

A note on the Σ_1 collection scheme and fragments of bounded arithmetic

Zofia Adamowicz*

Institute of Mathematics, Polish Academy of Sciences
Śniadeckich 8, 00-950 Warszawa, Poland
z.adamowicz@impan.gov.pl

Leszek Aleksander Kołodziejczyk*

Institute of Mathematics, University of Warsaw
Banacha 2, 02-097 Warszawa, Poland
lak@mimuw.edu.pl

December 19, 2008

Abstract

We show that for each $n \geq 1$, if T_2^n does not prove the weak pigeonhole principle for Σ_n^b functions, then the collection scheme $B\Sigma_1$ is not finitely axiomatizable over T_2^n . The same result holds with S_2^n in place of T_2^n .

The collection scheme $B\Sigma_1$ is

$$\forall v (\forall x < v \exists y \varphi(x, y) \Rightarrow \exists w \forall x < v \exists y < w \varphi(x, y))$$

for all bounded formulae φ (or equivalently, for all $\varphi \in \Sigma_1$, as the initial existential quantifiers may be absorbed by $\exists y$).

An intriguing open problem, mentioned already in [WP89], concerns the provability of $B\Sigma_1$ in $I\Delta_0 + \neg\text{exp}$. It is well known that $I\Delta_0$ does not

*The authors were partially supported by grant N N201 382234 of the Polish Ministry of Science and Higher Education.

prove $B\Sigma_1$, but all known proofs (cf. e.g. [PK78], [CFL07]) make use of the universal formula for Σ_1 , and hence need the totality of exponentiation. It is widely believed that $B\Sigma_1$ remains unprovable even if we assume $\neg \text{exp}$, but so far, no proof or even promising proof strategy has emerged. $B\Sigma_1$ is, however, known to be unprovable in $T_2^n + \neg \text{exp}$, where T_2^n is the finite fragment of Buss' S_2 (essentially a notational variant of $I\Delta_0 + \Omega_1$) axiomatized by induction for Σ_n^b formulae. Here, the universal formula for a restricted fragment of Σ_1 is enough.

The proofs of independence of $B\Sigma_1$ from $I\Delta_0 + \text{exp}$ and from $T_2^n + \neg \text{exp}$ have very much in common. In this note, we point out that the behaviour of $B\Sigma_1$ in these theories is nevertheless probably quite different. In $I\Delta_0 + \text{exp}$, collection is finitely axiomatizable, by the existence of the universal Σ_1 formula. On the other hand, we show that under a plausible assumption, $B\Sigma_1$ is not only unprovable, but even not finitely axiomatizable over T_2^n .

The “plausible assumption” we need is that T_2^n does not prove the weak pigeonhole principle $\text{WPHP}(\Sigma_n^b)$, i.e. that the existence of a Σ_n^b definable injection from a^2 to a for some $a > 1$ is consistent with T_2^n . Our proof goes through if T_2^n is replaced in both the assumption and the conclusion by the presumably weaker theory S_2^n . It is worth noting that until a breakthrough occurs, we cannot hope to prove non-finite axiomatizability of $B\Sigma_1$ unconditionally: it follows easily from [Bus95] and [Zam96] that if $T_2^n \vdash S_2^{n+1}$, then $B\Sigma_1$ is finitely axiomatized over T_2^n . The assumption about unprovability of WPHP seems reasonable, as it is true for all n in the relativized world ([BK94]) and, for $n = 1$ and S_2^1 , follows from the hardness of integer factoring ([Jeř07]).

Our result does have some bearing on the problem whether $I\Delta_0 + \neg \text{exp} \vdash B\Sigma_1$, in that it casts doubt on one possible line of attack. If $B\Sigma_1$ were finitely axiomatized over T_2^n for some n , the answer to the problem would be negative. This is because the unprovability of $B\Sigma_1$ in $T_2^m + \neg \text{exp}$ for each $m \geq n$ would imply unprovability of a fixed finite fragment, which would then be independent from $S_2 + \neg \text{exp}$ by compactness. However, if we are to believe the assumption about WPHP , then finite axiomatizability should not be hoped for.

We assume familiarity with basic notions and results concerning bounded arithmetic, which may be found e.g. in [HP93], [Kra95] or [Bus98]. For a brief review of relevant facts about WPHP , see e.g. [KT08] or [Tha02]. One important fact we need is that in S_2^n the failure of $\text{WPHP}(\Sigma_n^b)$ can be

amplified, that is, a Σ_n^b injection from a^2 into a can be used to obtain an injection from b into a for larger b .

