Unambiguous Büchi is weak

Henryk Michalewski Michał Skrzypczak

DLT 2016
Montreal

Languages of **infinite trees**

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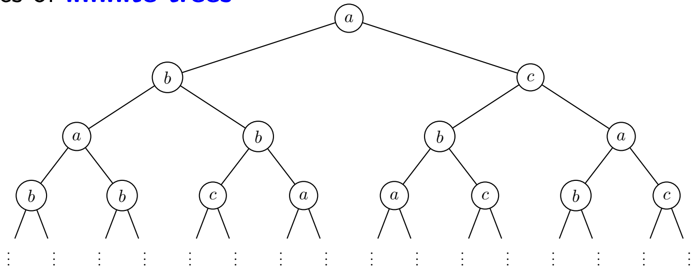

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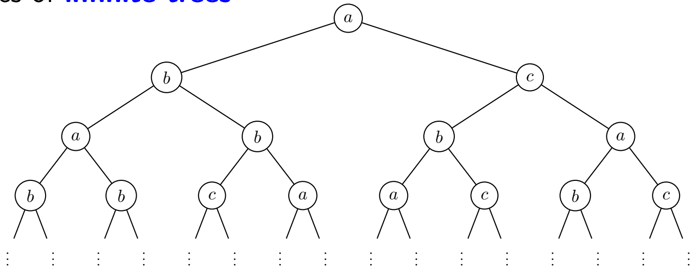

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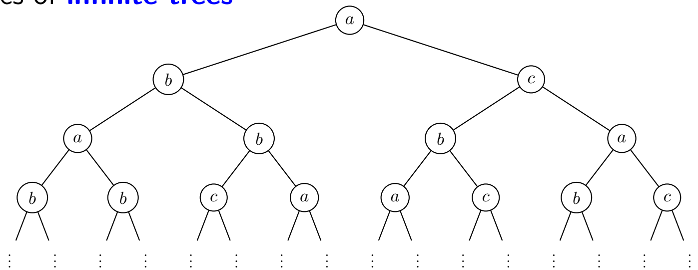

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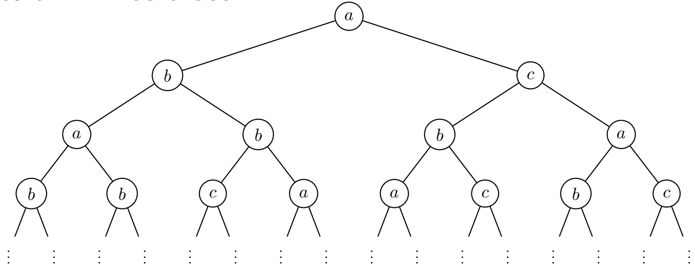

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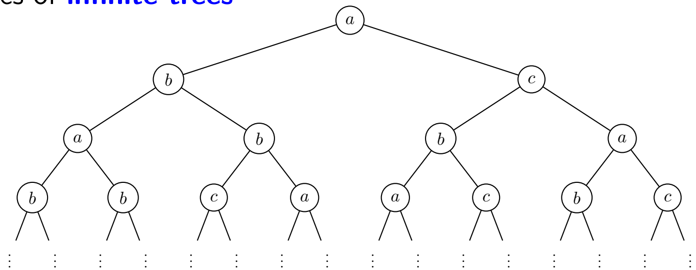

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Proof

Automata...

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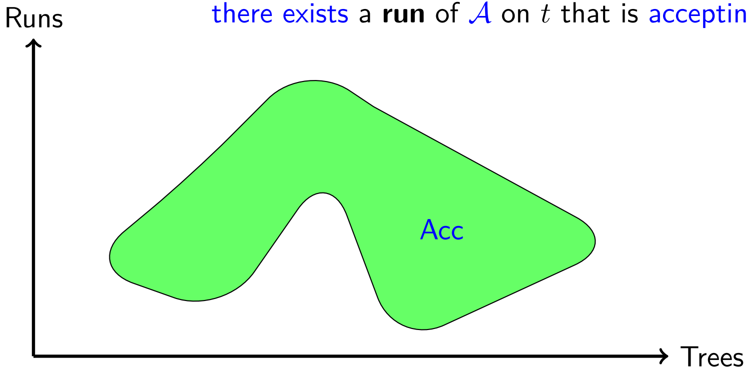

