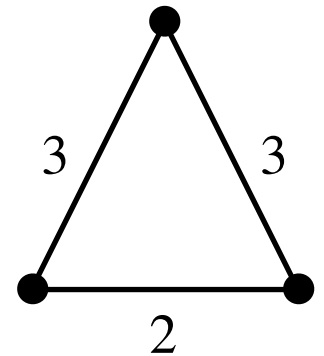
$$\underbrace{abababab\dots}_{m \text{ letters}} = \underbrace{babababa\dots}_{m \text{ letters}}$$

for each edge that connects vertices a and b and is labelled m .

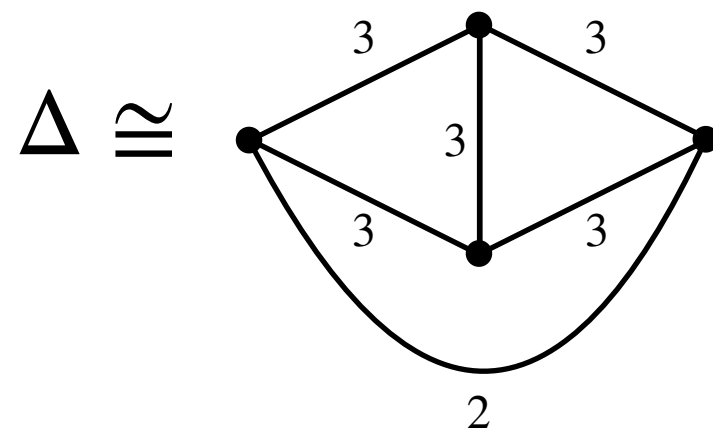
Examples

– RAAGs = Artin groups where all labels equal 2.

– Every braid group B_n is an Artin group e.g. $B_4 = A(\Gamma)$ for $\Gamma =$



– $A(\Delta)$ has presentation



$$\langle a, b, x, y \mid ab = ba, axa = xax, aya = yay, \\ bxb = xbx, byb = yby, xyx = yxy \rangle$$

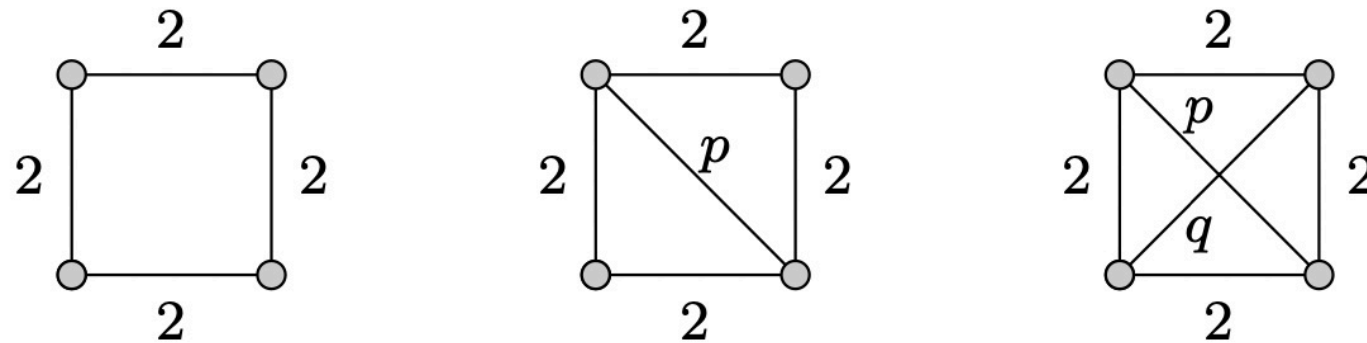
Artin groups

Theorem (RDG and Nyberg-Brodde 2025)

Let $A = A(\Gamma)$ be an Artin group. Then the following are equivalent:

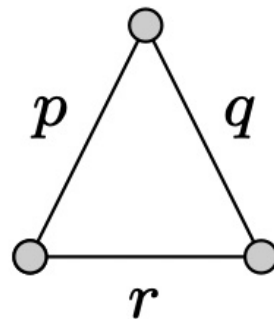
- (i) A has decidable membership problem
- (ii) the graph Γ does not embed any of the following as an induced subgraph:

- (a) a (generalized) square of one of the following three forms



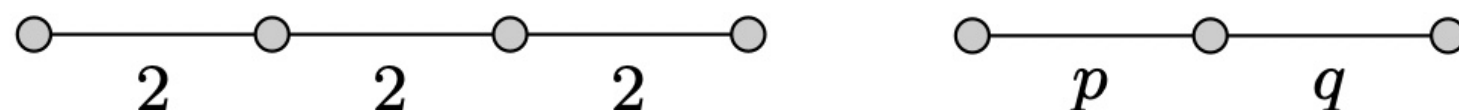
with $p > 2$ and $q > 2$,

- (b) a triangle of the form to 2, or



where at most one of $\{p, q, r\}$ is equal

- (c) a path of one of the following two forms

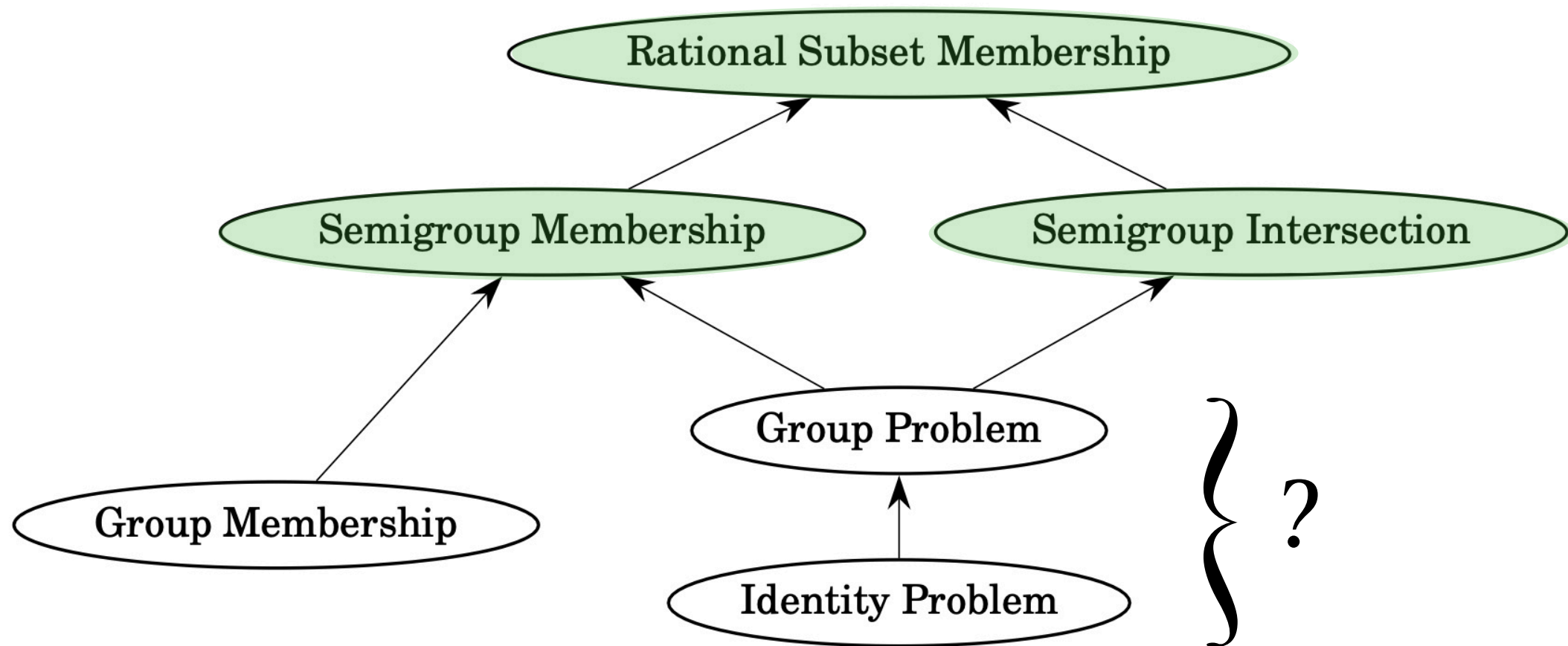


where at most one of $\{p, q\}$ is equal to 2.

Artin groups - notes on the proof

- A group G is **subgroup separable** (also called LERF) if for every finitely generated subgroup H of G , and element $g \in G \setminus H$ there is a finite index subgroup K of G , which contains H but $g \notin K$.
- In general for finitely presented groups:
[subgroup separable \nRightarrow decidable membership problem] and
[decidable membership problem \nRightarrow subgroup separable].
- However, a key step in the proof of our result above is showing that for Artin groups subgroup separable is equivalent to having decidable membership problem.
- Then we may apply results of [Almeida and Lima \(2021 + 2024\)](#) who classified the subgroup separable Artin groups.
- See also [Foniqi \(2024\)](#) for further exploration of these connections.

Other properties for Artin groups



Theorem (RDG and Nyberg-Brodde 2025)

Let $A = A(\Gamma)$ be an Artin group. Then the following are equivalent:

- (i) A has decidable (subsemigroup) membership problem;
- (ii) A has decidable rational subset membership problem;
- (iii) A has decidable semigroup intersection problem;
- (iv) Γ does not embed any of the forbidden subgraphs in the above theorem.