We recall that the class $\hat{\Sigma}_n^b$, the prenex version of Σ_n^b , consists of formulae of the form

$$\exists y_1 < t_1 \forall y_2 < t_2 \dots Q y_n < t_n \psi,$$

where ψ is sharply bounded. The classes Σ_n^b and $\hat{\Sigma}_n^b$ coincide w.r.t. equivalence in S_2^n , but in weaker theories Σ_n^b might be strictly larger. $\hat{\Pi}_n^b$ is defined dually to $\hat{\Sigma}_n^b$, and $\exists \hat{\Pi}_n^b$ is the class of $\hat{\Pi}_n^b$ formulae preceded by existential quantifiers. It is easily checked that collection for $\exists \hat{\Pi}_n^b$ formulae, $B\exists \hat{\Pi}_n^b$, is equivalent to $B\hat{\Pi}_n^b$.

We also introduce one piece of notation: for a number a , $\#^m a$ is $a \# a \dots \# a$, where a appears m times. Given a model \mathcal{A} , $\#^{\mathbb{N}} a$ is the cut in \mathcal{A} determined by the numbers $\#^m a$ for standard m .

We now state and prove our theorem. Our proof is essentially a combination of slightly stronger variants of arguments from [AK07] and [KT08], and we assume the reader has access to those two papers.

Theorem 1. *Let T_n be T_2^n or S_2^n . If $T_n \not\vdash \text{WPHP}(\Sigma_n^b)$, then the collection scheme $B\Sigma_1$ is not finitely axiomatizable over T_n .*

We prove the theorem through a series of lemmas. Our starting point is a countable model $\mathcal{A} \models T_n$ containing a number a such that $\mathcal{A} = \#^{\mathbb{N}} a$ and the WPHP for Σ_n^b functions fails at a , in the sense that there is a Σ_n^b definable injection from a^2 into a . Such a model exists if T_n does not prove $\text{WPHP}(\Sigma_n^b)$.

To prove that $B\Sigma_1$ is not finitely axiomatizable over T_n , we need to show that there is no k such that over T_2^n , collection for Σ_1 formulae follows from collection for $\exists \hat{\Pi}_k^b$ formulae. W.l.o.g., we may consider only k for which $k + 2 \geq n + 1$.

There is a standard way of building a cofinal and $\hat{\Sigma}_{k+2}^b$ -elementary extension of \mathcal{A} to a $\hat{\Sigma}_{k+3}^b$ -maximal model \mathcal{A}_+ of T_n , i.e. one all of whose $\hat{\Sigma}_{k+2}^b$ -elementary extensions to models of T_n are actually $\hat{\Sigma}_{k+3}^b$ -elementary. By tweaking the construction a little, we are able to make \mathcal{A}_+ satisfy $\exists \hat{\Pi}_k^b$ collection (Lemma 2). By $\hat{\Sigma}_{n+1}^b$ -elementarity, $\text{WPHP}(\Sigma_n^b)$ still fails at a in \mathcal{A}_+ .

To complete the proof of the theorem, we show that \mathcal{A}_+ does not satisfy $B\exists \hat{\Pi}_{k+2}^b$. The argument is as follows. We observe that in \mathcal{A}_+ , each $\hat{\Sigma}_{k+3}^b$ formula is equivalent to a $\forall \hat{\Sigma}_{k+2}^b$ formula, with a as parameter (Lemma 3). If \mathcal{A}_+ satisfied $B\exists \hat{\Pi}_{k+2}^b$, this collapse would translate into a ‘‘local’’ collapse

of $\hat{\Sigma}_{k+2}^b$ to $\hat{\Pi}_{k+2}^b$ (Corollary 4), which falls just short of implying $\Sigma_{k+2}^p \subseteq \Pi_{k+2}^p/\text{poly}$. But in any model of S_2^n , such a collapse is incompatible with $\neg\text{WPHP}(\Sigma_n^b)$ (Lemma 5).

The remainder of the note contains proofs of the lemmas and a concluding remark.