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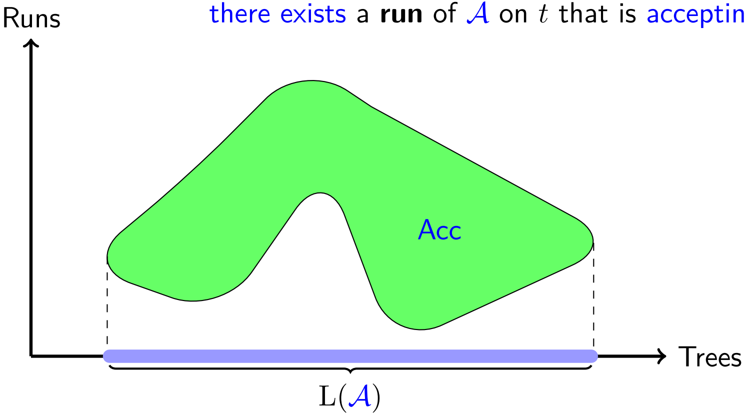

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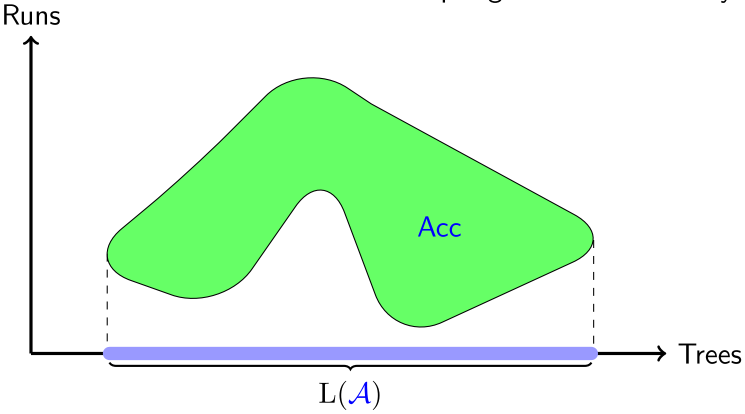

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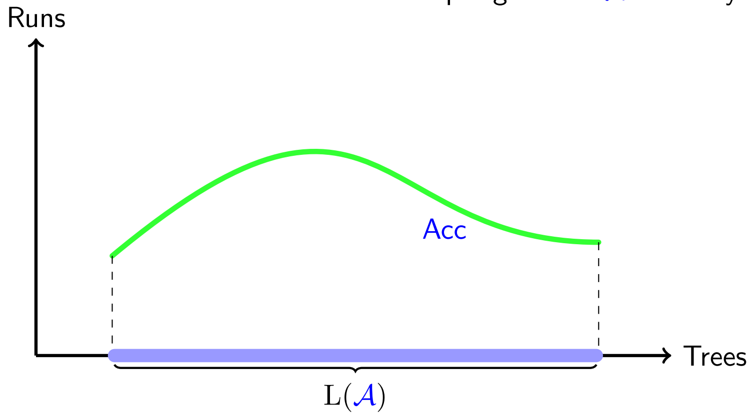

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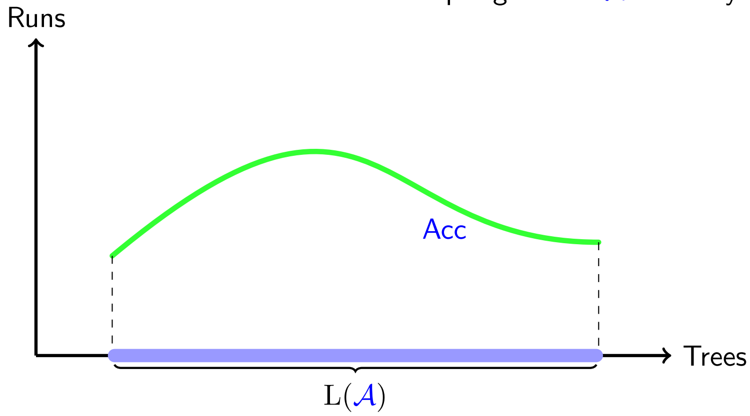

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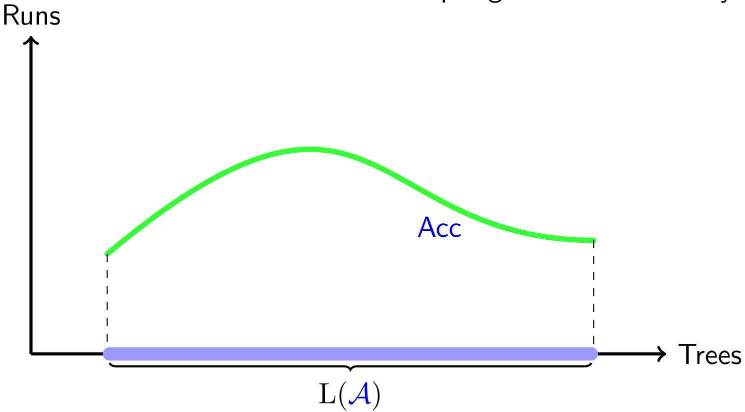