The identity problem for B_n

Definition the braid group B_n has **decidable identity problem** if there is an algorithm solving:

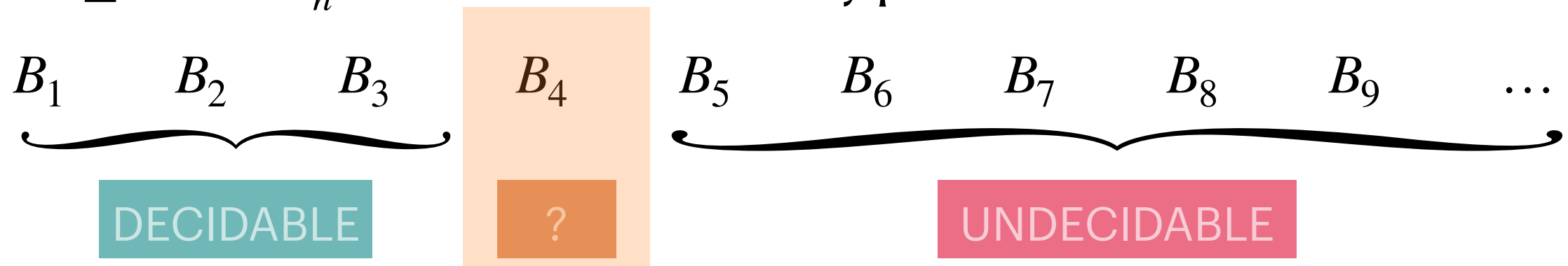
INPUT: The identity element $1 \in B_n$ and elements $\delta_1, \dots, \delta_k \in B_n$.

QUESTION: $1 \in \langle \delta_1, \dots, \delta_k \rangle \leq B_n$?

where $\langle \delta_1, \dots, \delta_k \rangle \leq B_n$ is the subsemigroup generated by $\{\delta_1, \dots, \delta_k\}$.

Theorem (Potapov 2013)

1. If $n \leq 3$ then B_n has decidable identity problem.
2. If $n \geq 5$ then B_n has undecidable identity problem.



Question (Potapov 2013 + Ko & Potapov 2017)

Does B_4 has decidable identity problem?

This problem remains open!

Undecidable fixed target membership problem

Theorem (RDG and Nyberg-Brodde 2025)

Let τ_0 be the following fixed element from the braid group B_4

$$\tau_0 = \begin{array}{c} \text{Diagram of a braid with 4 strands and 6 crossings} \end{array} = \sigma_2^2 \sigma_1^2 \sigma_2^{-2} \sigma_1^{-2}$$

Then there is **no algorithm** that can solve the following:

INPUT: The fixed element τ_0 above, and elements $\delta_1, \dots, \delta_k \in B_4$.

QUESTION: $\tau_0 \in \langle \delta_1, \dots, \delta_k \rangle \leq B_n$?

where $\langle \delta_1, \dots, \delta_k \rangle \leq B_n$ is the subsemigroup generated by $\{\delta_1, \dots, \delta_k\}$.

Remark

The identity problem for B_4 asks whether we can replace τ_0 by the identity braid $\begin{array}{c} \text{Diagram of the identity braid with 4 strands} \end{array} \in B_4$ in the above result?

Fixed target membership problem - proof ideas

Lemma (RDG and Nyberg-Brodde 2025)

Let G be a finitely generated group, and let $F_{x,y,z}$ be the free group with basis $\{x, y, z\}$. If the fixed-target membership problem for xy^{-1} in the free product $G * F_{x,y,z}$ is decidable, then the (subsemigroup) membership problem in G is decidable.

Corollary

The fixed-target membership problem for xy^{-1} in $A(P_4) * F_{x,y,z}$ is undecidable.

Final step of the B_4 undecidable fixed target membership problem proof

Adapt results of [Droms, Lewin & Servatius \(1991\)](#) to find appropriate embeddings

$$A(P_4) * F_{x,y,z} \hookrightarrow A(P_\infty) \hookrightarrow A(P_4) \hookrightarrow B_4$$

where $\tau_0 = xy^{-1}$ and P_∞ is the bi-infinite line



Open problems / future directions

- Does B_4 has decidable identity problem?
- Classify the Artin groups with decidable identity problem.
- Does B_4 have decidable subgroup membership problem?
 - **Note:** The RAAG $A(C_5) \hookrightarrow B_4$ and it is unknown whether $A(C_5)$ has decidable subgroup membership problem.
- Classify Coxeter groups with decidable subsemigroup membership problem.
- In particular does the the right-angled Coxeter group of the pentagon with presentation (subscripts taken mod 5)
$$\langle a_0, a_1, a_2, a_3, a_4 \mid a_i^2 = 1, \quad a_i a_{i+1} = a_{i+1} a_i \quad (i \in \{0,1,2,3,4\}) \rangle$$
have decidable subsemigroup membership problem?

- **Note:** This group contains a hyperbolic surface subgroup of finite index and the subsemigroup membership problem is open for surface groups.