Lemma 2. *Let $m \geq n$. Let $\mathcal{A} \models T_n$ be countable and of the form $\#^{\mathbb{N}}a$ for some a . There exists a cofinal countable extension $\mathcal{A}_+ \succeq_{\hat{\Sigma}_m^b} \mathcal{A}$ which is a $\hat{\Sigma}_{m+1}^b$ -maximal model \mathcal{A}_+ of T_n and (if $m \geq 2$) satisfies $\exists\hat{\Pi}_{m-2}^b$ collection.*

Proof. The construction of a cofinal $\hat{\Sigma}_m^b$ -elementary $\hat{\Sigma}_{m+1}^b$ -maximal extension of \mathcal{A} is a routine variant of the general model-theoretic construction of an existentially closed model. Starting with $\mathcal{A}_0 = \mathcal{A}$, we build a chain $\mathcal{A}_0 \preceq_{\hat{\Sigma}_m^b} \mathcal{A}_1 \preceq_{\hat{\Sigma}_m^b} \dots$ of countable cofinal models of T_n . \mathcal{A}_{l+1} arises from \mathcal{A}_l by adding a witness for the initial existential quantifier in a given $\hat{\Sigma}_{m+1}^b$ formula with a given choice of parameters from \mathcal{A}_l , whenever that is possible without losing $\hat{\Sigma}_m^b$ -elementarity. \mathcal{A}_+ is the union of the chain, and $\hat{\Sigma}_m^b$ -elementarity guarantees that \mathcal{A}_+ satisfies T_n . (See the proof of Lemma 2.2 in [AK07] for details.)

To ensure that \mathcal{A}_+ satisfies the right amount of collection, we slightly modify our method of constructing \mathcal{A}_{l+1} from \mathcal{A}_l . As before, we add a witness for a given $\hat{\Sigma}_{m+1}^b$ formula with given parameters in a $\hat{\Sigma}_m^b$ -elementary way. However, we also make sure that the model thus obtained, say $\tilde{\mathcal{A}}_{l+1}$, is *not* a cofinal extension of \mathcal{A}_l , but satisfies overspill for all bounded formulae. This can be achieved by a standard compactness argument. We then take \mathcal{A}_{l+1} to be the cut $\#^{\mathbb{N}}a$ in $\tilde{\mathcal{A}}_{l+1}$.

In this way, \mathcal{A}_{l+1} is a proper initial segment of the form $\#^{\mathbb{N}}a$ in a model of Σ_∞^b overspill. It is now easy to show $\mathcal{A}_{l+1} \models B\Sigma_1$. This is done by mirroring the well-known proof that a proper initial segment of a model of $I\Delta_0$ satisfies $B\Sigma_1$ (cf. [WP89]).

It remains to check that the fact that $\mathcal{A}_l \models B\Sigma_1$ for all l implies $\mathcal{A}_+ \models B\exists\hat{\Pi}_{m-2}^b$. Let $b, \bar{p} \in \mathcal{A}_+$, let $\psi(x, y, \bar{p})$ be a $\hat{\Pi}_{m-2}^b$ formula, and assume that

$$\mathcal{A}_+ \models \forall w \exists x < b \forall y < w \neg\psi(x, y, \bar{p}).$$

In particular, for each $i \in \mathbb{N}$ we have

$$\mathcal{A}_+ \models \exists x < b \forall y < \#^i a \neg\psi(x, y, \bar{p}).$$

Take l such that \mathcal{A}_l contains b and \bar{p} . By $\hat{\Sigma}_m^b$ -elementarity, we get

$$\mathcal{A}_l \models \exists x < b \forall y < \#^i a \neg \psi(x, y, \bar{p})$$

for each i , and thus

$$\mathcal{A}_l \models \forall w \exists x < b \forall y < w \neg \psi(x, y, \bar{p})$$

since \mathcal{A}_l is of the form $\#^{\mathbb{N}} a$. By $B\exists\hat{\Pi}_{m-2}^b$ in \mathcal{A}_l , there exists $c < b$ such that $\mathcal{A}_l \models \forall y \neg \psi(c, y, \bar{p})$. Applying $\hat{\Sigma}_m^b$ -elementarity once again, we obtain $\mathcal{A}_+ \models \forall y < \#^i a \neg \psi(c, y, \bar{p})$ for each i . But \mathcal{A}_+ is also of the form $\#^{\mathbb{N}} a$, which means that $\mathcal{A}_+ \models \forall y \neg \psi(c, y, \bar{p})$. \square

Lemma 3. *Let $m \in \mathbb{N}$ and let \mathcal{A} be a $\hat{\Sigma}_{m+1}^b$ -maximal model of T_n of the form $\#^{\mathbb{N}} a$. Then each $\hat{\Sigma}_{m+1}^b$ formula is equivalent in \mathcal{A} to a $\forall\hat{\Sigma}_m^b$ formula with a as an additional parameter.*

Proof. We sketch the proof omitting some details which are essentially the same as in Section 3 of [AK07].