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(e.g. Stearns, Hunt [1985])

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$\rightsquigarrow$ **Goal**: understand unambiguous languages of infinite trees

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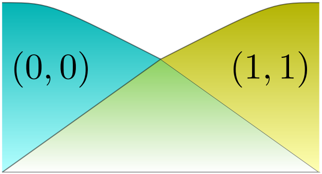

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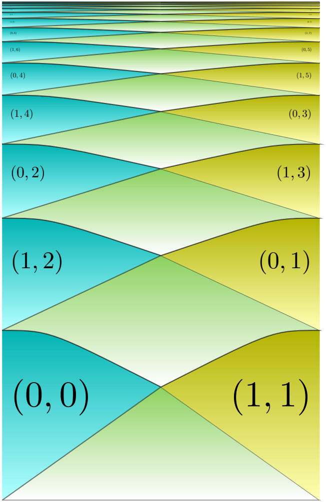

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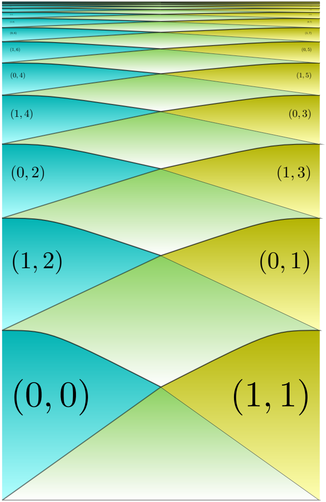

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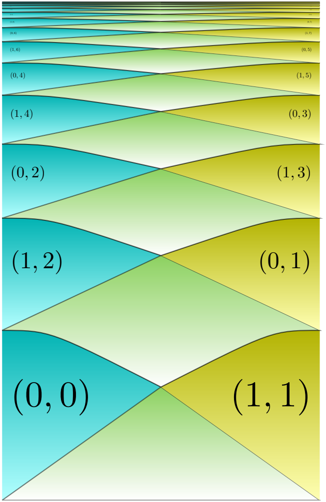

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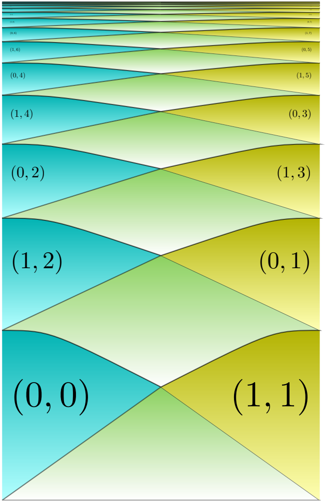

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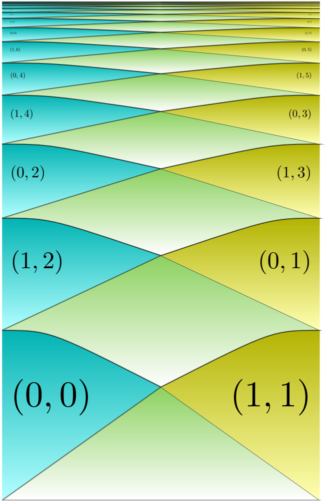

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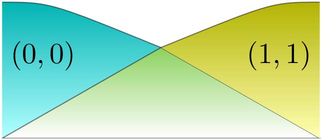

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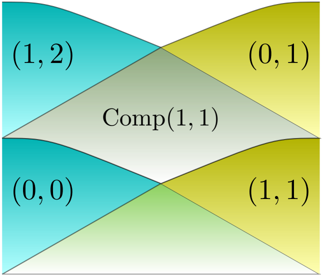

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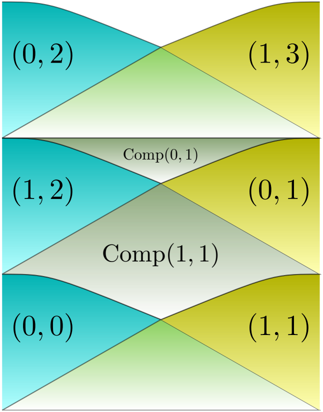

