

Sequent calculus

Sequents have the form $\Gamma \longrightarrow \Delta$, where Γ, Δ are finite sequences of formulas. The intuitive meaning of a sequent $\Gamma \longrightarrow \Delta$ is: “the conjunction of all formulas in Γ implies the disjunction of all formulas in Δ ”.

Axiom

$$\frac{}{\varphi \longrightarrow \varphi} \text{ (axiom)}$$

Structural rules

$$\frac{\Gamma_1, \varphi, \psi, \Gamma_2 \longrightarrow \Delta}{\Gamma_1, \psi, \varphi, \Gamma_2 \longrightarrow \Delta} \text{ (exchange:L)}$$

$$\frac{\Gamma \longrightarrow \Delta_1, \varphi, \psi, \Delta_2}{\Gamma \longrightarrow \Delta_1, \psi, \varphi, \Delta_2} \text{ (exchange:R)}$$

$$\frac{\Gamma, \varphi, \varphi \longrightarrow \Delta}{\Gamma, \varphi \longrightarrow \Delta} \text{ (contraction:L)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi, \varphi}{\Gamma \longrightarrow \Delta, \varphi} \text{ (contraction:R)}$$

$$\frac{\Gamma \longrightarrow \Delta}{\Gamma, \varphi \longrightarrow \Delta} \text{ (weakening:L)}$$

$$\frac{\Gamma \longrightarrow \Delta}{\Gamma \longrightarrow \Delta, \varphi} \text{ (weakening:R)}$$

Rules for connectives

$$\frac{\Gamma, \varphi, \psi \longrightarrow \Delta}{\Gamma, \varphi \wedge \psi \longrightarrow \Delta} \text{ (\wedge:L)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi \quad \Gamma \longrightarrow \Delta, \psi}{\Gamma \longrightarrow \Delta, \varphi \wedge \psi} \text{ (\wedge:R)}$$

$$\frac{\Gamma, \varphi \longrightarrow \Delta \quad \Gamma, \psi \longrightarrow \Delta}{\Gamma, \varphi \vee \psi \longrightarrow \Delta} \text{ (\vee:L)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi, \psi}{\Gamma \longrightarrow \Delta, \varphi \vee \psi} \text{ (\vee:R)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi}{\Gamma, \neg \varphi \longrightarrow \Delta} \text{ (\neg:L)}$$

$$\frac{\Gamma, \varphi \longrightarrow \Delta}{\Gamma \longrightarrow \Delta, \neg \varphi} \text{ (\neg:R)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi \quad \Gamma, \psi \longrightarrow \Delta}{\Gamma, \varphi \Rightarrow \psi \longrightarrow \Delta} \text{ (\Rightarrow:L)}$$

$$\frac{\Gamma, \varphi \longrightarrow \Delta, \psi}{\Gamma \longrightarrow \Delta, \varphi \Rightarrow \psi} \text{ (\Rightarrow:R)}$$

The cut rule

$$\frac{\Gamma, \varphi \longrightarrow \Delta \quad \Gamma \longrightarrow \Delta, \varphi}{\Gamma \longrightarrow \Delta} \text{ (cut)}$$

Rules for quantifiers

$$\frac{\Gamma, \varphi(t/x) \longrightarrow \Delta}{\Gamma, \forall x \varphi \longrightarrow \Delta} \text{ (\forall:L)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi(y/x)}{\Gamma \longrightarrow \Delta, \forall x \varphi} \text{ (\forall:R)}$$

$$\frac{\Gamma, \varphi(y/x) \longrightarrow \Delta}{\Gamma, \exists x \varphi \longrightarrow \Delta} \text{ (\exists:L)}$$

$$\frac{\Gamma \longrightarrow \Delta, \varphi(t/x)}{\Gamma \longrightarrow \Delta, \exists x \varphi} \text{ (\exists:R)}$$

(restriction: in the (\forall :R) and (\exists :L) rules, the variable y does not appear in $\Gamma \cup \Delta \cup \{\varphi\}$)

Equality axioms (in first-order logic with equality)

$$\frac{}{\longrightarrow t = t} \text{ (AE1)}$$

$$\frac{}{t = s \longrightarrow s = t} \text{ (AE2)}$$

$$\frac{}{t = s, s = u \longrightarrow t = u} \text{ (AE3)}$$

$$\frac{}{t_1 = s_1, \dots, t_n = s_n, R(\bar{t}) \longrightarrow R(\bar{s})} \text{ (AE4)}$$

$$\frac{}{t_1 = s_1, \dots, t_n = s_n \longrightarrow f(\bar{t}) = f(\bar{s})} \text{ (AE5)}$$