It is easy to see that if \mathcal{A} is $\hat{\Sigma}_{m+1}^b$ -maximal for T_n , $\psi(x)$ is a $\hat{\Sigma}_{m+1}^b$ formula, and $d \in \mathcal{A}$, then $\psi(d)$ holds iff it is consistent with T_n plus the $\hat{\Pi}_m^b$ theory of \mathcal{A} in the language $L(\mathcal{A})$ (that is, L_{BA} expanded by constants for all elements of \mathcal{A}). Thus, it remains to check that “ $\psi(x)$ is consistent with T_n plus the $\hat{\Pi}_m^b$ theory of $\mathcal{A}_{L(\mathcal{A})}$ ” can be expressed in \mathcal{A} using a $\forall\hat{\Sigma}_m^b$ formula with a as a parameter.

Formalize $L(\mathcal{A})$ in some reasonable way, e.g. by letting the first few odd numbers represent the symbols of L_{BA} , letting $2d$ represent a constant symbol \underline{d} standing for $d \in \mathcal{A}$, and then coding syntax as usual. Our formula will say the following:

$$\begin{aligned} & \forall y \forall l \forall s [l \in \mathbb{N} \ \& \ y = 2^{|a|^l} \\ & \ \& \text{“}s \text{ is a sequence of formulae”} \ \& \ \sum_{i < \text{lh}(s)} \text{lh}((s)_i) \leq |l| \\ & \ \& \text{“no } (s)_i \text{ contains a constant for a number greater than } y\text{”} \\ & \ \& \ \forall i < \text{lh}(s) ((s)_i \in T_n \vee \text{“}(s)_i \text{ is a true } \hat{\Pi}_m^b \text{ formula”} \\ & \ \vee \text{“}(s)_i \text{ is derived from previous elements of } s \text{ by an inference rule”}) \\ & \Rightarrow (s)_{\text{lh}(s)-1} \neq \ulcorner \neg \psi(\underline{x}) \urcorner] \end{aligned}$$

We need to see that this is equivalent in \mathcal{A} to a $\forall \hat{\Sigma}_m^b$ formula, which amounts to checking that each conjunct in the antecedent of the implication may be stated in $\exists \hat{\Pi}_m^b$ form. The conjunct $l \in \mathbb{N}$ is not really needed, as it is implied by $y = 2^{|a|^l}$ because \mathcal{A} is of the form $\#^{\mathbb{N}}a$. The only other problematic conjunct is:

$$\forall i < \text{lh}(s) (\dots \vee \text{“(}s)_i \text{ is a true } \hat{\Pi}_m^b \text{ formula”} \vee \dots),$$

but a universal formula for $\hat{\Pi}_m^b$ formulae of length $\leq |l|$ and arguments below $2^{|a|^l}$ is $\hat{\Pi}_m^b$ with a bounding parameter, which can be any number above $2^{|a|^{l^2}}$. \square

Corollary 4. *Let \mathcal{A} be a $\hat{\Sigma}_{m+1}^b$ -maximal model of T_n of the form $\#^{\mathbb{N}}a$ and $\mathcal{A} \models B\exists \hat{\Pi}_m^b$. Then for each $d \in \mathcal{A}$, each $\hat{\Sigma}_{m+1}^b$ formula is equivalent on $[0, d]$ to a $\hat{\Pi}_{m+1}^b$ formula with a as an additional parameter.*

Proof. Let $\psi(x)$ be a $\hat{\Sigma}_{m+1}^b$ formula and $d \in \mathcal{A}$. By Lemma 3, $\psi(x)$ is equivalent in \mathcal{A} to $\forall y \eta(x, y, a)$, where η is $\hat{\Sigma}_m^b$. Thus, we have

$$\mathcal{A} \models \forall x (\psi(x) \vee \exists y \neg \eta(x, y, a)).$$

If $\mathcal{A} \models B\exists \hat{\Pi}_m^b$, then there exists w such that for $x \in [0, d]$, a witness for either $\psi(x)$ or $\exists y \neg \eta(x, y)$ may be bounded by w . We may take w to be of the form $\#^l a$ for some $l \in \mathbb{N}$, so on $[0, d]$, $\psi(x)$ is equivalent to $\forall y < \#^l a \eta(x, y, a)$. \square