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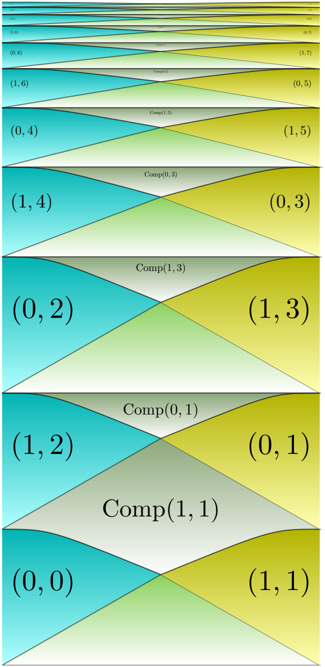

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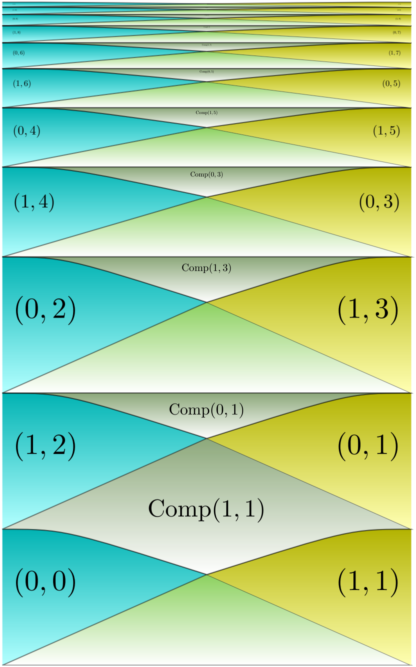

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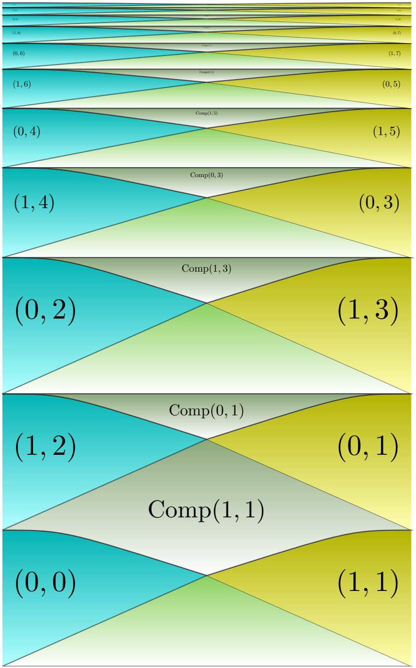

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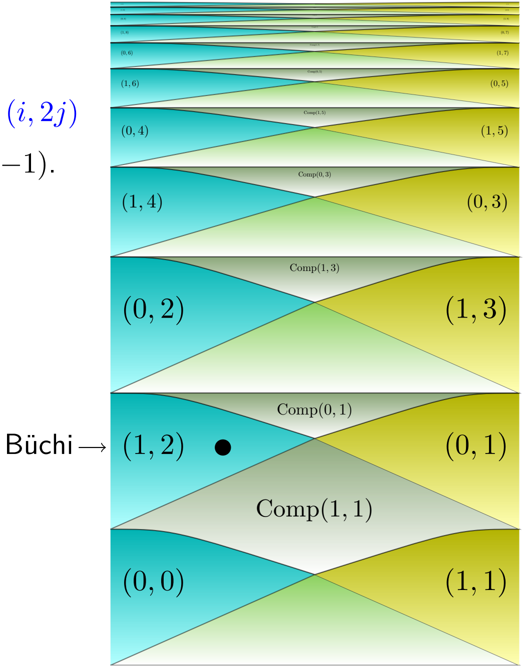

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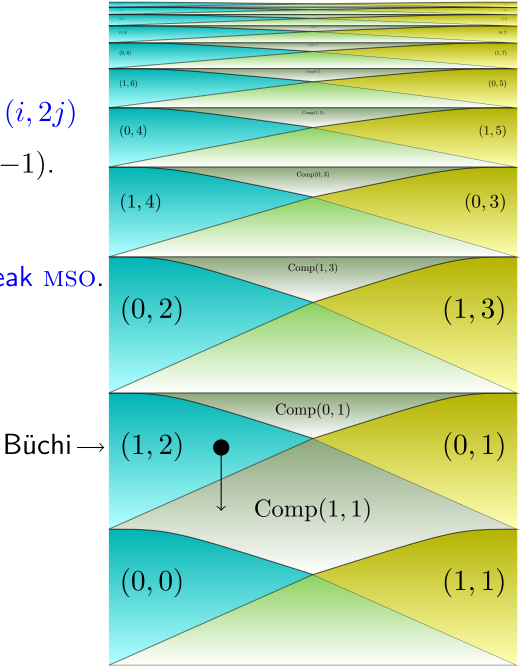