Lemma 5. *Let $\mathcal{A} \models S_2^n + \neg \text{WPHP}(\Sigma_n^b)$ be of the form $\#^{\mathbb{N}}a$. Then for each $m \geq 1$ there exists $d \in \mathcal{A}$ and a $\hat{\Sigma}_m^b$ formula $\psi(x)$ which is not equivalent on $[0, d]$ to a $\hat{\Pi}_m^b$ formula, even with parameters.*

Proof. The proof is based on an argument from Sections 4 and 5 of [KT08], used there to show that in a model of $S_2^n + \neg \text{WPHP}(\Sigma_n^b)$ the bounded formula hierarchy does not collapse, even with parameters (Theorem 5.1 of [KT08]). We will check that the argument is actually strong enough to show that in the case of models of the form $\#^{\mathbb{N}}a$, failure of $\text{WPHP}(\Sigma_n^b)$ excludes even a collapse on a large enough initial segment with a top.

Assume that m is such that on each interval $[0, d]$, every $\hat{\Sigma}_m^b$ formula is equivalent to a $\hat{\Pi}_m^b$ formula with a parameter (which may depend on d). It can then be easily checked that on each $[0, d]$, every bounded formula is equivalent to a $\hat{\Pi}_m^b$ formula with a parameter.

By increasing a and amplifying the function violating WPHP if necessary, we may assume that a is a power of 2, f is a Σ_n^b injection from $a\#a$ into a and that the definition of f involves only a parameter $q < a$ and quantifiers bounded by $a\#a$. We now use compactness to extend \mathcal{A} elementarily to a model \mathcal{A}' additionally containing an element $t > \mathcal{A}$ and an element $b > \#^{\mathbb{N}}t$ of the form $\#^c a$ for some small nonstandard c . Let \mathcal{B} be the cut $\#^{\mathbb{N}}a$ in \mathcal{A}' . The reason for the amplification and for introducing the new models is that the relation between \mathcal{A}' and \mathcal{B} is now exactly the same as between the models \mathcal{A} and \mathcal{B} in Section 4 of [KT08], and we will be able to apply the results of that section.

Note that because $\mathcal{A}' \succeq \mathcal{A}$ and \mathcal{A} is cofinal in \mathcal{B} , on each interval $[0, d]$ in \mathcal{B} every bounded formula is equivalent to a $\hat{\Sigma}_m^b$ formula with a parameter from \mathcal{B} (even from \mathcal{A}).

Since $t > \mathcal{B}$, there is a universal $\hat{\Sigma}_m^b$ formula U_m such that for all $x, y \in \mathcal{B}$ and $\hat{\Sigma}_m^b$ formulae ψ , $\psi(x, y)$ is equivalent in \mathcal{A}' to $U_m(x, \langle \ulcorner \psi \urcorner, y \rangle, t)$. Now, $b > \#^{\mathbb{N}}t$ and U_m is bounded, so for $x, y \in \mathcal{B}$ and for standard ψ the quantifiers in $U_m(x, \langle \ulcorner \psi \urcorner, y \rangle, t)$ range only over numbers below b . By Lemma 4.3 of [KT08], this means that we can use the failure of WPHP to translate $U_m(x, \langle \ulcorner \psi \urcorner, y \rangle, t)$ into a bounded formula with parameters from \mathcal{B} . More precisely, there is a bounded (even linearly bounded) formula U_m^{lin} such that for all $x, y \in \mathcal{B}$ and all ψ , $U_m(x, \langle \ulcorner \psi \urcorner, y \rangle, t)$ is equivalent to $U_m^{\text{lin}}(x, \langle \ulcorner \psi \urcorner, y \rangle, p)$, where p is a parameter bounded by some standard power of $a\#a$ (p is actually a tuple $(\hat{t}, a\#a, c, q)$, where c, q are as above and \hat{t} is a number below a which codes t in a certain way).

On each interval $[0, d]$ in \mathcal{B} , the bounded formula $\neg U_m^{\text{lin}}(x, x, p)$ is equivalent to a $\hat{\Sigma}_m^b$ formula $\varphi(x, r)$. A priori, the size of r depends on d , but for sufficiently large d , we can assume $\langle \ulcorner \varphi \urcorner, r, p \rangle < d$. This is because r is certainly bounded by $\#^i a$ for some i , and $\#^i a$ can be mapped injectively into a by an amplified version $f^{(i-1)}$ of f , which is also Σ_n^b definable; thus, we can replace the original $\varphi(x, r)$ by $\tilde{\varphi}(x, \tilde{r}) := \exists y < \#^m a (f^{(i-1)}(y) = \tilde{r} \ \& \ \varphi(x, y))$, where $\tilde{r} = f^{(i-1)}(r)$ is a number below a .

Consider $\varphi(\langle \ulcorner \varphi \urcorner, r \rangle, r)$. By the properties of U_m , this is equivalent to $U_m(\langle \ulcorner \varphi \urcorner, r \rangle, \langle \ulcorner \varphi \urcorner, r \rangle, t)$ and hence to $U_m^{\text{lin}}(\langle \ulcorner \varphi \urcorner, r \rangle, \langle \ulcorner \varphi \urcorner, r \rangle, p)$. On the other hand, for a large enough d , $\neg U_m^{\text{lin}}(\langle \ulcorner \varphi \urcorner, r \rangle, \langle \ulcorner \varphi \urcorner, r \rangle, p)$ is also equivalent to $\varphi(\langle \ulcorner \varphi \urcorner, r \rangle, r)$, which gives a contradiction. \square

Remark. The model \mathcal{A}_+ obtained in our construction has the property that each $\exists \hat{\Pi}_{k+3}^b$ formula is equivalent to an $\exists \hat{\Pi}_{k+2}^b$ formula with a parameter

(Lemma 3), but is not in general equivalent to an $\exists\hat{\Pi}_k^b$ formula, even with parameters (because collection holds for the latter class).

References

- [AK07] Z. Adamowicz and L. A. Kołodziejczyk, *Partial collapses of the Σ_1 complexity hierarchy in models for fragments of bounded arithmetic*, Annals of Pure and Applied Logic **145** (2007), 91–95.
- [BK94] S. R. Buss and J. Krajíček, *An application of boolean complexity to separation problems in bounded arithmetic*, Proceedings of the London Mathematical Society **s3-69** (1994), 1–21.
- [Bus95] S. R. Buss, *Relating the bounded arithmetic and polynomial time hierarchies*, Annals of Pure and Applied Logic **75** (1995), 67–77.
- [Bus98] ———, *First-order proof theory of arithmetic*, Handbook of Proof Theory (S. R. Buss, ed.), Elsevier, 1998, pp. 79–147.
- [CFL07] A. Cordon Franco, A. Fernández Margarit, and F. F. Lara Martín, *A note on Σ_1 -maximal models*, Journal of Symbolic Logic **72** (2007), 1072–1078.
- [HP93] P. Hájek and P. Pudlák, *Metamathematics of first-order arithmetic*, Springer-Verlag, 1993.
- [Jeř07] E. Jeřábek, *On independence of variants of the weak pigeonhole principle*, Journal of Logic and Computation **17** (2007), 587–604.
- [Kra95] J. Krajíček, *Bounded arithmetic, propositional logic, and complexity theory*, Cambridge University Press, 1995.
- [KT08] L. A. Kołodziejczyk and N. Thapen, *The polynomial and linear hierarchies in models where the weak pigeonhole principle fails*, Journal of Symbolic Logic **73** (2008), 578–592.
- [PK78] J. B. Paris and L. A. S. Kirby, *Σ_n collection schemas in arithmetic*, Logic Colloquium '77, Studies in Logic and the Foundations of Mathematics, vol. 96, North Hollandg, 1978, pp. 199–209.

- [Tha02] N. Thapen, *A model-theoretic characterization of the weak pigeonhole principle*, *Annals of Pure and Applied Logic* **118** (2002), 175–195.
- [WP89] A. J. Wilkie and J. B. Paris, *On the existence of end-extensions of models of bounded induction*, *Logic, Methodology, and Philosophy of Science VIII (Moscow 1987)* (J.E. Fenstad, I.T. Frolov, and R. Hilpinen, eds.), North-Holland, 1989, pp. 143–161.
- [Zam96] D. Zambella, *Notes on polynomially bounded arithmetic*, *Journal of Symbolic Logic* **61** (1996), 942–966.