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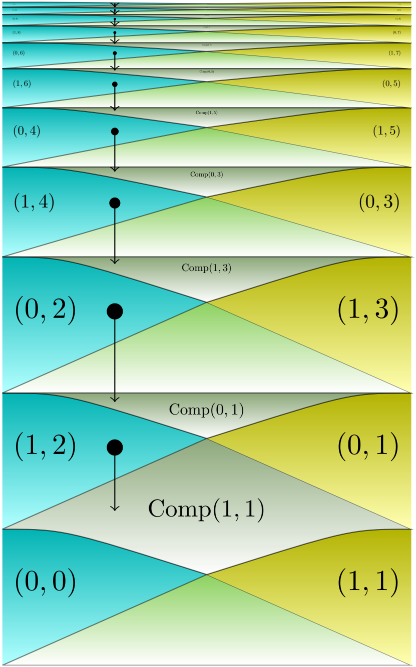

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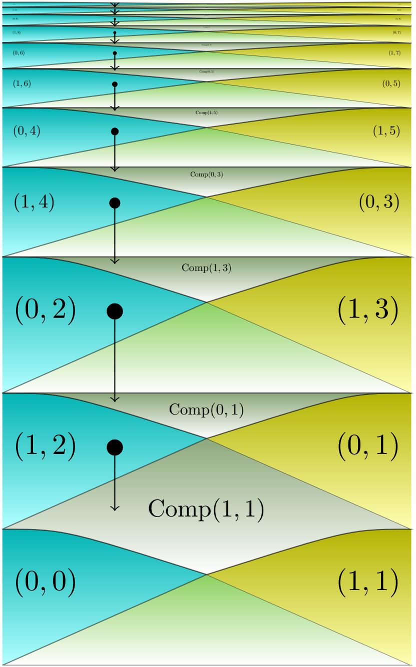

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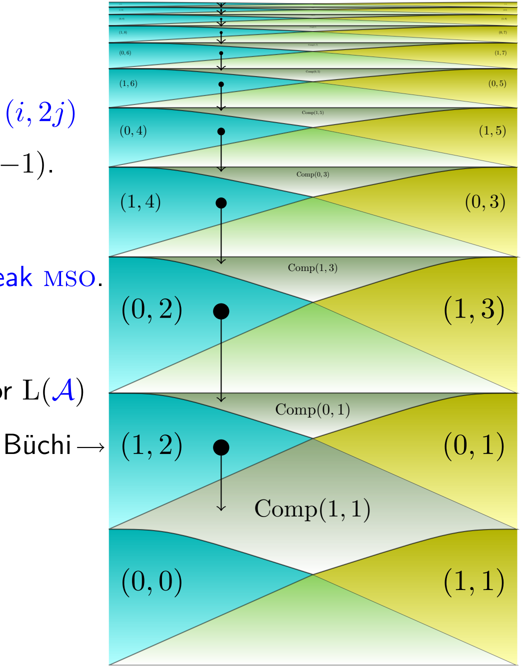

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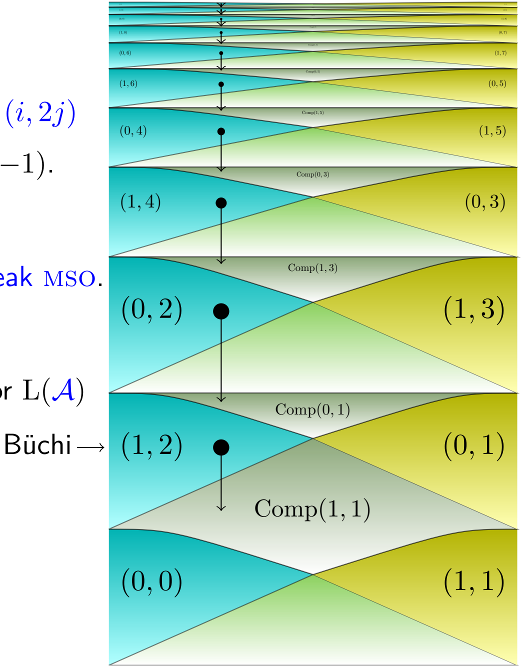

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